

**STUDIES IN SET-THEORETIC HIERARCHIES:
FROM BOREL SETS TO R-SETS**

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NOTATION

ω :	set of natural numbers
ω^ω :	sequence of natural numbers called reals
$\omega^{<\omega}$:	finite sequence of natural number
Seq :	set of Gödel numbers of finite sequence of natural numbers
η, ξ, \dots	usually denote subsets of ω
$\alpha, \beta, \gamma, \dots$	usually denote reals
s, t, u, v, \dots	usually vary over Seq
$\langle a_0, \dots, a_{n-1} \rangle$	Gödel number or sequence number of (a_0, \dots, a_{n-1}) , sometimes identified with (a_0, \dots, a_{n-1}) .
$s^*t, s^{\wedge}t$	concatenation of two sequence numbers or concatenation of two finite sequences; otherwise taken to be 0.
sn	$s^*\langle n \rangle$
ns	$\langle n \rangle * s$
$s \sqsubseteq t$	s is an initial segment of t if $s, t \in \text{Seq}$ or $\omega^{<\omega}$
$s^*\alpha$	$(a_0, a_1, \dots, a_{n-1}, \alpha(0), \alpha(1), \dots)$ if $s = \langle a_0, \dots, a_{n-1} \rangle$
$lh(s)$	length of the finite sequence or sequence number s
χ_A	characteristic function of A
$\mathcal{P}(A)$	power set of A
$N_s, \Sigma(s)$	set of reals extending $s \in \text{Seq}$
$\bar{\alpha}(n)$	$\langle \alpha(0), \dots, \alpha(n-1) \rangle$
$E^x(E_y)$	horizontal (vertical) section of E at x .
$A^c, \neg A, \sim A$	complement of A
$e, \langle \rangle$	code of the empty sequence

$A B$	$A - B$
$\text{Dom}(f)$	domain of f
$f(x) \downarrow$	$x \in \text{Dom}(f)$
$c \mathcal{F}, \neg \mathcal{F}$	family of complements of sets in \mathcal{F}
B_A	Borel σ -field on A
$\sigma(\mathcal{F})$	σ -field generated by \mathcal{F}
\leq_α	$\{(m, n) : \alpha(\langle m, n \rangle) = 0\}$
WO	set of reals α such that \leq_α is a wellordering on ω

For unexplained notation and terminology from descriptive set theory see Moschovakis [38]. Notation and terminology from recursion theory is from Hinman [22].

SUMMARY OF THESIS

The present thesis consists of three parts.

Part I of the thesis is devoted to the study of regularity properties of classes of sets obtained from δ -s operations applied to the class of clopen sets. In particular, the classes of sets obtained by applying the hierarchy of increasingly complex R-operations $\{R_\alpha\}$, are investigated. Our points of departure in this study are the recent works on R-sets by Hinman and Burgess. In this thesis we have unified the methods of Hinman and Burgess by extensively using the game quantifier and methods of inductive definability. Using these tools we are able not only to considerably simplify the proofs of known results in the theory (e.g. the comparison of indices theorem on which so much of Lyapunov's work depends as well as the category and measure formulas for R-sets) but also to obtain a number of new results for R-sets. Among these new results we mention the measurable selection property for different levels of the hierarchy of R-sets, the scale property at all levels of the R-hierarchy (this was known only for the first two levels) and the uniformization property at all levels of the R-hierarchy (again this was known only for the first level). These last mentioned results depend on the notion of presentability introduced by Burgess in his penetrating study of the measurable selection property of various families of

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sets. Our proof turns on the observation that the presentability property holds at every level of the R -hierarchy. Finally, by combining analysis of the associated games with an adaptation of the Kechris Game Formula and the existence of measurable winning strategies, we are able to extend a recent result of Srivatsa on C -sets to R -sets. Briefly stated this result ensures that an R -set in the product of two Polish spaces can be approximated in the sense of category by means of sets in appropriate product σ -fields, the importance of this result resting on the fact that sets in product σ -field have a much more tractable structure than R -sets. Following Srivatsa we use this result to obtain a selection theorem of Burgess for R -sets.

While Part I is devoted to the study of classical (bold-face) R -sets, in Part II we take up the study of the effective analogues of these classes at the finite levels. Using again the game theoretic methods of Moschovakis we are able to establish various regularity properties of these effective classes. For instance, we show that each cR_n is a Spector pointclass. We also observe that various other results like the Measure and Category Formulas, the scale property hold in the effective case also. We also take another look at the problem of effectivizing C -sets. Hinman, by generalizing Addison's construction of the effective

Borel hierarchy, obtained a hierarchy of effective C -sets on ω which exhaust ${}^1sc(\mathbb{E}_1)$, the class of subsets of ω recursive in \mathbb{E}_1 . We, however, obtain an effective hierarchy in ${}^1sc(\mathbb{E}_1)$ in a different manner. We first exhibit a set $D(T, \alpha, x)$, where T varies over wellfounded trees on ω , which is universal for C -sets $\uparrow X$. When $X = \omega$, by taking sections of D by wellfounded trees and reals of increasing complexity, we obtain a hierarchy whose scope is ${}^1sc(\mathbb{E}_1)$.

Part III of our thesis is different in spirit and content. This part is divided into two sections. In the first section we work with hereditary families of subsets of reals satisfying a weak definability condition and think of members of these families as small sets. We show that, at the first level of the analytical hierarchy, a number of measure-theoretic and category-theoretic results of Kechris and others are instances of more general results much in the spirit of Kechris' Basis Theorem. We also obtain the Basis Theorem of Kechris for large Π_1^1 sets under a weaker definability condition. For the basis theorem we need to consider σ -ideals instead of hereditary families. Some σ -ideals we come across in descriptive theory, like the σ -ideals of meager sets and of null sets, satisfy both the Σ_1^1 and Π_1^1 computational formulas. However, we show that

for a σ -ideal which is generated by a Π_1^1 -coded ideal of closed sets, small Σ_1^1 sets have the covering property and, consequently, the Σ_1^1 computational formula holds.

In Section 2, we undertake a study of the structure of hyperarithmetical sets which are of ambiguous Borel classes. There is a natural hierarchy, called the difference hierarchy, for Borel sets of ambiguous class $\xi+1$ ($\Delta_{\xi+1}^0$). We show that there is an analogous effective hierarchy of $\Delta_1^1 \cap \Delta_{\xi+1}^0$ which coincides with the difference hierarchy of $\Delta_{\xi+1}^0$ restricted to Δ_1^1 . (Earlier, Louveau has obtained similar results in a more general setting in his paper in Cabal Seminar, 1979-81). As a consequence we have been able to obtain a complete generalization of a recent result of Debs for every level of the difference hierarchy. Debs proved that a Borel set in the product of two Polish spaces, with vertical sections the difference of two G_δ -sets, is the difference of two Borel sets with G_δ sections. The effective result mentioned above enables us to extend Debs' result. We obtain this effective result via a separation theorem for Σ_1^1 sets. From another separation theorem, we are able to establish the faithful separation property of Δ_ξ^0 for every ξ . This extends Burgess' result to every level. Finally, we are able to establish the non-existence of Borel sets universal for $\Delta_{\xi+1}^0$, $\xi < \omega_1$ — an unpublished result of Addison and Harrington.

INTRODUCTION TO PART I

Descriptive set theory is concerned with the study of those sets of real numbers which are of interest and significance for analysis. The first large class of sets studied was the class of Borel sets which was classified, according to the degree of complexity, into a hierarchy of ω_1 levels. The discovery of operation \mathcal{A} and the analytic sets (which result from application of \mathcal{A} to countable sequences of open sets) led to the definition of \mathcal{C} -sets and their classification into a hierarchy [43]. The importance of this discovery can be gauged by the fact that it led directly to Hausdorff's δ -s operations and to the general notion of a set-theoretic operation. It was then natural to search for set-theoretic operations more powerful than operation \mathcal{A} which would allow a still further enlargement of the class of sets for which an "explicit construction" could be given. This led to the discovery of the operator \mathcal{R} by Kolmogorov [23]. However, it was only years later that a systematic application of it was made to enlarge the class of explicitly constructible sets [33].

In [23], Lyapunov constructs for each $\rho < \omega_1$ a positive analytical set operation \mathcal{R}_ρ and the class of sets \mathcal{R}_ρ consisting of those sets obtained by the application of \mathcal{R}_ρ to sequences of open sets. We have $\mathcal{R}_0 = \mathcal{A}$ and hence \mathcal{R}_0 = class of analytic sets; while $\mathcal{R}_1 \cap \subset \mathcal{R}_1$ properly includes the

class of C -sets. Members of $\mathcal{R} = \bigcup_{\rho < \omega_1} \mathcal{R}^\rho$ are classically known as the R -sets. \mathcal{R} was also shown to consist of measurable sets and to be a proper subclass of Δ_2^1 .

Although the theory of R -sets and the \mathcal{R} -operator has been studied extensively by Russian mathematicians [23],[33],[34], and most of the basic properties have been deduced by them; it is only very recently that interest in the theory has been revived due mainly to the work of Hinman [20,21] who developed the effective counterpart and showed that the effective hierarchies have deep interconnections with recursion-theoretic hierarchies. The work of Burgess [12,13] has added a new dimension to the theory. He has shown that the entire hierarchy of R -sets can be obtained by applying the game-quantifier to the "difference hierarchy" of Δ_3^0 (obtained via decreasing sequences of G_δ -sets). Burgess and Lockhart [14] have also shown that the class of sets obtained from Borel-programmable sets (of Blackwell cf.[7]) by "iteration" gives precisely the R -sets. That seemingly different definitions yield the same class of sets suggests that the R -sets form a natural class of subsets of the reals.

The first few sections of Part I will be mainly devoted to the study of R -sets and the \mathcal{R} -operator. In fact, we shall study these in much greater generality by considering the classes $\Sigma_1^{\Phi^*}$, $\Pi_1^{\Phi^*}$, $\nabla(\Phi^*)$ etc. much in the spirit of Hinman (cf.[20]). Thus

we have tried to unify the work of Burgess and Hinman, and have formulated many results appropriate for our purpose in, perhaps, a more unified manner than is available in the literature. For instance, we have obtained the prewellordering (p.w.o) property of $\mathcal{C}\mathcal{R}^0$ by showing that for most operations Φ , $\Pi_1^{\Phi^*}$ has the p.w.o. property (Section 3). This is proved via a comparison of indices lemma whose proof is much along the lines of the Kunen-Aczel theorem (cf. [37], see also [34]). The comparison of indices lemma is crucial for our purpose since it helps in computing the complexity of the winning strategy for the existential player in the game associated with the operator \mathcal{R} .

In Section 4, we have obtained the partition-selection property of \mathcal{R} -sets by showing, à la Burgess, that each $\mathcal{V}(\mathcal{R}_\rho)$ is presentable. This we exhibit by showing that the property of being presentable is "transferred" by the operator \mathcal{R} . The notion of presentability not only yields the partition-selection property but also the scale and uniformization properties of \mathcal{R} -sets. We also obtain a decomposition of $E^* = \{x : E^x \text{ is comeager}\}$ for sets $E \in \Sigma_1^{\Phi^*}$, analogous to the one obtained by Vaught for Σ_1^1 sets [51] and for \mathcal{R}^1 -sets obtained by Burgess and Miller [8]. This immediately gives us the transfer theorem, which essentially computes E^* for sets $E \in \Sigma_1^{\Phi^*}$, whenever the computation for F^* for sets $F \in \Sigma_1^{\Phi}$

is known. This is done in Section 6. An immediate application of the transfer theorem is the computation of E^* for R -sets E — a result due to Burgess and Miller [8]. These computations yield sets in product σ -fields which "approximate" section-wise (in the sense of category) R -sets in two dimensions. This is done in Section 7. Incidentally, in this section, we also prove a slightly stronger version of the Game Formula of Kechris (cf. [26]).

PART I

R-SETS

1. Positive analytical operations : In the section we shall introduce positive analytical operations and Hausdorff's δ -s operations and discuss some of their elementary properties which we will need in the sequel. The papers of Kantorovitch and Livenson [23] give a detailed exposition of these operations.

Let X be a non-empty set and $N \subseteq \mathcal{P}(\omega)$. (As usual, we shall identify $\mathcal{P}(\omega)$ with 2^ω).

1.1 Definition : The δ -s operation with base N is defined by

$$\mathbb{I}_N(\{E_n : n \in \omega\}) = \bigcup_{\eta \in N} \bigcap_{n \in \eta} E_n,$$

where $\{E_n : n \in \omega\}$ is any family (sequence) of subsets of X .

To avoid trivialities we shall always assume $\emptyset \notin N$ and $N \neq \emptyset$.

Examples : If $N = \{\{n\} : n \in \omega\}$, then $\mathbb{I}_N = U$ (countable).

If $N = \{\eta \subseteq \omega : \eta \text{ is infinite}\}$, then $\mathbb{I}_N = \text{linsup}$.

If $N = \{\bar{\alpha}(n) : n \in \omega : \alpha \in \omega^\omega\}$, then $\mathbb{I}_N = \text{operation } A$.

By an operation we shall mean a function $\mathbb{I} : \mathcal{P}(X)^\omega \rightarrow \mathcal{P}(X)$ for all X .

1.2 Definition : An operation \mathbb{I} is said to be a positive analytical operation (p.a.o) if

(a) \mathbb{I} is non-constant and

(b) for any families $\{E_n : n \in \omega\}$ and $\{F_n : n \in \omega\}$ of subsets of a set X ,

$$x \in \Phi(\{E_n : n \in \omega\}) \& y \notin \Phi(\{F_n : n \in \omega\}) \rightarrow (\exists n)[x \in E_n \& y \notin F_n].$$

Clearly, $U, \mathcal{A}, \text{limesup}$ are positive analytical operations. A p.a.o. Φ on constant sequences takes on the constant value and is isotone i.e. for any families $\{F_n : n \in \omega\}$ and $\{G_n : n \in \omega\}$ if $(\forall n)(F_n \subseteq G_n)$, then $\Phi(\{F_n : n \in \omega\}) \subseteq \Phi(\{G_n : n \in \omega\})$.

Notice that a δ -s operation is a p.a.o. The converse is true and follows from the following.

Proposition: Let Φ be a p.a.o. Then $\Phi = \Phi_N$ for some $N \subseteq \mathcal{P}(\omega)$.

Proof: Let $D_i = \{\eta \subseteq \omega : i \in \eta\}$. Put $N = \Phi(\{D_i : i \in \omega\})$. It is easy to check that $\Phi = \Phi_N$.

The set N obtained above is called the canonical base for Φ .

1.3 Definition: A base $N \subseteq \mathcal{P}(\omega)$ is complete if

$$\eta \in N \& \eta \subseteq \eta' \subseteq \omega \rightarrow \eta' \in N.$$

If $N = \Phi_N(\{D_i : i \in \omega\})$, where D_i is as above, then \bar{N} is complete (called the completion of N) and $\Phi_N = \Phi_{\bar{N}}$. The canonical base for a p.a.o. is therefore complete.

1.4 Definition: For any operation Φ , the dual Φ° is defined by:

$$\Phi^\circ(\{E_n : n \in \omega\}) = [\Phi(\{E_n^\circ : n \in \omega\})]^\circ.$$

For example $U^\circ = \Pi$, $(\text{limesup})^\circ = \text{liminf}$, $\mathcal{A}^\circ = \mathcal{J}$ (the Suslin operation),

where $\mathcal{J}(\{E_n : n \in \omega\}) = \{x : (\forall \alpha)(\exists n)(x \in E_{\bar{\alpha}(n)})\}$.

If $\bar{\Phi}_N$ is a δ -s operation with base N , then the canonical base N° of its dual is given by

$$\begin{aligned} N^\circ &= \{ \eta \in 2^\omega : \eta \cap \eta' \neq \emptyset \text{ for every } \eta' \in N \} \quad \dots(1) \\ &= \{ \eta \in 2^\omega : \eta^\circ \notin N \}, \text{ if } N \text{ is complete} \end{aligned}$$

Plainly, N° is always complete. Thus, for any family $\{E_n : n \in \omega\}$,

$$(\exists \eta \in N^\circ)(\forall n \in \omega) [x \in E_n] \leftrightarrow (\forall n \in \omega)(\exists \eta \in N)(\eta \cap E_n \neq \emptyset). \quad \dots(2)$$

If N is complete, $N^{\circ\circ} = N$ and hence

$$(\forall \eta \in N^\circ)(\exists n \in \omega) [x \in E_n] \leftrightarrow (\exists \eta \in N)(\forall n \in \omega) [x \in E_n]. \quad \dots(3)$$

2. The \mathbb{R} -operator : Although the \mathbb{R} -operator was first introduced by Kolmogorov, the first published account of the theory appeared in [23] and further results obtained in [33, 34]. Lyapunov also studied the hierarchies of \mathbb{R} -sets (cf. § 3) and obtained most of their properties. The interconnection between the \mathbb{R} -operator and games was first noticed by Hinman [20] (and also independently by Aczel [1]). Hinman also developed the effective theory and did most of the groundwork. Much of the material in this section is adapted from these sources.

2.1 Definition : Let $\mathcal{N} = \{N_p \subseteq 2^\omega : p \in \omega\}$ be a sequence of nonempty bases. $\theta \subseteq \omega$ is called an \mathcal{N} -chain if for any $s, t \in \text{Seq}$,

$$(a) \quad s \in \theta,$$

$$(b) \quad s \in \theta \text{ and } t \subseteq s \rightarrow t \in \theta,$$

$$(c) \quad s \in \theta \rightarrow \{n : s \prec n \in \theta\} \in N_s.$$

Put $\textcircled{H} = \{\theta : \theta \text{ is an } \mathcal{N}\text{-chain}\}.$

\mathcal{R}_X is the set operation defined by

$$\mathcal{R}_X(\{E_n : n \in \omega\}) = \bigcup_{\theta \in \mathbb{H}} \bigcap_{s \in \theta} E_s.$$

Clearly, \mathcal{R}_X is a δ -s operation with base \mathbb{H} .

If $N_p = N$ for each p , then \mathcal{R}_X is denoted by $\mathbb{R}\bar{\Phi}_N$ and its base by $\mathbb{R}N$. An X -chain will then be called an N -chain.

Example: Let $N = \{ \{n\} : n \in \omega \}$ so that $\bar{\Phi}_N = U$ and put $X = \{N\}$. Clearly, an X -chain contains a set of the form $\{ \bar{a}(n) : n \in \omega \}$. Thus $\mathbb{R}U = \mathbb{R}\bar{\Phi}_N = \mathcal{A}$.

2.2 Definition: Let $\bar{\Phi}_N$ and $\bar{\Phi}_M$ be two δ -s operations with bases $N, M \subset 2^\omega$. The composed operation Ψ is given by

$$\Psi(\{F_n : n \in \omega\}) = \bar{\Phi}_N(\{\bar{\Phi}_M(\{F_{\langle p, n \rangle} : n \in \omega\}) : p \in \omega\}).$$

Ψ is sometimes denoted by $\bar{\Phi}_N \bar{\Phi}_M$. By the characterization lemma, Ψ is a δ -s operation whose canonical base we shall denote by NM . Thus $\Psi = \bar{\Phi}_N \bar{\Phi}_M = \bar{\Phi}_{NM}$ and

$$n \in NM \iff (\exists n_1 \in N)(\forall n_1 \in n_1)(\exists \xi_1 \in M)(\forall m_1 \in \xi_1) [\langle n_1, m_1 \rangle \in n].$$

Henceforth, for convenience, we shall work in ω^ω (or, $(\omega^\omega)^k \times \omega^\lambda$), the space of reals, since this can be identified with the irrationals in $(0,1)$. Henceforth we shall denote such spaces by X, Y unless otherwise stated.

2.3 Definition: For any operation $\bar{\Phi}$, let $\Sigma_1^{\bar{\Phi}}$ be the class of relations of the form $\bar{\Phi}(\{F_n : n \in \omega\})$ with all F_n clopen, $\Pi_1^{\bar{\Phi}}$ the class

of complements of such relations and $\Delta_1^{\bar{\Phi}} = \Sigma_1^{\bar{\Phi}} \cap \Pi_1^{\bar{\Phi}}$. Thus $\Sigma_1^U = \Sigma_1^{\bar{\Phi}}$ and $\Sigma_1^A = \Sigma_1^1$, the class of analytic sets. If $\bar{\Phi} = \bar{\Phi}_N$, then put $\bar{\Phi}^* = \mathbb{R}\bar{\Phi}_{\mathbb{N}\mathbb{N}^0}$.

By taking F_n to be recursive uniformly in n one obtains the corresponding lightface classes $\Sigma_1^{\bar{\Phi}}$, $\Pi_1^{\bar{\Phi}}$ etc.

The next two lemmas show a close connection between \mathbb{R} -operators, inductive definability and games.

2.4 Lemma : (a) Suppose $F = \mathbb{R}\bar{\Phi}_N(\{F_n : n \in \omega\})$. Then

$$x \in F \leftrightarrow (\exists \eta_0 \in N)(\forall n_0 \in \eta_0)(\exists \eta_1 \in N)(\forall n_1 \in \eta_1) \dots \dots (\forall k) [x \in F_{\langle n_0, n_1, \dots, n_{k-1} \rangle}] \quad (*)$$

(b) If $E = \bar{\Phi}_N^*(\{E_n : n \in \omega\})$, then

$$x \in E \leftrightarrow (\exists \eta_0 \in N)(\forall n_0 \in \eta_0)(\forall \xi_0 \in N)(\exists m_0 \in \xi_0)(\exists \eta_1 \in N)(\forall n_1 \in \eta_1) \dots (\forall \xi_1 \in N)(\exists m_1 \in \xi_1) \dots (\forall k) [x \in E_{\langle \langle n_0, m_0 \rangle, \dots, \langle n_{k-1}, m_{k-1} \rangle \rangle}]$$

(The right-hand side of each equivalence is interpreted in terms of games between two players \forall and \exists . For instance, the right-hand expression of (*) is true just in case \exists has a winning strategy in the game G_x played as follows : \exists (player I) chooses $\eta_0 \in N$, then \forall (player II) chooses $n_0 \in \eta_0$, then \exists chooses $\eta_1 \in N$, etc. Thus a strategy for \exists is a function from Seq into $\mathcal{P}(\omega)$, while a strategy for \forall is a function from $\cup_k (2^\omega)^k \rightarrow \omega$. If $\eta_0, n_0, \eta_1, \dots$ is the sequence produced when both players have played as described above, then \exists wins iff

$$(\forall k) [x \in F_{\langle n_0, n_1, \dots, n_{k-1} \rangle}]$$

Proof : Clearly (b) follows from (a) and the fact that

$$\eta \in NN^0 \leftrightarrow (\exists \eta_1 \in N)(\forall n_1 \in \eta_1)(\forall \xi_1 \in N)(\exists m_1 \in \xi_1)[\langle n_1, m_1 \rangle \in \eta].$$

To prove the first assertion, fix x and suppose $x \in F$. Get an N -chain $\theta \in RN$ such that $(\forall \varepsilon \in \theta)[x \in F_\varepsilon]$. Now, consider the following strategy for \exists . As his first move \exists plays $\eta_0 = \{n : \langle n \rangle \in \theta\}$ which is clearly in N . Any response $n_0 \in \eta_0$ by \forall gives a set $\eta_1 = \{n : \langle n_0, n \rangle \in \theta\} \in N$ and \exists should play η_1 as his next move. If \forall then plays $n_1 \in \eta_1$, we still have $\langle n_0, n_1 \rangle \in \theta$ and \exists responds with $\eta_2 = \{n : \langle n_0, n_1, n \rangle \in \theta\}$. If \exists follows this strategy, then clearly for any k , $\langle n_0, n_1, \dots, n_{k-1} \rangle \in \theta$ and so $x \in F_{\langle n_0, \dots, n_{k-1} \rangle}$. Hence it is a winning strategy for \exists in the game (*).

For the converse implication suppose σ is a winning strategy for \exists . Let θ be the set of sequences $\langle n_0, \dots, n_{k-1} \rangle$ of first k possible moves of player \forall when \exists follows this strategy σ . Clearly θ is an N -chain and since σ is a winning strategy for \exists ,

$$(\forall k) [\langle n_0, \dots, n_{k-1} \rangle \in \theta \rightarrow x \in F_{\langle n_0, n_1, \dots, n_{k-1} \rangle}].$$

Hence $x \in \mathbb{R} \bigcap_N (\{F_n : n \in \omega\}) = F$.

Remark : It follows from Lemma 2.4 that our definition of \mathbb{I}^* is equivalent (cf. Definition 2.16) to that introduced in [22v.4].

A set relation $\Gamma(w, x, A)$, where w varies over W and A varies over subset of W , is said to be monotone if

$$\Gamma(w, x, A) \ \& \ A \subset B \rightarrow \Gamma(w, x, B).$$

Such a monotone set relation induces for each x a monotone set operator given by

$$\Gamma_x(A) = \{w \in \omega : \Gamma(w, x, A)\}.$$

Γ_x^μ denotes the μ -th iterate viz.,

$$\Gamma_x^\mu = \Gamma_x \left(\bigcup_{\lambda < \mu} \Gamma_x^\lambda \right).$$

In particular, $\Gamma_x^0 = \Gamma_x(\varphi)$. We define

$$\Gamma_x^{<\mu} = \bigcup_{\lambda < \mu} \Gamma_x^\lambda$$

and put $\Gamma_x^\omega = \bigcup_{\mu} \Gamma_x^\mu$.

Clearly Γ_x^ω is the least fixed point of the operator Γ_x . We shall use elementary facts on inductive definability as found in [38].

The following result which shows how the operator IR gives rise to inductive operators is due to Hinman [22].

2.5 Theorem: For any p.a.o. $\underline{\mathbb{I}}$ and any $E \subseteq X$ in $\prod_1^{\underline{\mathbb{I}}^*}$, there exists a monotone set relation $\Gamma(s, x, A)$ operative on ω such that for all x

$$x \in E \iff s \in \Gamma_x^\omega.$$

Proof: Let $\{E_s : s \in \omega\}$ be a family of clopen subsets of X such that $E^\circ = \underline{\mathbb{I}}^* (\{E_s : s \in \omega\})$. Without loss of generality take $E_\circ = \varphi$. Let $N \subseteq 2^\omega$ be the canonical base for $\underline{\mathbb{I}}$. Then $\underline{\mathbb{I}}^* = \text{IR}_{NN^\circ} \underline{\mathbb{I}}$. Define a set relation operative on ω as follows:

$$s \in \Gamma_x^0(A) \iff x \notin E_s \vee (\forall \eta \in N)(\exists \eta' \in N)(\exists \xi \in N)(\forall m \in \xi) [s \ast \langle \langle n, m \rangle \rangle \in A] \dots (4)$$

By convention we take $s \ast t = 0$ if one of s, t is not in Seq .

Clearly, $\Gamma(s, x, A)$ is a monotone set relation. Put $E^s = \bigcap^* (\{E_{s \ast t} : t \in \omega\})$.

One can easily see that $E^0 = E^s, E^s \subseteq E_s$. We claim that for all s ,

$$x \notin E^s \iff s \in \Gamma_x^\omega \dots (5)$$

and the result follows by putting $s = s$. We shall prove the implication (\leftarrow) by induction. If $s \in \Gamma_x^0$, then $x \notin E_s \supseteq E^s$. Now suppose $s \in \Gamma_x^\mu, \mu > 0$. Then

$$x \notin E_s \vee (\forall \eta \in N)(\exists \eta' \in N)(\exists \xi \in N)(\forall m \in \xi) [s \ast \langle \langle n, m \rangle \rangle \in \Gamma_x^{\mu-1}].$$

If $s \notin \text{Seq}$ or $x \notin E_s$, we are done; otherwise by induction hypothesis

$$(\forall \eta \in N)(\exists \eta' \in N)(\exists \xi \in N)(\forall m \in \xi) [x \notin E_{s \ast \langle \langle n, m \rangle \rangle}]$$

which by Lemma 2.4 (and determinacy of closed games) implies

$$(\forall \eta \in N)(\exists \eta' \in N)(\exists \xi \in N)(\forall m \in \xi) \left\{ (\forall \eta_1 \in N)(\exists \eta'_1 \in N) \right. \\ \left. (\exists \xi_1 \in N)(\forall m_1 \in \xi_1) \dots (\exists k)(\forall \eta_k \in N) [x \notin E_{s \ast \langle \langle n, m \rangle \rangle \ast \langle \langle \eta_1, m_1 \rangle \rangle \ast \dots \ast \langle \langle \eta_k, m_k \rangle \rangle}] \right\}.$$

This clearly implies

$$(\forall \eta_0 \in N)(\exists \eta'_0 \in N)(\exists \xi_0 \in N)(\forall m_0 \in \xi_0)(\forall \eta_1 \in N)(\exists \eta'_1 \in N) \\ (\exists \xi_1 \in N)(\forall m_1 \in \xi_1) \dots (\exists k)(\forall \eta_k \in N) [x \notin E_{s \ast \langle \langle n_0, m_0 \rangle \rangle \ast \dots \ast \langle \langle \eta_{k-1}, m_{k-1} \rangle \rangle}]$$

and thus $x \notin \bigcap_{NN}^* (\{E_{s \ast t} : t \in \omega\}) = E^s$, by Lemma 2.4 again.

Conversely, let $s \notin \Gamma_x^\omega$. We shall show that $x \in E^s$ i.e.,

$$(\exists \eta_0 \in N)(\forall n_0 \in \eta_0)(\forall \xi_0 \in N)(\exists m_0 \in \xi_0)(\exists \eta_1 \in N)(\forall n_1 \in \eta_1) \quad (**)$$

$$\dots (\forall k) [x \in E_{s * \langle \langle n_0, m_0 \rangle, \dots, \langle n_{k-1}, m_{k-1} \rangle \rangle}].$$

Since $s \notin \Gamma_x(\Gamma_x^\omega)$, by definition of Γ , $s \in \text{Seq}$, $x \in E_s$ and moreover,

$$(\exists \eta_0 \in N)(\forall n_0 \in \eta_0)(\forall \xi_0 \in N)(\exists m_0 \in \xi_0) [s * \langle \langle n_0, m_0 \rangle \rangle \notin \Gamma_x^\omega].$$

Now, \exists can win the game (**) by adopting the following strategy. He picks $\eta_0 \in N$ such that for any choice of $n_0 \in \eta_0$ and $\xi_0 \in N$, there is an $m_0 \in \xi_0$ such that $s * \langle \langle n_0, m_0 \rangle \rangle \notin \Gamma_x^\omega = \Gamma_x(\Gamma_x^\omega)$. \exists plays such an m_0 when n_0, ξ_0 are chosen by \forall . Thus $x \in E_{s * \langle \langle n_0, m_0 \rangle \rangle}$ and

$$(\exists \eta_1 \in N)(\forall n_1 \in \eta_1)(\forall \xi_1 \in N)(\exists m_1 \in \xi_1) [s * \langle \langle n_0, m_0 \rangle, \langle n_1, m_1 \rangle \rangle \notin \Gamma_x^\omega].$$

\exists then picks $\eta_1 \in N$ such that for any choice of $n_1 \in \eta_1$ and $\xi_1 \in N$ made by \forall , there is an $m_1 \in \xi_1$ (and \exists plays such an m_1) such that $s * \langle \langle n_0, m_0 \rangle, \langle n_1, m_1 \rangle \rangle \notin \Gamma_x^\omega$. Proceeding this way, \exists has a strategy which ensures that for all k , $x \in E_{s * \langle \langle n_0, m_0 \rangle, \dots, \langle n_{k-1}, m_{k-1} \rangle \rangle}$ and so \exists wins the game (**). Consequently, $x \in E^s$.

Remark : If E is such that $x \notin E \leftrightarrow x \in \mathcal{R}\mathcal{I}_N(\{E_s : s \in \omega\})$, then the inductive operator takes a simpler form, viz.,

$$s \in \Gamma_x(A) \leftrightarrow x \notin E_s \vee (\forall \eta \in N)(\exists n \in \eta) [s * \langle n \rangle \in A]. \quad \dots (6)$$

More generally, if $\mathcal{X} = \{N_p : p \in \omega\}$ is a sequence of bases and

$$E^0 = \mathcal{R}\mathcal{X}(\{E_s : s \in \omega\}),$$

then we take the following inductive operator

$$s \in \Gamma_x(A) \leftrightarrow x \notin E_s \vee (\forall n \in \mathbb{N}_s)(\exists n \in \eta)[s * \langle n \rangle \in A], \dots (7)$$

and the conclusion of the above theorem still holds.

The inductive operator (4) (or (6) or (7)) is called the canonical inductive operator associated with $\{E_s : s \in \omega\}$ and N (or χ).

The above characterization theorem yields a decomposition of sets in $\prod_{i=1}^{\infty} \bar{\Phi}_i^*$ ($\Sigma_{i=1}^{\infty} \bar{\Phi}_i^*$) into simpler sets as is evident from the next theorem.

2.6 Definition: For any operation $\bar{\Phi}$, $\nabla(\bar{\Phi})$ is the smallest class of relations containing clopen relations and closed under $\bar{\Phi}$ and $\bar{\Phi}^{\circ}$.

$$\text{Thus } \nabla(U) = \nabla(\cap) = \Delta_1^1.$$

Let $E \in \Sigma_{i=1}^{\infty} \bar{\Phi}_i^*$ and suppose $E = \bar{\Phi}^*(\{E_s : s \in \omega\})$. Let N and Γ be as in Theorem 2.5. Then by 2.5, $E(x) \leftrightarrow s \notin \Gamma_x^{\omega}$. Set

$$E_s^{\mu} = \{x : s \notin \Gamma_x^{\mu}\}.$$

Then $E = \bigcap_{\mu < \omega_1} E_s^{\mu}$. Now define

$$T^{\mu} = \bigcup_{s \in \omega} (E_s^{\mu} - E_s^{\mu+1}).$$

It is easy to prove by induction on μ that for all $\mu < \omega_1$, $s \in \omega$, E_s^{μ} and T^{μ} are in $\nabla(\bar{\Phi})$ if it is closed under countable union. It is not hard to check that $E_s^{\mu} \downarrow$ while $(E_s^{\mu} - T^{\mu}) \uparrow$.

$$\text{2.7 Theorem: } E = \bigcup_{\mu < \omega_1} (E_s^{\mu} - T^{\mu}) = \bigcap_{\mu < \omega_1} E_s^{\mu}.$$

Proof: Let $x \in E$. Define

$$\beta(s) = \begin{cases} \text{least ordinal } \rho \text{ such that } x \notin E_s^{\rho} & \text{if } (\exists \rho)(x \notin E_s^{\rho}) \\ 0 & \text{otherwise.} \end{cases}$$

Let $\omega_1 > \rho_0 > \beta(s)$, $\forall s$. Then $(\forall s) [x \in E_s^{\rho_0} \leftrightarrow x \in E_s^{\rho_0+1}]$. Thus, $x \notin T^{\rho_0}$ and thus $x \in E_s^{\rho_0} - T^{\rho_0}$.

Conversely, suppose for some $\rho_0 < \omega_1$, $x \in E_s^{\rho_0} - T^{\rho_0}$. Then

$$(\forall s) [x \in E_s^{\rho_0} \leftrightarrow x \in E_s^{\rho_0+1}].$$

One can check by induction that

$$(\forall \rho > \rho_0)(\forall s) [x \in E_s^{\rho_0} \leftrightarrow x \in E_s^{\rho}].$$

So in particular, $x \in \bigcap_{\rho > \rho_0} E_s^{\rho} = \bigcap_{\rho < \omega_1} E_s^{\rho} = E$.

Standard arguments using the above decomposition and the countable chain condition yield the following (cf. [22]).

2.8 Theorem: If $\bar{\Phi}$ preserves measurability (preserves the Baire property), then so does $\bar{\Phi}^*$.

Remark (1): More generally, if $\mathcal{N} = \{N_p : p \in \omega\}$ is a family of bases such that each $\bar{\Phi}_{N_p}$ is measurability (Baire property) preserving, then a similar argument shows that the operation $\bar{\Phi}_{\mathcal{N}}$ is also measurability (Baire property) preserving.

Remark (2): The above proof uses the countable chain condition (ccc) of the Boolean algebra of sets with Baire property (measurability property) modulo meager (null) sets. Below we give a different proof which avoids the use of ccc and is an adaptation of the classical proof of Marozewski (see [28, § 11.VII]). Our result generalizes Marozewski's result to the R -operations.

2.8 (bis) Theorem: Let \mathcal{B} be a σ -algebra of subsets of an arbitrary set X such that for every $P \subseteq X$ there is some $\tilde{P} \supseteq P$ in \mathcal{B} with the property that if $P \subseteq B \in \mathcal{B}$ and $A \subseteq \tilde{P} - B$ then $A \in \mathcal{B}$. Then if \mathcal{B} is closed under $\bar{\cap}$ then \mathcal{B} is also closed under $\bar{\cap}^*$.

Proof: Let N be the canonical base for the p.a.o. $\bar{\cap}$. Let $\{E_n : n \in \omega\}$ be a family of sets in \mathcal{B} , and let $E = \bar{\cap}^* (\{E_n : n \in \omega\})$. Put, for each a ,

$$E^a = \bar{\cap}^* (\{E_{a * t} : t \in \omega\}),$$

so that $E^a = E$ and $E^a \subseteq E_a$. For each a , find \tilde{E}^a in \mathcal{B} such that $E^a \subseteq \tilde{E}^a$ and

$$E^a \subseteq B \in \mathcal{B} \ \& \ A \subseteq \tilde{E}^a - B \longrightarrow A \in \mathcal{B}.$$

We may, without loss of generality, assume that

$$\tilde{E}^a \subseteq E_a.$$

We shall show that

$$\tilde{E}^a - E \subseteq \bigcup_a [\tilde{E}^a - \bar{\cap}_{NN^0} (\{E^{a * \langle n \rangle} : n \in \omega\})] \quad (*)$$

Suppose not. Then there is an $x \in \tilde{E}^a - E$ such that

$$(\forall s) [x \in \tilde{E}^a \longrightarrow x \in \bar{\cap}_{NN^0} (\{E^{a * \langle n \rangle} : n \in \omega\})].$$

We claim that this implies

$$\begin{aligned} & (\exists \eta_0 \in N) (\forall \eta_0 \in \eta_0) (\forall \xi_0 \in N) (\exists m_0 \in \xi_0) \dots \dots \dots \\ & \dots \dots (\forall k) [x \in E^{\langle \langle \eta_0, m_0 \rangle, \dots, \langle \eta_{k-1}, m_{k-1} \rangle \rangle}] \end{aligned} \quad (**)$$

To see this first observe that $x \in \bigcap_{n \in \omega} (\{E^{<n>} : n \in \omega\})$.

Hence $(\exists \eta_0 \in N)(\forall n_0 \in \eta_0)(\forall \xi_0 \in N)(\exists m_0 \in \xi_0) [x \in E^{<<n_0, m_0>>}]$.

Thus in the game (**) \exists plays according to the following strategy.

He plays $\eta_0 \in N$ such that whatever be $n_0 \in \eta_0$ and $\xi_0 \in N$ played by \forall there is an $m_0 \in \xi_0$ such that $x \in E^{<<n_0, m_0>>}$. Then \exists plays such an m_0 . Since

$$x \in E^{<<n_0, m_0>>}, x \in \bigcap_{n \in \omega} (\{E^{<<n_0, m_0>, n>} : n \in \omega\})$$

and so $(\exists \eta_1 \in N)(\forall n_1 \in \eta_1)(\forall \xi_1 \in N)(\exists m_1 \in \xi_1) [x \in E^{<<n_0, m_0>, <n_1, m_1>>}]$.

Consequently \exists plays as described above. Proceeding in this manner it is clear that \exists wins the game (**). But since $\tilde{E}^s \subseteq E^s$, by Lemma 2.4 and (**) we have $x \in E$. This is a contradiction. Thus (*) holds. Now for each s ,

$$F^s \stackrel{\text{def}}{=} \tilde{E}^s - \bigcap_{n \in \omega} (\{E^{s* <n>} : n \in \omega\})$$

$$\subseteq \tilde{E}^s - \bigcap_{n \in \omega} (\{E^{s* <n>} : n \in \omega\}) = \tilde{E}^s - E^s.$$

By hypothesis, each F^s is in \mathcal{B} and since $(F^s)^c \supseteq E^s$, therefore, every subset of $F^s \in \mathcal{B}$. Hence, by (*), $\tilde{E}^s - E \in \mathcal{B}$. So,

$$E = \tilde{E}^s - (\tilde{E}^s - E) \in \mathcal{B}.$$

Remark : The family of sets with Baire property in any topological space satisfies the hypothesis of the above theorem ([20] § 11.IV 1).

Consequently, in any topological space the family of sets with Baire property is closed under every R -operation (cf. § 3). (The family of universally measurable sets also has this closure property and this follows directly from Theorem 2.8).

2.9 Definition: Call a p.a.o. $\bar{\Phi}$ a Δ_r^1 operation if Δ_r^1 is closed under $\bar{\Phi}$.

2.10 Theorem (Hinman): For any p.a.o. $\bar{\Phi}$, and any $r \geq 1$, $\bar{\Phi}$ is a Δ_r^1 operation iff $B \in \Delta_r^1$, where B is the canonical base for $\bar{\Phi}$.

Proof: If $\bar{\Phi}$ is a Δ_r^1 operation, then clearly $B \in \Delta_r^1$, since $B = \bar{\Phi}(\{D_i : i \in \omega\})$, where D_i is as in the Proposition of § 1. Suppose now that $B \in \Delta_r^1$ and let $\{E_n : n \in \omega\}$ be a family of Δ_r^1 relations. First notice that the canonical base B^0 of $\bar{\Phi}^0$ is in Δ_r^1 by (1). Then we have

$$\begin{aligned} x \in \bar{\Phi}(\{E_n : n \in \omega\}) &\leftrightarrow (\exists \eta \in B)(\forall n \in \eta)[x \in E_n] \\ &\leftrightarrow (\forall \eta \in B^0)(\exists n \in \eta)[x \in E_n]. \end{aligned}$$

This shows that $\bar{\Phi}$ is a Δ_r^1 relation.

We shall omit the proof of the next theorem which can be found in [22].

2.11 Theorem: For any p.a.o. $\bar{\Phi}$, $\Delta_1^{\bar{\Phi}^*}$ is closed under both $\bar{\Phi}$ and $\bar{\Phi}^0$.

Hence $\Sigma(\bar{\Phi}) \subseteq \Delta_1^{\bar{\Phi}^*}$. \square

2.12 Theorem: Let $\mathcal{N} = \{N_p : p \in \omega\}$ be a family of complete bases such that each $N_p \in \Delta_r^1$, with $r \geq 2$. Then $\mathcal{R}_{\mathcal{N}}$ is a Δ_r^1 operation.

In particular, if Φ is a Δ_r^1 operation, then Φ^* is also Δ_r^1 .

Proof : This follows immediately from the fact that if $\{E_n : n \in \omega\}$ is a family of Δ_r^1 sets, then the canonical inductive operator (7) is both pos \prod_r^1 and pos Σ_r^1 . The result then follows from Theorem 2.5 and [38:7E.2].

A standard diagonal argument gives us the following (cf [22; V.4.11]).

2.13 Corollary : For all $r \geq 2$ and any Δ_r^1 p.a.o. Φ ,

$$\nabla(\Phi)(\subseteq \Delta_r^1) \subseteq \Delta_r^1 \quad \square$$

Remark : An effective analogue of the above corollary can also be deduced (see [22; V.5]).

A pointclass Γ is said to be closed under $\Phi = \Phi_N$ if for any relation $E \subseteq \omega \times X$ in Γ the relation

$$(\Phi E)(x) \leftrightarrow (\exists n \in N)(\forall \eta \in \eta)E(n, x)$$

is also in Γ .

The next result is due to A. Maitra (see also [1]).

2.14 Theorem : Let Γ be a Spector pointclass closed under $\forall^{\omega\omega}$ or $\exists^{\omega\omega}$. Let $\mathcal{N} = \{N_p : p \in \omega\}$ be a sequence of complete bases which are (uniformly) in Γ . Then Γ is closed under the operation $\mathcal{R}_{\mathcal{N}}$.

Proof : We say, that a set relation $\varphi(w, x, A)$, operative on ω , is Γ on Γ if for every $Q \subseteq Z \times \omega$, the relation

$$P(x, z) \leftrightarrow \varphi(\omega, z, \{w : Q(z, w)\})$$

is also in Γ .

Now suppose $E(s,x)$ is a relation in Γ and let

$$P = \mathcal{R}_\Gamma^0 (\{E_s : s \in \omega\}).$$

We shall show that $P \in \Gamma$. First observe that if N_P^0 is the base for the dual operation, then N_P^0 is (uniformly) in Γ . Now consider the following set relation:

$$\begin{aligned} s \in \varphi_x(A) &\leftrightarrow x \in E_s \vee (\forall \eta \in N_s) (\exists n \in \eta) [s^{\wedge \langle n \rangle} \in A] \\ &\leftrightarrow x \in E_s \vee (\exists \eta \in N_s^0) (\forall n \in \eta) [s^{\wedge \langle n \rangle} \in A]. \end{aligned}$$

By the remark following 2.5, we have

$$P(x) \leftrightarrow s \in \varphi_x^{\omega}.$$

Moreover, it is easy to check that in either case φ is Γ on Γ .

Hence by [38: 7C.8], $P \in \Gamma$. \square

The proof of the above also yields the following.

2.15 Corollary (Aczel): If a Spector pointclass Γ is closed under both $\bar{\Phi}$ and $\bar{\Phi}^0$, then it is closed under $\bar{\Phi}^*$. \square

2.16 Definition: Let $\bar{\Phi}$ and Ψ be two p.a.o.'s. Say that $\bar{\Phi}$ subsumes Ψ (in symbols, $\bar{\Phi} \geq \Psi$ or $\Psi \leq \bar{\Phi}$) if there is a recursive function $f : \omega \rightarrow \omega$ such that for any family $\{F_n : n \in \omega\}$,

$$\Psi(\{F_n : n \in \omega\}) = \bar{\Phi}(\{F_{f(n)} : n \in \omega\}).$$

Say that $\bar{\Phi}$ and Ψ are equivalent ($\bar{\Phi} \sim \Psi$) if $\bar{\Phi} \geq \Psi$ and $\Psi \geq \bar{\Phi}$.

For example, \mathcal{A} subsumes both \cap and \cup .

A positive analytical operation $\bar{\Phi}$ is said to be normal if there is recursive a function g such that for any family $\{F_n : n \in \omega\}$,

$$\bar{\Phi}(\{\bar{\Phi}(\{F_{\langle p, q \rangle} : q \in \omega\}) : p \in \omega\}) = \bar{\Phi}(\{F_{g(n)} : n \in \omega\}).$$

We shall omit the proof of the next lemma which can be found in [20, 21].

2.17 Lemma : For any operations $\bar{\Phi}$ and ψ

$$(a) \quad \bar{\Phi}^{\circ\circ} = \bar{\Phi};$$

$$(b) \quad \bar{\Phi} \geq \psi \rightarrow \bar{\Phi}^{\circ} \geq \psi^{\circ};$$

$$(c) \quad \bar{\Phi}\bar{\Phi}^{\circ} \geq \bar{\Phi}, \bar{\Phi}^{\circ};$$

$$(d) \quad R\bar{\Phi} \geq \bar{\Phi};$$

$$(e) \quad \bar{\Phi} \geq \psi \rightarrow R\bar{\Phi} \geq R\psi;$$

$$(f) \quad R\bar{\Phi} \geq \psi \text{ and } R\bar{\Phi} \geq \psi^{\circ} \rightarrow R\bar{\Phi} \geq \psi\psi^{\circ};$$

$$(g) \quad \bar{\Phi}R\bar{\Phi} \rightsquigarrow R\bar{\Phi};$$

$$(h) \quad R\bar{\Phi} \rightsquigarrow RR\bar{\Phi};$$

$$(i) \quad R\bar{\Phi} \text{ is normal.}$$

3. The R-sets : We shall first construct a sequence $\{R_{\rho} : \rho < \omega_1\}$ of positive analytical operations by induction as follows. Put $R_0 = \mathcal{A}$ and having defined R_{ρ} , put

$$R_{\rho+1} = R_{\rho}^*.$$

If λ is a limit, choose a sequence $\rho_i \uparrow \lambda$ and set, for any family $\{E_n : n \in \omega\}$,

$$\bar{\Phi}(\{E_n : n \in \omega\}) = \bigcap_{i=0}^{\infty} \bar{\Phi}_{N_{\rho_i}} N_{\rho_i}^{\circ} (\{E_{\langle i, m \rangle} : m \in \omega\}),$$

where N_{ρ_i} is the canonical base for R_{ρ_i} . Then define

$$R_\lambda = \mathbb{I}^*.$$

Note that any other sequence $\rho'_i \uparrow \lambda$ gives rise to an equivalent operation by Lemma 2.17. Also, it is easy to check that $R_\rho \geq R_{\rho'}$ if $\rho \geq \rho'$.

For each $\rho < \omega_1$, let $\mathcal{R}^\rho = \Sigma_1^{R_\rho}$ and $\mathcal{BR}^\rho = \Delta_1^{R_\rho}$. Let \mathcal{BR}^ρ be the least class containing clopen relations and closed under R_ρ and complementation. Thus, for instance, $\mathcal{R}^0 = \Sigma_1^1$, $\mathcal{BR}^0 = \Delta_1^1$ and $\mathcal{BR}^0 = \mathcal{C}$ -sets of Selivanovskii. Finally, set

$$\mathcal{R} = \bigcup_{\rho < \omega_1} \mathcal{R}^\rho.$$

Members of \mathcal{R} are known classically as the R -sets. It is not difficult to see that for each ρ ,

$$\mathcal{R}^\rho \subseteq \mathcal{BR}^\rho \subseteq \mathcal{BR}^{\rho+1} \subseteq \mathcal{R}^{\rho+1}.$$

In fact, the inclusions can be shown to be strict (cf. [34]). For instance, to show that the middle inclusion is strict one observes that \mathcal{BR}^ρ consists precisely of those sets in $\mathcal{R}^{\rho+1}$ with bounded constituents. That the other two inclusions are also strict can be proved by the standard universal set argument.

The next result follows immediately from 2.12, 2.13, 2.8 and the following remarks and 2.17.

3.1 Theorem : (a) For each $\rho < \omega_1$, $\mathcal{R}^\rho \subsetneq \Delta_2^1$.

(b) R -sets are absolutely measurable and have the Baire property.

(c) For any arbitrary (T_1) topological space X , the family of sets with Baire property, \mathcal{B} , is closed under R_ρ , for each $\rho < \omega_1$.

(d) For every $\rho < \omega_1$, R_ρ is normal and $RR_\rho \hookrightarrow R_\rho$. \parallel

Remark: Theorem 3.1 (b) also follows from the fact that R -sets are provably Δ^1_2 , while 3.1 (c) follows from a recent result of Schilling [42] which states that absolutely Δ^1_2 operations preserve the Baire property in all topological spaces and one can show that each R_ρ is an absolute Δ^1_2 -operation.

For any class of sets \mathcal{F} and any operation \mathbb{I} , let $\mathbb{I}[\mathcal{F}]$ denote the class of sets which result on applications of \mathbb{I} to families of sets from \mathcal{F} ; $C\mathcal{F}$ denotes the class of complements of sets in \mathcal{F} .

Just as in the case of the Borel class we decompose, for each ρ , the class BR^ρ into a hierarchy as follows. Let $R_\rho = \mathbb{I}^*$. We set $\mathcal{R}_0^\rho = \mathcal{R}^\rho$ and define, for $\mu < \omega_1$

$$\mathcal{R}_\mu^\rho = R_\rho \left[C \left(\bigcup_{\lambda < \mu} \mathcal{R}_\lambda^\rho \right) \right],$$

$$BR_\mu^\rho = \{E \mid E, E^c \in \mathcal{R}_\mu^\rho\},$$

and BR_μ^ρ as the smallest class containing \mathcal{R}_μ^ρ and closed under \cap and complementation. It is immediate that these classes are included in BR^ρ and indeed that $BR^\rho = \bigcup_{\mu < \omega_1} \mathcal{R}_\mu^\rho$. As above we have, for each $\rho, \mu < \omega_1$,

$$\mathcal{R}_\mu^\rho \subsetneq BR_\mu^\rho \subsetneq BR_{\mu+1}^\rho \subsetneq \mathcal{R}_{\mu+1}^\rho.$$

We shall now show that each class $\mathcal{R}^\rho = \prod_{\rightarrow 1}^{\mathcal{R}^\rho}$ has the prewellordering property. The key to this is the following lemma.

If $E = \mathcal{R}_{\mathcal{N}}(\{E_n : n \in \omega\})$, where $\mathcal{N} = \{N_p : p \in \omega\}$ and Γ the canonical inductive operator associated with $\{E_n\}$ and \mathcal{N} , then we have

$$x \notin E \iff \exists \rho \in \Gamma_x^\omega.$$

Put

$$\begin{aligned} |(s, x)|_\Gamma &= \text{least } \rho \text{ such that } s \in \Gamma_x^\rho, \text{ if such } \rho \text{ exists.} \\ &= \omega_1, \text{ otherwise.} \end{aligned}$$

3.2 Lemma (Comparison of Indices) : Let $\mathcal{E} = \{E_n : n \in \omega\}$ and $\mathcal{F} = \{F_n : n \in \omega\}$ be two families of subsets of X and further assume that \mathcal{F} is regular, i.e., $F_t \subseteq F_s$ if $s \subseteq t$. Let $\mathcal{N} = \{N_p : p \in \omega\}$, $\mathcal{M} = \{M_p : p \in \omega\}$ be two sequences of bases. Define a sequence of bases $\{K_s : s \in \omega\}$ as follows. If $s = \langle \langle n_0, m_0 \rangle, \dots, \langle n_k, m_{k-1} \rangle \rangle$, then K_s is the canonical base for the positive analytical operation Φ defined by

$$\Phi(\{G_n : n \in \omega\}) = \Phi_{M_{\langle m_0, \dots, m_{k-1} \rangle}}^{\Phi_{N_{\langle n_0, \dots, n_{k-1} \rangle}}}(\{G_{\langle n, m \rangle} : n \in \omega\} : m \in \omega).$$

Otherwise, $K_s = \{\omega\}$. Let $\mathcal{K} = \{K_s : s \in \omega\}$,

$$H_s = \begin{cases} E_{\langle n_0, \dots, n_{k-1} \rangle} \cup F_{\langle m_0, \dots, m_{k-1} \rangle} & ; \text{ if } s = \langle \langle n_0, m_0 \rangle, \dots, \langle n_{k-1}, m_{k-1} \rangle \rangle \\ X, & \text{ otherwise.} \end{cases}$$

Suppose Γ is the canonical inductive operator associated with \mathcal{E} and \mathcal{N} ; Δ the inductive operator associated with \mathcal{F} and \mathcal{M} . Then,

$$\{x : |(\sigma, x)|_{\Gamma} < |(\sigma, x)|_{\Delta}\}^{\sigma} = \mathcal{R}_{\mathcal{K}}(\{H_s : s \in \omega\}).$$

Proof : First note that

$$\begin{aligned} \eta \in \mathcal{K}_{\langle\langle n_0, m_0 \rangle, \dots, \langle n_{i-1}, m_{i-1} \rangle\rangle} \\ \iff (\exists \eta' \in \mathcal{M}_{\langle\langle m_0, \dots, m_{i-1} \rangle\rangle}^{\sigma}) (\forall m \in \eta') (\exists \eta'' \in \mathcal{N}_{\langle\langle n_0, \dots, n_{i-1} \rangle\rangle}) \\ (\forall n \in \eta'') [\langle n, m \rangle \in \eta]. \end{aligned}$$

The operators Γ and Δ are as follows.

$$\begin{aligned} s \in \Gamma_x(A) &\iff x \notin E_s \vee (\forall \eta \in \mathcal{N}_s) (\exists n \in \eta) [s * \langle n \rangle \in A] \\ s \in \Delta_x(A) &\iff x \notin F_s \vee (\forall \eta \in \mathcal{M}_s) (\exists n \in \eta) [s * \langle n \rangle \in A]. \end{aligned}$$

To obtain the result look at the canonical inductive operator associated with \mathcal{K} and $\mathcal{K} = \{H_s : s \in \omega\}$;

$$s \in \Lambda_x(A) \iff x \notin H_s \vee (\forall \eta \in \mathcal{K}_s) (\exists n \in \eta) [s * \langle n \rangle \in A].$$

We claim that for all $x \in X$ and $n_0, m_0, \dots, n_{i-1}, m_{i-1}$

$$\begin{aligned} |(\langle n_0, \dots, n_{i-1} \rangle, x)|_{\Gamma} < |(\langle m_0, \dots, m_{i-1} \rangle, x)|_{\Delta} \\ \iff \langle\langle n_0, m_0 \rangle, \dots, \langle n_{i-1}, m_{i-1} \rangle\rangle \in \Lambda_x^{\omega} \end{aligned} \quad \dots(8)$$

and the result follows by taking $i = 0$.

We shall prove the implication (\rightarrow) by induction on

$|(\langle n_0, \dots, n_{i-1} \rangle, x)|_{\Gamma}$. Suppose

$$\rho = |(\langle n_0, \dots, n_{i-1} \rangle, x)|_{\Gamma} < |(\langle m_0, \dots, m_{i-1} \rangle, x)|_{\Delta}$$

and assume, to the contrary, $\langle\langle n_0, m_0 \rangle, \dots, \langle n_{i-1}, m_{i-1} \rangle\rangle \notin \Lambda_x^{\omega}$. Then

$$x \in (E_{\langle n_0, \dots, n_{i-1} \rangle} \cup F_{\langle m_0, \dots, m_{i-1} \rangle}^0) \& (\exists \eta \in K_{\langle \langle n_0, m_0 \rangle, \dots, \langle n_{i-1}, m_{i-1} \rangle \rangle}) \dots (i)$$

$$(\forall s \in \eta) [\langle \langle n_0, m_0 \rangle, \dots, \langle n_{i-1}, m_{i-1} \rangle \rangle * \langle s \rangle \notin \Delta_x^0].$$

Now, $\langle m_0, \dots, m_{i-1} \rangle \notin \Delta_x^\rho$ and so

$$x \in F_{\langle m_0, \dots, m_{i-1} \rangle} \& (\exists \eta' \in M_{\langle m_0, \dots, m_{i-1} \rangle}) (\forall m \in \eta') [\langle m_0, \dots, m_{i-1}, m \rangle \notin \Delta_s^{\langle \rho \rangle}].$$

This implies

$$x \in F_{\langle m_0, \dots, m_{i-1} \rangle} \& \dots (ii)$$

$$(\exists \eta' \in M_{\langle m_0, \dots, m_{i-1} \rangle}) (\forall m \in \eta') [\rho \leq |(\langle m_0, \dots, m_{i-1}, m \rangle, x)|_\Delta].$$

Clearly from (i) and (ii), $x \in E_{\langle n_0, \dots, n_{i-1} \rangle}$ and since $\langle n_0, \dots, n_{i-1} \rangle \in \Gamma_x^\rho$,

$$(\forall \eta'' \in N_{\langle n_0, \dots, n_{i-1} \rangle}) (\exists n \in \eta'') [\langle n_0, \dots, n_{i-1}, n \rangle \in \Gamma_x^{\langle \rho \rangle}]$$

which gives

$$(\forall \eta'' \in N_{\langle n_0, \dots, n_{i-1} \rangle}) (\exists n \in \eta'') [|(\langle n_0, \dots, n_{i-1}, n \rangle, x)|_\Gamma < \dots (iii)$$

$$\rho = |(\langle n_0, \dots, n_{i-1} \rangle, x)|_\Gamma].$$

Fix $\eta \in K_{\langle \langle n_0, m_0 \rangle, \dots, \langle n_{i-1}, m_{i-1} \rangle \rangle}$ to satisfy (i). Get $\eta' \in M_{\langle m_0, \dots, m_{i-1} \rangle}$

such that

$$(\forall m \in \eta') (\exists \eta'' \in N_{\langle n_0, \dots, n_{i-1} \rangle}) (\forall n \in \eta'') [\langle n, m \rangle \in \eta]. \dots (iv)$$

Clearly, (ii) and (1) of δ_1 yield $m^* \in \eta'$ such that

$$\rho \leq |(\langle m_0, \dots, m_{i-1}, m^* \rangle, x)|_\Delta.$$

By (iv) corresponding to m^* get $\eta'' \in N_{\langle n_0, \dots, n_{i-1} \rangle}$ such that
 $(\forall n \in \eta'') [\langle n, m^* \rangle \in \eta]$. By (iii) get $n^* \in \eta''$ such that
 $l(\langle n_0, \dots, n_{i-1}, n^* \rangle, x) \uparrow \Gamma < \rho$. Clearly, $\langle n^*, m^* \rangle \in \eta$ and

$$l(\langle n_0, \dots, n_{i-1}, n^* \rangle, x) \uparrow \Gamma < \rho \leq l(\langle m_0, \dots, m_{i-1}, m^* \rangle, x) \uparrow \Delta$$

and by the induction hypothesis this implies

$$\langle \langle n_0, m_0 \rangle, \dots, \langle n_{i-1}, m_{i-1} \rangle, \langle n^*, m^* \rangle \rangle \in \Lambda_x^\omega.$$

This clearly contradicts our choice of η . Hence

$$\langle \langle n_0, m_0 \rangle, \dots, \langle n_{i-1}, m_{i-1} \rangle \rangle \in \Lambda_x^\omega.$$

To prove the reverse implication, set

$$A^* = \left\{ \langle \langle n_0, m_0 \rangle, \dots, \langle n_{i-1}, m_{i-1} \rangle \rangle ; \right. \\ \left. l(\langle n_0, \dots, n_{i-1} \rangle, x) \uparrow \Gamma < l(\langle m_0, \dots, m_{i-1} \rangle, x) \uparrow \Delta \right\} \\ \cup \left\{ t : t \text{ is not of the form } \langle \langle n_0, m_0 \rangle, \dots, \langle n_{i-1}, m_{i-1} \rangle \rangle \right\}.$$

We shall show that $\Lambda_x(A^*) \subseteq A^*$, from which it will follow that

$\Lambda_x^\omega \subseteq A^*$. So let

$$\langle \langle n_0, m_0 \rangle, \dots, \langle n_{i-1}, m_{i-1} \rangle \rangle \in \Lambda_x(A^*). \quad \dots (v)$$

We will have to show that

$$l(\langle n_0, \dots, n_{i-1} \rangle, x) \uparrow \Gamma < l(\langle m_0, \dots, m_{i-1} \rangle, x) \uparrow \Delta.$$

Assume to the contrary that

$$l(\langle m_0, \dots, m_{i-1} \rangle, x) \uparrow \Delta \leq l(\langle n_0, \dots, n_{i-1} \rangle, x) \uparrow \Gamma.$$

From (v) we have

$$x \in E_{\langle n_0, \dots, n_{i-1} \rangle}^{\square} \cap F_{\langle m_0, \dots, m_{i-1} \rangle} \vee (\forall \eta \in K_{\langle \langle n_0, m_0 \rangle, \dots, \langle n_{i-1}, m_{i-1} \rangle \rangle}) (\exists \varepsilon \in \eta) [\langle \langle n_0, m_0 \rangle, \dots, \langle n_{i-1}, m_{i-1} \rangle, \varepsilon \rangle \in A^*] .$$

If $x \in E_{\langle n_0, \dots, n_{i-1} \rangle}^{\square} \cap F_{\langle m_0, \dots, m_{i-1} \rangle}$, then $l(\langle n_0, \dots, n_{i-1} \rangle, x) \uparrow \Gamma = 0$ and

$$l(\langle m_0, \dots, m_{i-1} \rangle, x) \uparrow \Delta > 0 .$$

and we are done. So assume

$$x \in (E_{\langle n_0, \dots, n_{i-1} \rangle} \cup F_{\langle m_0, \dots, m_{i-1} \rangle}^{\square}) \& \dots (vi)$$

$$(\forall \eta \in K_{\langle \langle n_0, m_0 \rangle, \dots, \langle n_{i-1}, m_{i-1} \rangle \rangle}) (\exists \varepsilon \in \eta) [\langle \langle n_0, m_0 \rangle, \dots, \langle n_{i-1}, m_{i-1} \rangle, \varepsilon \rangle \in A^*] .$$

If $x \in F_{\langle m_0, \dots, m_{i-1} \rangle}^{\square}$, then by regularity $x \in F_{\langle m_0, \dots, m_{i-1}, m \rangle}^{\square}$ for all

m , and hence $(\forall m) [l(\langle m_0, \dots, m_{i-1}, m \rangle, x) \uparrow \Delta = 0]$. But this is not

possible by (vi). Therefore,

$$x \in (E_{\langle n_0, \dots, n_{i-1} \rangle} \cap F_{\langle m_0, \dots, m_{i-1} \rangle}) \& \dots (vii)$$

$$(\forall \eta \in K_{\langle \langle n_0, m_0 \rangle, \dots, \langle n_{i-1}, m_{i-1} \rangle \rangle}) (\exists \varepsilon \in \eta) [\langle \langle n_0, m_0 \rangle, \dots, \langle n_{i-1}, m_{i-1} \rangle, \varepsilon \rangle \in A^*] .$$

Case 1: $l(\langle n_0, \dots, n_{i-1} \rangle, x) \uparrow \Gamma = \omega_1$. In this case $\langle n_0, \dots, n_{i-1} \rangle \notin \Gamma_x^{\square}$

and hence

$$(\exists \eta'' \in N_{\langle n_0, \dots, n_{i-1} \rangle}) (\forall n \in \eta'') [l(\langle n_0, \dots, n_{i-1}, n \rangle, x) \uparrow \Gamma = \omega_1] .$$

Fix such an $\eta'' \in N_{\langle n_0, \dots, n_{i-1} \rangle}$. Pick any $\eta' \in M_{\langle m_0, \dots, m_{i-1} \rangle}^{\square}$ and put

$$\eta^* = \{ \langle n, m \rangle : n \in \eta'' \& m \in \eta' \} .$$

Clearly $\eta^* \in K_{\langle\langle n_0, m_0 \rangle, \dots, \langle n_{i-1}, m_{i-1} \rangle\rangle}$ and, moreover,

$$(\forall \langle n, m \rangle \in \eta^*) [|(\langle m_0, \dots, m_{i-1}, m \rangle, x)|_{\Delta} \leq \omega_1 = |(\langle n_0, \dots, n_{i-1}, n \rangle, x)|_{\Gamma}].$$

This contradicts (vii).

Case 2: $|(\langle n_0, \dots, n_{i-1} \rangle, x)|_{\Gamma} = \rho < \omega_1$. Here we have, by our assumption, $\langle m_0, \dots, m_{i-1} \rangle \in \Delta_x^{\rho}$ and since $x \in F_{\langle m_0, \dots, m_{i-1} \rangle}$,

$$(\forall \eta' \in M_{\langle m_0, \dots, m_{i-1} \rangle}^{\rho}) (\exists m \in \eta') [|(\langle m_0, \dots, m_{i-1}, m \rangle, x)|_{\Delta} < \rho]$$

which implies by (2) that

$$(\exists \eta' \in M_{\langle m_0, \dots, m_{i-1} \rangle}^{\rho}) (\forall m \in \eta') [|(\langle m_0, \dots, m_{i-1}, m \rangle, x)|_{\Delta} < \rho],$$

hence

$$(\exists \eta' \in M_{\langle m_0, \dots, m_{i-1} \rangle}^{\rho}) (\forall m \in \eta') [\langle n_0, \dots, n_{i-1} \rangle \notin \Gamma [|(\langle m_0, \dots, m_{i-1}, m \rangle, x)|_{\Delta}]].$$

Consequently,

$$(\exists \eta' \in M_{\langle m_0, \dots, m_{i-1} \rangle}^{\rho}) (\forall m \in \eta') (\exists \eta'' \in N_{\langle n_0, \dots, n_{i-1} \rangle}) (\forall n \in \eta'') [|(\langle m_0, \dots, m_{i-1}, m \rangle, x)|_{\Delta} \leq |(\langle n_0, \dots, n_{i-1}, n \rangle, x)|_{\Gamma}].$$

This clearly contradicts (vii). Thus in either case we have a contradiction and so $\Lambda_x^{\omega} (A^*) \subseteq A^*$. Thus $\Lambda_x^{\omega} \subseteq A^*$. This proves the other implication of (8).

The following trick is due to Lyapunov.

3.3 Lemma (Increasing the Index by 1): Let Γ be the canonical inductive operator associated with $\mathcal{E} = \{E_p : p \in \omega\}$ and $\mathcal{N} = \{N_k : k \in \omega\}$.

Define

$$E_s^* = X,$$

$$E_s^* = \begin{cases} E_{\langle n_0, \dots, n_{k-1} \rangle} & \text{if } s = \langle 0, n_0, \dots, n_{k-1} \rangle, \\ \varnothing & \text{otherwise.} \end{cases}$$

Put

$$N_s^* = \{\{0\}\},$$

$$N_s^* = \begin{cases} N_{\langle n_0, \dots, n_{k-1} \rangle} & \text{if } s = \langle 0, n_0, \dots, n_{k-1} \rangle, \\ \{\{n\} : n \in \omega\} & \text{otherwise.} \end{cases}$$

If $\mathcal{N}^* = \{N_k^* : k \in \omega\}$, then

$$\mathcal{R}_{\mathcal{N}^*}(\{E_s^* : s \in \omega\}) = \mathcal{R}_{\mathcal{N}}(\{E_s : s \in \omega\})$$

and $|(\alpha, x)|_{\Gamma^*} = |(\alpha, x)|_{\Gamma} + 1$, where Γ^* is the inductive operator associated with

$$\mathcal{C}^* = \{E_s^* : s \in \omega\} \text{ and } \mathcal{N}^*$$

Proof : We shall prove by induction on ρ that

$$\langle 0, n_0, \dots, n_{k-1} \rangle \in \Gamma_x^{*\rho} \iff \langle n_0, \dots, n_{k-1} \rangle \in \Gamma_x^\rho.$$

Now,

$$\langle 0, n_0, \dots, n_{k-1} \rangle \in \Gamma_x^{*\rho} \iff x \notin E_{\langle 0, n_0, \dots, n_{k-1} \rangle}^* \vee$$

$$(\forall n \in N_{\langle 0, n_0, \dots, n_{k-1} \rangle}^*) (\exists n \in \eta) [\langle 0, n_0, \dots, n_{k-1}, n \rangle \in \Gamma_x^{*\rho}]$$

$$\leftrightarrow x \notin E_{\langle n_0, \dots, n_{k-1} \rangle}^V$$

$$(\forall n \in N_{\langle n_0, \dots, n_{k-1} \rangle}) (\exists n \in n) [\langle n_0, \dots, n_{k-1}, n \rangle \in \Gamma_x^{<\rho}] \text{ by induction hypothesis}$$

$$\leftrightarrow \langle n_0, \dots, n_{k-1} \rangle \in \Gamma_x^\rho.$$

Consequently,

$$s \in \Gamma_x^{*\rho+1} \leftrightarrow (\forall n \in N_s^*) (\exists n \in n) [\langle n \rangle \in \Gamma_x^{*\rho}]$$

$$\leftrightarrow \langle 0 \rangle \in \Gamma_x^{*\rho}$$

$$\leftrightarrow s \in \Gamma_x^\rho.$$

Hence $\mathcal{R}_{N^*}(\{E_s^* : s \in \omega\}) = \mathcal{R}_N(\{E_s : s \in \omega\})$ and $l(s, x) |_{\Gamma^*} = l(s, x) |_{\Gamma+1}$.

The following lemma follows from above. One has only to observe that for any positive analytical operation $\bar{\Phi}$, $R\bar{\Phi}$ is normal, $R\bar{\Phi}^0 \bar{\Phi} \sim R\bar{\Phi}\bar{\Phi}^0$ and if

$$\mathcal{N} = \{N_s : s \in \omega\}$$

is a family of bases such that $\bar{\Phi} \geq \bar{\Phi}_{N_s}$ for each s , then $R\bar{\Phi} \geq \mathcal{R}_{\mathcal{N}}$ (cf [0:§7]).

3.4 Lemma: Let $\bar{\Phi} \geq U$ be a positive analytical operation and $\mathcal{N} = \{N_s : s \in \omega\}$, $\mathcal{M} = \{M_s : s \in \omega\}$ be two families of bases such that for each s , $\bar{\Phi}^*$ subsumes $\bar{\Phi}_{N_s}$ and $\bar{\Phi}_{M_s}$. Suppose $\{E_s : s \in \omega\}$ is a family of sets in $\Sigma_1^{\bar{\Phi}^*}$ and $\{F_s : s \in \omega\}$ a regular family in $\Pi_1^{\bar{\Phi}^*}$. Let Γ be the canonical inductive operator associated with $\mathcal{E} = \{E_s : s \in \omega\}$ and \mathcal{N} , and Δ that associated with $\mathcal{F} = \{F_s : s \in \omega\}$ and \mathcal{M} . Put $\beta_1(x) = l(s, x) |_{\Gamma}$ and $\beta_2(x) = l(s, x) |_{\Delta}$.

Then

$$(a) \{x : \beta_1(x) < \beta_2(x)\} \in \prod_1^{\Phi^*},$$

$$(b) \{x : \beta_1(x) < \omega_1 \ \& \ \beta_1(x) \leq \beta_2(x)\} \in \prod_1^{\Phi^*}.$$

Proof : The first assertion follows from the comparison of indices lemma and the observations made above. The second assertion follows from the first by increasing the index of Δ by 1.

By slightly modifying the inductive operator Δ in the proof of Lemma 3.2, one can obtain the following (see also the proof of Theorem II.1.1).

3.5 Corollary : Let $\Gamma, \Delta, \mathbb{I}$ be as in 3.4. Then

$$\{(s, t, x) : |(s, x)|_{\Gamma} < |(t, x)|_{\Delta}\} \in \prod_1^{\Phi^*}.$$

3.6 Theorem : For any positive analytical operation $\mathbb{I} \geq U$, $\prod_1^{\Phi^*}$ has the prewellordering property.

Proof : Let $E \in \prod_1^{\Phi^*}$ and suppose

$$E^{\square} = \mathbb{I}^*(\{A_s : s \in \omega\}) \text{ with } \{A_s\} \text{ clopen and regular.}$$

Let $\beta(x)$ be the norm on E induced by the canonical inductive operator.

Let N be the canonical base for \mathbb{I} and put $\mathcal{X} = \{NN^{\square}\}$, $\mathcal{K} = \{NN^{\square}\}$. For each s , set

$$E_s = A_s \times \mathcal{X}, \quad F_s = \mathcal{X} \times A_s,$$

and let $\beta_1(x, y), \beta_2(x, y)$ be the norms induced by the inductive operators associated with $\{E_s : s \in \omega\}, \mathcal{X}$ and $\{F_s : s \in \omega\}, \mathcal{K}$ respectively.

Since all the hypotheses of Lemma 3.4 are satisfied, the sets

$\{(x,y) : \beta_1(x,y) < \beta_2(x,y)\}$ and $\{(x,y) : \beta_1(x,y) < \omega_1 \ \& \ \beta_1(x,y) \leq \beta_2(x,y)\}$

are in $\prod_1 \mathbb{Q}^*$. But

$$\{(x,y) : \beta_1(x,y) < \beta_2(x,y)\} = \{(x,y) : \beta(x) < \beta(y)\}$$

and $\{(x,y) : \beta_1(x,y) < \omega_1 \ \& \ \beta_1(x,y) \leq \beta_2(x,y)\}$

$$= \{(x,y) : \beta(x) < \omega_1 \ \& \ \beta(x) \leq \beta(y)\}.$$

Consequently, $\prod_1 \mathbb{Q}^*$ is normed.

3.7 Corollary : For each $\rho < \omega_1$, \mathcal{R}^ρ has the prewellordering property.

3.8 Theorem : Let \mathbb{Q} be a positive analytical operation which subsumes both (countable) \cup and \cap . Let \mathcal{F} be the σ -field generated by Σ_1^* .

Let $E \in \Sigma_1^*$ such that

$$x \in E \iff (\exists n_0 \in \mathbb{N}^0)(\forall n_0 \in \eta_0)(\exists n_1 \in \mathbb{N}^0)(\forall n_1 \in \eta_1) \dots$$

$$\dots (\forall k) [x \in E_{\langle n_0, \dots, n_{k-1} \rangle}], \quad \dots (i)$$

N being the canonical base for \mathbb{Q} . Then, there is a \mathcal{F} -measurable function $x \rightarrow \sigma_x$ such that σ_x is a winning strategy for the player \exists , whenever $x \in E$.

Proof : Let Γ be the canonical inductive operator associated with \mathbb{N}^0 and $\{E_s : s \in \omega\}$ and put $\mathbb{N}^0 = M$. Define

$$\beta(s,x) = \begin{cases} \text{least } \rho \text{ such that } s \in \Gamma_x^\rho, & \text{if } s \in \Gamma_x^\omega \\ \omega_1, & \text{otherwise.} \end{cases}$$

Now suppose $x \in E$. So \exists wins the game (i).

If $\eta_0, \eta_1, \dots, \eta_{k-1}$ and n_0, n_1, \dots, n_{k-1} are the first k relevant moves of \exists and \forall , notice that \exists goes on to win the game (i), i.e., he is in a winning position iff $\langle n_0, \dots, n_{k-1} \rangle \notin \Gamma_x^{\omega}$ i.e., iff $\beta(\langle n_0, \dots, n_{k-1} \rangle, x) = \omega_1$. In such a case, \exists has to play an $\eta \in M$ such that $(\forall n \in \eta) [\beta(\langle n_0, \dots, n_{k-1}, n \rangle, x) = \omega_1]$. We therefore define for each x , the strategy σ_x for \exists as follows.

$$p \in \sigma_x(s) \iff \beta(s, x) \leq \beta(s * \langle p \rangle, x).$$

Clearly by 3.5, the map $x \rightarrow \sigma_x$ is \mathcal{F} -measurable. We shall now show that if $x \in E$, then σ_x is a winning strategy for \exists in the game (i). Suppose $\eta_0, n_0, \eta_1, n_1, \dots, \eta_{k-1}, n_{k-1}$ are the first k moves of \exists and \forall and assume that \exists is in a winning position. Consequently, we have $\beta(\langle n_0, n_1, \dots, n_{k-1} \rangle, x) = \omega_1$ and hence $\langle n_0, \dots, n_{k-1} \rangle \notin \Gamma_x^{\omega}$. Therefore,

$$(\exists \eta \in M)(\forall n \in \eta) [\beta(\langle n_0, n_1, \dots, n_{k-1}, n \rangle, x) = \omega_1] \quad \dots(ii)$$

By definition,

$$\begin{aligned} \sigma_x(\langle n_0, \dots, n_{k-1} \rangle) &= \{p : \beta(\langle n_0, \dots, n_{k-1} \rangle, x) \leq \beta(\langle n_0, \dots, n_{k-1}, p \rangle, x)\} \\ &= \{p : \beta(\langle n_0, \dots, n_{k-1}, p \rangle, x) = \omega_1\} = \eta, \text{ say.} \end{aligned}$$

Hence, by (ii) and the completeness of M ,

$$\sigma_x(\langle n_0, \dots, n_{k-1} \rangle) = \eta \in M,$$

and $(\forall p \in \eta)(\beta(\langle n_0, \dots, n_{k-1}, p \rangle, x) = \omega_1)$, so that \exists is still in a winning position, and, moreover, $x \in E_{\langle n_0, \dots, n_{k-1}, p \rangle}$ for all $p \in \eta$.

Remark : Notice that

$$\begin{aligned}
 x \notin E &\leftrightarrow (\forall n_0 \in M)(\exists n_0 \in \eta_0)(\forall n_1 \in M)(\exists n_1 \in \eta_1) \dots \\
 &\dots (\exists k) [x \notin E_{\langle n_0, \dots, n_{k-1} \rangle}] \\
 &\leftrightarrow (\exists n_0 \in M^0)(\forall n_0 \in \eta_0) \dots (\exists k) [x \notin E_{\langle n_0, \dots, n_{k-1} \rangle}] \quad \dots (iii)
 \end{aligned}$$

In this case also, a definable winning strategy for \exists can similarly be defined. Unlike the game (i), here \exists has to play such that at each stage the ordinal β is decreased. The following gives a \mathcal{F} -measurable winning strategy for \exists in the game (iii) whenever $x \notin E$.

$$p \in \tau_x(s) \leftrightarrow x \notin E_s \vee (\beta(s * \langle p \rangle, x) < \beta(s, x)).$$

4. The Ramsey property and the partition selection property of R -sets.

Let $|\omega|^\omega$ denote the set of infinite sets of natural numbers. A set $P \subseteq |\omega|^\omega$ is said to have the Ramsey property if it admits a homogeneous set H i.e. an infinite set $H \subseteq \omega$ such that every infinite subset of H belongs to P or every infinite subset of H belongs to the complement of P . Silver [45] has shown, in $ZF + DC$, that every Σ_1^1 set has the Ramsey property, while if measurable cardinals exist, Σ_2^1 sets have the Ramsey property. The question naturally arises whether, in $ZF + DC$, R -sets have the Ramsey property. A clue to the solution of this problem comes from Ellentuck's proof of Silver's theorem [17]. But first the relevant definitions which are as follow. Let s be a finite set of natural numbers and A an infinite one. The pair $\langle s, A \rangle$ is said to be a condition if every member of s is less than any member of A . The

Ellentuck neighbourhood (s, A) consists of all $X \in |\omega|^\omega$ such that $s \subseteq X \subseteq s \cup A$. Condition $\langle s, A \rangle$ extends $\langle t, B \rangle$ iff $t \subseteq s$, $s - t \subseteq B$ and $A \subseteq B$. A set $P \subseteq |\omega|^\omega$ is said to be completely Ramsey if for every condition $\langle s, A \rangle$ there is an extension $\langle s, A' \rangle$ such that $(s, A') \subseteq P$ or $(s, A') \subseteq |\omega|^\omega - P$. The set P , then, possesses the Ramsey property iff there is an $A \in |\omega|^\omega$ such that $(\emptyset, A) \subseteq P$ or $(\emptyset, A) \subseteq |\omega|^\omega - P$. In [17], Ellentuck has shown that a set $P \subseteq |\omega|^\omega$ is completely Ramsey iff it has the Baire property with respect to the Ellentuck topology. Thus to show that R-sets are completely Ramsey, and hence have the Ramsey property, we need only to show that R-sets have the Baire property. Since the Ellentuck topology is strictly finer than the usual one on $|\omega|^\omega$, \prod_1^0 sets have the Baire property. Moreover, by Theorem 3.1 sets with Baire property are closed under the operations R_ρ . Thus we have

4.1 Theorem : R-sets are completely Ramsey and hence have the Ramsey property. \parallel

Next we show that an equivalence relation generated by countably many R-sets admits a \mathcal{R} -measurable selector. This is a result of Burgess [10]. No proof appears in [10], and we give a complete proof primarily because our proof enables us to obtain some very important properties of the R-sets like the scale property and the uniformization property for the class $\mathcal{C}\mathcal{R}^P$ (see Theorems 5.4, 5.6). As in [10] we obtain our result via presentability, a notion introduced by Burgess.

4.2 Definition : Let X be an uncountable Polish space and \mathcal{E} an equivalence relation on X . Let $[x]$ denote the \mathcal{E} -equivalence class containing $x \in X$. Clearly \mathcal{E} gives rise to a partition $Q = \{ [x] : x \in X \}$ into disjoint equivalence classes. A selector for \mathcal{E} is a function $S : X \rightarrow X$ such that for each x , $S(x) \in x$ and $S(x) = S(x')$ whenever $x \in x'$. A cross-section for \mathcal{E} is a set $T \subseteq X$ consisting of exactly one element from each \mathcal{E} -equivalence class. Plainly, if S is a selector for \mathcal{E} , then $T = \text{Range}(S)$ is a cross-section. A countable family $\{A_n : n \in \omega\}$ of subsets of X generates \mathcal{E} if

$$x \mathcal{E} x' \iff (\forall n) (x \in A_n \iff x' \in A_n).$$

4.3 Definition : Let \mathcal{B} be a σ -field of subsets of X . The equivalence relation \mathcal{E} is said to be \mathcal{B} -generated if it is generated by a countable subfamily of \mathcal{B} . The σ -field \mathcal{B} is said to have the partition-selection property if every \mathcal{B} -generated equivalence relation admits a \mathcal{B} -measurable selector.

In order to show that various families have the partition-selection property Burgess [10] introduced the following technical notion.

4.4 Definition : Let \mathcal{B} and X be as above. For a set $A \subseteq X$, a \mathcal{B} -presentation of A is a quadruple (Y, \mathcal{B}, P, G) such that

- (i) Y is a Polish space,
- (ii) $B \subseteq X \times Y$ is clopen,
- (iii) $P \subseteq X \times Y$ is Π_1^1 ,
- (iv) $G : X \rightarrow Y$ is \mathcal{B} -measurable with graph $\subseteq P$,
- (v) $B \cap P = (A \times Y) \cap P$.

The σ -field \mathcal{B} is presentable if every $A \in \mathcal{B}$ admits a \mathcal{B} -presentation.

The next result follows immediately from a refinement of a theorem of Kaniewski [11].

4.5 Lemma (Burgess) : Every presentable family has the partition-selection property.

In [10], Burgess proved that the \mathcal{C} -sets, the Borel-programmable sets, the \mathcal{R} -sets, the absolute Δ_2^1 sets and the Lebesgue measurable sets have the partition selection property. We shall obtain the partition-selection property of \mathcal{R} -sets by showing that the property of being presentable is 'preserved' by the operator \mathcal{R} . (This is one of the several instances where a certain property is left invariant by the \mathcal{R} operator. Another such 'invariance property' will be exhibited in the next two sections). The presentability, and hence the partition-selection property, of $\mathcal{B}\mathcal{R}^p$ is then obtained by a simple inductive argument.

4.6 Theorem : Let \mathcal{Q} be a p.a.o. which subsumes operation (A) . If $\nabla(\mathcal{Q})$ is presentable, then $\nabla(\mathcal{Q}^*)$ is presentable.

Proof : We work in ω^ω .

Step 1. We first show that each $E \in \prod_1 \mathcal{Q}^*$ is presentable. Suppose $E^c = \mathcal{Q}^*(\{E_n : n \in \omega\})$, with each E_n clopen, and let N be the canonical base for \mathcal{Q} . The inductive operator associated with N and $\{E_n : n \in \omega\}$ is the following :

$$s \in \prod_x (A) \leftrightarrow x \notin E_s \vee (\forall n \in N)(\exists n \in \eta)(\exists \xi \in N)(\forall m \in \xi)[s * \langle \langle n, m \rangle \rangle \in A].$$

Then it is not hard to see that

$$\begin{aligned} s \notin \Gamma_x^\omega &\leftrightarrow (\exists \eta_0 \in N)(\forall \eta_0 \in \eta)(\forall \xi_0 \in N)(\exists m_0 \in \xi_0) \dots \dots \dots \\ &\dots \dots (\forall k) [x \in E_{s^*} \langle \langle \eta_0, m_0 \rangle, \dots, \langle \eta_{k-1}, m_{k-1} \rangle \rangle] \dots (i) \end{aligned}$$

In particular,

$$s \notin \Gamma_x^\omega \leftrightarrow x \notin E.$$

Let $Y = \omega^\omega \times 2^\omega \times \{0,1\}$. Define $C \subseteq X \times Y$ as follows :

$$C(x, \alpha, \beta, k) \stackrel{\text{def}}{\longleftrightarrow} \omega(\alpha) \ \&$$

$$(\forall i) [i \in \text{Field}(\leq_\alpha) \rightarrow (\beta)_i = \Gamma_x \left(\bigcup_{j <_\alpha i} (\beta)_j \right)]$$

$$\& \Gamma_x \left(\bigcup_{i \in \text{Field}(\leq_\alpha)} (\beta)_i \right) = \bigcup_{i \in \text{Field}(\leq_\alpha)} (\beta)_i \ \& \ k = d \left(\bigcup_{i \in \text{Field}(\leq_\alpha)} (\beta)_i \right);$$

where $d : 2^\omega \rightarrow \{0,1\}$ is the function defined by

$$d(\eta) = \begin{cases} 0 & \text{if } s \in \eta \\ 1 & \text{otherwise.} \end{cases}$$

Put $B = X \times \omega^\omega \times 2^\omega \times \{0\}$. Plainly B is clopen and C is in $\nabla(\mathbb{Q})$.

Moreover,

$$C \cap B = C \cap (E \times Y).$$

Let

$$\beta(n, x) = \begin{cases} \text{least ordinal } \lambda \text{ such that } n \in \Gamma_x^\lambda, & \text{if } n \in \Gamma_x^\omega \\ \omega_1 & \text{otherwise.} \end{cases}$$

Define $f : X \rightarrow \omega_1$ by

$$f(x)(\langle n, m \rangle) = 0 \iff \beta(n, x) < \beta(m, x) < \omega_1$$

$$\& (\forall i) [\beta(n, x) = \beta(i, x) \rightarrow n \leq i]$$

$$\& (\forall j) [\beta(m, x) = \beta(j, x) \rightarrow m \leq j].$$

Otherwise, put $f(x)(s) = 1$.

By Corollary 3.5, f is $\nabla(\Phi^*)$ -measurable. Now, define $g : X \rightarrow 2^\omega$ as follows :

$$(g(x))_n = \begin{cases} \Gamma_x^{\beta(n, x)} & \text{if } n \in \text{Field}(f(x)) \\ \varphi & \text{otherwise.} \end{cases}$$

For s not of the form $\langle n, m \rangle$, put $g(x)(s) = 1$.

Thus

$$g(x)(\langle n, m \rangle) = 0 \iff \beta(m, x) \leq \beta(n, x) \& n \in \text{Field}(f(x)).$$

Consequently, g is $\nabla(\Phi^*)$ -measurable by 3.5 again. Let

$$G(x) = (f(x), g(x), \chi_E(x)).$$

Then $G : X \rightarrow Y$ is $\nabla(\Phi^*)$ -measurable and $\text{graph}(G) \subseteq C$. By hypothesis, C is presentable. Get $(Z, \tilde{B}, \tilde{P}, \tilde{G})$ such that

- (i) $\tilde{P} \subseteq X \times Y \times Z$ is \prod_1^1
- (ii) $\tilde{B} \subseteq X \times Y \times Z$ is clopen
- (iii) $\tilde{G} : X \times Y \rightarrow Z$ is $\nabla(\Phi)$ -measurable with graph contained in P
- (iv) $\tilde{P} \cap \tilde{B} = \tilde{P} \cap (C \times Z)$.

Put

$$P' = \tilde{P} \cap \tilde{B}$$

$$B' = B \times Z$$

$$G'(x) = (G(x), \tilde{G}(x, G(x))).$$

It is not hard to check that $(Y \times Z, B', P', G')$ is a $\underline{\Sigma}(\Phi^*)$ -presentation of E .

Step 2. Let $\mathcal{F} \subseteq \underline{\Sigma}(\Phi^*)$ be the family of presentable sets in $\underline{\Sigma}(\Phi^*)$. As shown above, $\prod_1^{\Phi^*} \subseteq \mathcal{F}$ and trivially \mathcal{F} is closed under complementation. We shall show that \mathcal{F} is closed under Φ^* . Let $\{E_n : n \in \omega\}$ be a family of sets in \mathcal{F} and let $E = \Phi^*(\{E_n : n \in \omega\})$. Since each E_n is presentable, there is $(Y', \{B_n\}, P', G')$ which is a $\underline{\Sigma}(\Phi^*)$ -presentation of $\{E_n\}$ (see [11]). In other words,

- (i) Y' is Polish,
- (ii) each $B_n \subseteq X \times Y$ is clopen,
- (iii) $P' \subseteq X \times Y$ is \prod_1^1 ,
- (iv) $G : X \rightarrow Y$ is $\underline{\Sigma}(\Phi^*)$ -measurable with graph in P' ,
- (v) for each n , $P' \cap B_n = P' \cap (E_n \times Y)$.

Put $B = \Phi^*(\{B_n : n \in \omega\})$. Plainly, $B \in \underline{\Sigma}_1^{\Phi^*}$ and is therefore presentable. Let (Z, D, P'', H) be a $\underline{\Sigma}(\Phi^*)$ -presentation of B . Set

$$P = P'' \cap (P' \times Z)$$

$$\tilde{G}(x) = (G(x), H(x, G(x))).$$

Then $(Y \times Z, D, P, \tilde{G})$ is a $\underline{\Sigma}(\Phi^*)$ -presentation of E . \square

4.7 Lemma : Let $\{\Phi_i : i \in \omega\}$ be a sequence of p.a.o.'s with canonical bases N_i and suppose for each i , $\underline{\Sigma}(\Phi_i)$ is presentable. Define a p.a.o. ψ by

$$x \in \psi(\{E_n : n \in \omega\}) \leftrightarrow (\forall i) (\exists \eta \in N_i) (\forall n \in \eta) [x \in E_{\langle i, n \rangle}],$$

where $\{E_n : n \in \omega\}$ is any family of subsets. Then $\underline{\Sigma}(\psi)$ is also presentable.

Proof : We shall first show that every \sum_1^ψ set admits a $\nabla(\psi)$ -presentation.

So let $E \in \sum_1^\psi$. Find a family of clopen sets $\{E_n\}$ such that

$E = \psi(\{E_n : n \in \omega\})$. Set for every i

$$E^i = \bar{Q}_i(\{E_{\langle i, n \rangle} : n \in \omega\}).$$

Plainly each $E^i \in \nabla(\bar{Q}_i)$ and $E = \bigcap_i E^i$. By hypothesis, each E^i admits a $\nabla(\bar{Q}_i)$ -presentation and hence a $\nabla(\psi)$ -presentation. It is now easy to see that $E = \bigcap_i E^i$ admits a $\nabla(\psi)$ -presentation. Thus the class of sets which admit a $\nabla(\psi)$ -presentation contains \sum_1^ψ and is (trivially) closed under complementation. That it is also closed under ψ can be shown along the lines of the proof of Theorem 4.6. This shows that $\nabla(\psi)$ is presentable. \square

We now have the following result of Burgess announced in [10].

4.8 Theorem : For each $\rho < \omega_1$, \underline{BR}^ρ is presentable.

Consequently, each \underline{BR}^ρ has the partition selection property.

Proof : By [10], the class of \mathcal{C} -sets = \underline{BR}^0 is presentable. Moreover,

if \underline{BR}^ρ is presentable then by Theorem 4.6 $\underline{BR}^{\rho+1}$ is also presentable.

Finally if λ is limit and for each $\mu < \lambda$, \underline{BR}^μ is presentable, then

the presentability of \underline{BR}^λ can be established by Lemma 4.7 and Theorem 4.6. \square

Remark : It is significant to note that presentability is not preserved

by the game quantifier \mathfrak{G} (at least under some determinacy hypothesis).

To see this assume $\text{Det}(\Delta_2^1)$ and assume to the contrary that presentability is preserved by Θ . Since Δ_2^1 is presentable ([10]) and

$\Theta \Delta_2^1 = \Delta_3^1$ ([38; 6E.13]) it follows that Δ_3^1 is presentable and

hence has the partition-selection property. Now fix a set $P \in \Sigma_3^1 - \Delta_3^1$.

Let f be a continuous function and $A \subseteq \omega^\omega$ a Π_2^1 set such that

$f(A) = P$. Define an equivalence relation E on ω^ω as follows:

$$E(x,y) \leftrightarrow x \notin A \ \& \ y \notin A \vee x \in A \ \& \ y \in A \ \& \ f(x) = f(y).$$

Plainly E is a Δ_3^1 -generated equivalence relation. By the partition-

selection property, there is a Δ_3^1 -measurable selector S and hence a

Δ_3^1 -cross-section T for E . It is easy to check that $f \upharpoonright T \cap A$ is

one-one and $f(T \cap A) = P$. By [38; 6E.14], this implies that P is

Δ_3^1 which is a contradiction.

5. The scale property and the uniformization property of R-sets : The notion of scale is of great importance in descriptive set theory and it is extremely useful to know which pointclasses have the scale property. For any p.a.o. Φ , $\prod_1^{\Phi^*}$ behaves very much like \prod_1^1 and it is not surprising that for most operations Φ , $\prod_1^{\Phi^*}$ has the prewellordering property. It therefore seems natural to expect that the proof of the scale property should follow a similar pattern and that the key to this is the proof of the Scale Transfer Theorem of Moschovakis [38]. However, this method does not seem to work. On the other hand, by the game characterization of R-sets [13, §1], R-sets at the λ th level are obtained by applying the game quantifier \mathcal{G} to $\lambda - \Sigma_2^0$ sets (see [48] for definition); and hence Steel's method of transferring scale applies (cf. [48]). Although this method gives scales it does not seem to give the definable scales of right type. The failure of these methods led us to express R-sets in terms of games played on ω of length much greater than \aleph_0 and then using Martin's method (see [36]). Unfortunately, this method also failed to produce scales in the right definable class. By Martin's method we have, for instance, been able to show that cR^n -sets, $n \in \omega$, admit $\mathcal{B}R^{n+1}$ scales although we believe that this can be improved. This hopeless situation was retrieved by Srivatsa by an ingenious method. By modifying our proof of the presentability of cR^p , he has been able to 'transfer' the \prod_1^1 scales to cR^p sets and thus obtain the scale property. We present below his proof, slightly simplified by us.

The key result for the proof is the following extremely useful and important result of Burgess which refines a theorem of Kaniewski (see [11]). The notion of presentability also plays a crucial role.

5.1 Lemma (Burgess) : Let Z be a Polish space, $P \subseteq Z \prod_1^1$, $Q \subseteq Z \times Z \Sigma_1^1$, $D = Q \cap P \times P$. Suppose that D is an equivalence relation on P and that every D -equivalence class is relatively closed in P . Then D admits an analytically measurable selector whose associated cross-section is \prod_1^1 . \square

We first show that \underline{BR}^ρ is scaled.

5.2 Theorem : For each $\rho < \omega_1$, \underline{BR}^ρ is scaled.

Proof : Let $E \in \underline{BR}^\rho$. By Theorem 4.8 E is presentable. Get a \underline{BR}^ρ -presentation (Y, \mathcal{B}, P, G) as in 4.4. Define a Σ_1^1 (in fact, Borel) equivalence relation on $X \times Y$ such that its restriction on P has vertical lines as its equivalence classes. By Lemma 5.1 get an analytically measurable selector f such that the associated cross-section $P^* = \text{range}(f)$ is \prod_1^1 . Now fix a very good \prod_1^1 -scale ([38; 4E]) $\{\varphi_n\}$ on $P^* \cap \mathcal{B}$ and on E define a sequence of norms $\{\psi_n\}$ as follows.

$$\psi_n(x) = \varphi_n(f(x, G(x))).$$

$$\text{Plainly, } x \leq_{\psi_n}^* y \iff f(x, G(x)) \leq_{\varphi_n}^* f(y, G(y))$$

$$\text{and } x <_{\psi_n}^* y \iff f(x, G(x)) <_{\varphi_n}^* f(y, G(y)).$$

To check that $\{\psi_n\}$ is a scale on E is quite routine (see the proof of the next theorem). \square

To show $c\mathcal{R}^\rho$ has the scale property, one has to take more care. We work with partial functions and appropriately modify the proof of Theorem 4.6 to get the scale property.

5.3 Theorem : Let \mathbb{I} be a p.a.o. which subsumes operation A . If $\nabla(\mathbb{I})$ is presentable, then $\prod_1 \mathbb{I}^*$ is scaled.

Proof : Let N be the canonical base for \mathbb{I} and suppose $E \in \prod_1 \mathbb{I}^*$. Let $E^0 = \mathbb{I}^* (\{E_n : n \in \omega\})$, where $\{E_n : n \in \omega\}$ is a clopen family. The canonical inductive operator determining E is the following.

$$e \in \Gamma_x(A) \leftrightarrow x \notin E_s \vee (\forall \eta \in N)(\exists n \in \eta)(\exists \xi \in N)(\forall m \in \xi) [e * \langle\langle n, m \rangle\rangle \in A].$$

Then

$$e \notin \Gamma_x^\omega \leftrightarrow (\exists n_0 \in N)(\forall n_0 \in \eta_0)(\forall \xi_0 \in N)(\exists m_0 \in \xi_0) \dots \dots \dots (\forall k) [x \in E_{s * \langle\langle n_0, m_0 \rangle\rangle, \dots, \langle n_1, \dots, m_{k-1} \rangle\rangle}].$$

In particular, $e \in \Gamma_x^\omega \leftrightarrow x \in E$.

Define $C \subseteq \mathcal{X} \times \omega^\omega \times 2^\omega$ as follows.

$$C(x, \alpha, \gamma) \leftrightarrow \text{WO}(\alpha) \ \&$$

$$(\forall i) [i \in \text{Field}(\leq_\alpha) \rightarrow (\gamma)_i = \Gamma_x(\bigcup_{j <_\alpha i} (\gamma)_j)]$$

$$\& e \in \bigcup_{i \in \text{Field}(\leq_\alpha)} (\gamma)_i$$

Plainly, C is in $\nabla(\mathbb{I})$ and $\text{proj}_\mathcal{X}(C) = E$.

Let $\beta(n, x)$ denote the least ordinal λ such that $n \in \Gamma_x^\lambda$, otherwise $\beta(n, x) = \omega_1$. Define a partial function on E into WO as follows.

For every $x \in E$,

$$f(x)(s) = 0 \leftrightarrow \beta((s)_0, x) \leq \beta(e, x) \ \& \ \beta((s)_1, x) \leq \beta(e, x)$$

$$\ \& \ \beta((s)_0, x) < \beta((s)_1, x)$$

$$\ \& \ (\forall i) [\beta((s)_0, x) = \beta(i, x) \rightarrow (s)_0 \leq i]$$

$$\ \& \ (\forall j) [\beta((s)_1, x) = \beta(j, x) \rightarrow (s)_1 \leq j].$$

Otherwise, put $f(x)(s) = 1$.

Observe that, since $\beta(e, x) < \omega_1$, by 3.5 it follows easily that f is a $\prod_1^{\aleph^*}$ -recursive partial function. Now define a partial function g on E into 2^ω as follows

$$g(x)(s) = 0 \leftrightarrow \beta((s)_0, x) \leq \beta(e, x)$$

$$\ \& \ \beta((s)_1, x) \leq \beta(e, x) \ \& \ \beta((s)_1, x) \leq \beta((s)_0, x).$$

Otherwise, $g(x)(s) = 1$.

Thus when $x \in E$ and $n \in \text{Field}(\leq_{f(x)})$,

$$\begin{aligned} (g(x))_n &= \{m : \beta(m, x) \leq \beta(n, x)\} \\ &= \Gamma_x \beta(n, x). \end{aligned}$$

Furthermore, by 3.5 again, g is $\prod_1^{\aleph^*}$ -recursive on E .

Put, for $x \in E$,

$$G(x) = (f(x), g(x)).$$

Plainly, G is a partial $\prod_1^{\aleph^*}$ -recursive function whose graph is

contained in C . By hypothesis, C is presentable. Get (Z, B, P, \tilde{G}) such that

- (i) Z is Polish
- (ii) $P \subseteq X \times \omega^\omega \times 2^\omega \times Z$ is \prod_1^1
- (iii) $B \subseteq X \times \omega^\omega \times 2^\omega \times Z$ is clopen
- (iv) $\tilde{G} : X \times \omega^\omega \times 2^\omega \rightarrow Z$ is $\nabla(\mathbb{Q})$ -measurable with graph contained in P .
- (v) $P \cap B = P \cap (C \times Z)$.

As in the proof of Theorem 5.2 get a selector F measurable with respect to the analytical σ -field which from each line $\{x\} \times (P \cap B)_x$ picks a unique point and such that the corresponding cross-section P^* is \prod_1^1 . Observe that $\text{proj}_X(P^*) = E$.

Let $\{\varphi_n\}$ be a very good \prod_1^1 -scale on P^* . Define for any $x \in E$ and for every n ,

$$\psi_n(x) = \varphi_n(F(x, G(x), \tilde{G}(x, G(x)))).$$

We claim that

- (a) $x \leq_{\psi_n}^* y \iff F(x, G(x), \tilde{G}(x, G(x))) \downarrow \&$
 $(\forall \alpha)(\forall \gamma)(\forall z) [F(x, G(x), \tilde{G}(x, G(x))) \leq_{\varphi_n}^* (y, \alpha, \gamma, z)]$
- (b) $x <_{\psi_n}^* y \iff F(x, G(x), \tilde{G}(x, G(x))) \downarrow \&$
 $(\forall \alpha)(\forall \gamma)(\forall z) [F(x, G(x), \tilde{G}(x, G(x))) <_{\varphi_n}^* (y, \alpha, \gamma, z)]$

To see (a), first assume $x \leq_{\psi_n}^* y$. Then $x \in E$ and hence

$F(x, G(x), G(x, G(x))) \downarrow$. If $y \notin E$ then clearly for any α, γ, z , $(y, \alpha, \gamma, z) \notin P^*$ and so $F(x, G(x), G(x, G(x))) \leq_{\varphi_n}^* (y, \alpha, \gamma, z)$. Now assume $y \in E$ and $\psi_n(x) \leq \psi_n(y)$. Fix any α, γ, z . If

$$(y, \alpha, \gamma, z) = F(y, G(y), \tilde{G}(y, G(y))),$$

then

$$\varphi_n(F(x, G(x), \tilde{G}(x, G(x)))) \leq \varphi_n(y, \alpha, \gamma, z).$$

If $(y, \alpha, \gamma, z) \neq F(y, G(y), \tilde{G}(y, G(y)))$, then observe that $(y, \alpha, \gamma, z) \notin P^*$ and so

$$F(x, G(x), \tilde{G}(x, G(x))) \leq_{\varphi_n}^* (y, \alpha, \gamma, z)$$

Conversely, suppose $F(x, G(x), G(x, G(x))) \downarrow$ and

$$(\forall \alpha) (\forall \gamma) (\forall z) \left[F(x, G(x), G(x, G(x))) \leq_{\varphi_n}^* (y, \alpha, \gamma, z) \right].$$

Plainly $x \in E$ and assume $y \in E$. Choose α, γ, z such that

$$(y, \alpha, \gamma, z) = F(y, G(y), \tilde{G}(y, G(y))).$$

Then

$$\varphi_n(F(x, G(x), \tilde{G}(x, G(x)))) \leq \varphi_n(F(y, G(y), \tilde{G}(y, G(y))))$$

and so $\psi_n(x) \leq \psi_n(y)$. Thus we have (a); (b) is proved similarly. Now notice that

$$S(n, x, \alpha_1, \gamma_1, z_1, y) \stackrel{\text{def}}{\longleftrightarrow} (\forall \alpha) (\forall \gamma) (\forall z) \left[(x, \alpha_1, \gamma_1, z_1) \leq_{\varphi_n}^* (y, \alpha, \gamma, z) \right]$$

is \prod_1^1 since $\{\varphi_n\}$ is a \prod_1^1 -scale. Since $\prod_1^{\mathbb{Q}^*}$ has the substitution property (this can be easily established),

$$x \leq_{\psi_n}^* y \iff F(x, G(x), \vec{G}(x, G(x))) \downarrow \& S(n, F(x, G(x), \vec{G}(x, G(x))), y)$$

is $\prod_1 \vec{\psi}_i^*$. Similarly, it can be shown that the relation $x <_{\psi_n}^* y$ is $\prod_1 \vec{\psi}_i^*$. It now remains to show that $\{\psi_n\}$ is a scale on E . So let $\{x_n\}$ be a sequence in E and suppose $x_n \rightarrow x$ and for each i ,

$$\psi_i(x_n) = \lambda_i \text{ ultimately.}$$

Thus, for each i ,

$$\varphi_i(F(x_n, G(x_n), G(x_n, \vec{G}(x_n)))) = \lambda_i$$

for all large n . Since $\{\varphi_i\}_{i \in \omega}$ is a very good scale, there is $\vec{z} \in P^*$ such that

$$F(x_n, G(x_n), G(x_n, \vec{G}(x_n))) \rightarrow \vec{z}$$

and $\varphi_i(\vec{z}) \leq \lambda_i$ (c)

Since $F(x, \vec{w})$ is of the form (x, \vec{w}) , it is clear that \vec{z} is of the form (x, \vec{w}) for some \vec{w} . Thus $x \in E$ and moreover,

$$(x, \vec{w}) = F(x, G(x), \vec{G}(x, G(x))).$$

Hence $\psi_i(x) = \varphi_i(F(x, G(x), \vec{G}(x, G(x)))) \leq \lambda_i$ by (c).

This completes the proof.

As an immediate consequence of the above we have

5.4 Theorem : For each $\rho < \omega_1$, $\rho \geq 1$, $c\mathcal{R}^\rho$ has the scale property. \square

The proofs of the scale property give us the following uniformization results.

5.5 Theorem : Let \mathbb{I} be a p.a.o. which subsumes \mathcal{A} . If $\underline{\nabla}(\mathbb{I})$ is presentable, then every $E \subseteq X \times Y$ in $\underline{\nabla}(\mathbb{I})$ can be uniformized by a $\underline{\nabla}(\mathbb{I})$ set E^* .

In particular, for each $\rho < \omega_1$, \underline{BR}^ρ has the uniformization property.

Proof : Let $E \subseteq X \times Y$ be in $\underline{\nabla}(\mathbb{I})$. By hypothesis E has a $\underline{\nabla}(\mathbb{I})$ -presentation (Z, B, P, G) . As above, let $F_3 : P \cap B \rightarrow P \cap B$ be an analytically measurable function which picks up a unique point from every vertical line $\{(x, y)\} \times (P \cap B)_{(x, y)}$ and such that the associated cross-section P^* is \prod_1^1 . Again, let $F_2 : P^* \rightarrow P^*$ be another analytically measurable function which from every vertical line $\{x_1\} \times P_{x_1}^*$ picks up a unique point and whose associated cross-section P^{**} is \prod_1^1 . Let $E^{**} = \text{proj}_{X \times Y} (P^{**})$. Clearly, E^{**} is a graph which uniformizes E . We claim that $E^{**} \in \underline{\nabla}(\mathbb{I})$. To see this define $H : X \times Y \rightarrow X \times Y \times Z$ by

$$H(x, y) = F_3(x, y, G(x, y)).$$

Plainly, H is $\underline{\nabla}(\mathbb{I})$ -measurable, since $\mathbb{I} \geq \mathcal{A}$.

Moreover, $E^{**} = H^{-1}(P^{**})$,

and hence $E^{**} \in \underline{\nabla}(\mathbb{I})$. \square

Remark : Since Borel-programmable (absolutely Δ_2^1 ; Σ_2^1) sets are presentable (cf [10]) the above proof shows that Borel-programmable (absolutely Δ_2^1 ; Σ_2^1) sets can be uniformized by a Borel-programmable (absolutely Δ_2^1 ; Σ_2^1) graph.

5.6 Theorem : Suppose $\mathbb{I} \geq \mathcal{A}$ and $\Sigma(\mathbb{I})$ is presentable. Then every $E \subseteq \mathcal{X}_1 \times \mathcal{X}_2$ in $\prod_1^{\mathbb{I}^*}$ can be uniformized by a set in $\prod_1^{\mathbb{I}^*}$.

In particular, for every $\rho < \omega_1$, $c\mathcal{R}^\rho$ has the uniformization property.

Proof : Let $E \subseteq \mathcal{X}_1 \times \mathcal{X}_2$ be in $\prod_1^{\mathbb{I}^*}$. Let C, G be as in the proof of Theorem 5.3 with \mathcal{X} replaced by $\mathcal{X}_1 \times \mathcal{X}_2$. Suppose (Z, B, P, \tilde{G}) be a $\nabla(\mathbb{I})$ -presentation of C as in the proof there. By Lemma 5.1, let $F_3 : P \cap B \rightarrow P \cap B$ be an analytically measurable function such that it picks up a unique point from every vertical line $\{(x_1, x_2)\} \times (P \cap B)_{x_1, x_2}$ and whose associated cross-section P^* is \prod_1^1 . By Lemma 5.1 again, let $F_2 : P^* \rightarrow P^*$ be a similar function picking up a point from each vertical line $\{x_1\} \times (P^*)_{x_1}$ such that $P^{**} = \text{range}(F_2)$ is \prod_1^1 . Let

$$E^{**} = \text{proj}_{\mathcal{X}_1 \times \mathcal{X}_2} (P^{**}).$$

Plainly, E^{**} is a graph which uniformizes E . We shall show that $E^{**} \in \prod_1^{\mathbb{I}^*}$. Define a partial function

$$H : \mathcal{X}_1 \times \mathcal{X}_2 \rightarrow \mathcal{X}_1 \times \mathcal{X}_2$$

with domain E as follows

$$H(x_1, x_2) = F_3(x_1, x_2, G(x_1, x_2), \tilde{G}(x_1, x_2, G(x_1, x_2))) ; (x_1, x_2) \in E.$$

Note that H is a partial $\prod_1^{\mathbb{I}^*}$ -recursive function. We claim that

$$E^{**}(x_1, x_2) \leftrightarrow H(x_1, x_2) \downarrow \& P^{**}(H(x_1, x_2)) \quad (i)$$

To see this assume $(x_1, x_2) \in E^{**}$. Then there exist (unique) α, γ, z

such that $(x_1, x_2, \alpha, \gamma, z) \in P^{**}$. Now since $(x_1, x_2) \in E$, $H(x_1, x_2) \downarrow$.

Further,

$$H(x_1, x_2) = F_3(x_1, x_2, G(x_1, x_2), \overline{G}(x_1, x_2, G(x_1, x_2))) \in P^*.$$

So for some α', γ', z' ,

$$H(x_1, x_2) = (x_1, x_2, \alpha', \gamma', z') \in P^*.$$

Since $(x_1, x_2, \alpha, \gamma, z)$ and $(x_1, x_2, \alpha', \gamma', z')$ are in P^* it follows that $\alpha = \alpha', \gamma = \gamma'$ and $z = z'$. Thus

$$H(x_1, x_2) = (x_1, x_2, \alpha, \gamma, z) \in P^{**}.$$

Conversely, suppose $H(x_1, x_2) \in P^{**}$ so that for some α, γ, z ,

$(x_1, x_2, \alpha, \gamma, z) \in P^{**}$. But

$$E^{**} = \text{proj}_{X_1 \times X_2} (P^{**}),$$

and so $(x_1, x_2) \in E^{**}$. Thus we have (i). By the substitution property

this shows that $E^{**} \in \prod_1 \mathbb{I}^*$. \square

6. The Category and the Measure Formulas for R-sets : As promised

earlier we shall exhibit another property which is left invariant by the operator \mathbb{R} . Roughly speaking, we shall show that the $*$ -transform of Vaught (see 6.1) is preserved by the operator \mathbb{R} . More explicitly, we shall show that if a set $E \subseteq X \times Y$ obtained by the application of \mathbb{I} is such that $E^* = \{x : E^x \text{ is comeager}\}$ is computed by the application of Ψ , then for any set F computed by \mathbb{I}^* , F^* is computed by Ψ^* .

This invariance property is established by first obtaining a decomposition

of E^* for sets $E \in \Sigma_1^1 \bar{\mathcal{Q}}^*$, analogous to the one obtained by Vaught for Σ_1^1 sets [51] and for \mathcal{R}^1 -sets obtained by Burgess and Miller [8]. Although our methods for computing E^* are implicit in Burgess' proof for the same (cf [8]), we have, by isolating the core of his proof (viz. Theorem 6.3 below), been able to compute E^* for all levels of the hierarchy of \mathcal{R} -sets by a simple inductive argument; and thus, been able to avoid the notational complexities involved in Burgess' proof.

We first introduce the following definition due to Vaught [51].

6.1 Definition: A set $E \subseteq X \times Y$, the product of two Polish spaces, is said to be normal if for each $x \in X$, $E^x = \{y \in Y : (x, y) \in E\}$ has the Baire property. (In this section E^x denotes the vertical section at x). If E is normal and $U \subseteq Y$ is open, then define

$$E^{*U} = \{x \in X : E^x \text{ is comeager in } U\}.$$

If $U = Y$, we write E^* instead of E^{*U} .

Next we get a decomposition of E^* for sets E obtained by the application of $\bar{\mathcal{Q}}^*$ on a normal family. (Recall that $\Sigma(s)$ denotes the set of all reals α extending s).

6.2 Lemma: Let $\bar{\mathcal{Q}}$ be a p.a.o. which preserves the Baire property and let $E = \bar{\mathcal{Q}}^*(\{E_s : s \in \omega\})$, with each $E_s \subseteq \omega^\omega \times \omega^\omega$ normal. Define E_s^μ and T_s^μ as in 2.6. Then for any $s \in \text{Seq}$,

$$E^{*\Sigma(s)} = \bigcap_{\mu < \omega_1} [E_s^\mu]^{*\Sigma(s)} = \bigcup_{\mu < \omega_1} [E_s^\mu - T_s^\mu]^{*\Sigma(s)}.$$

Proof: As in Theorem 2.8, one can easily check that E_s^μ and T_s^μ are

normal for each μ . Since $E = \bigcap_{\mu < \omega_1} E_\theta^\mu$, it follows that

$$E^{*\Sigma(s)} \subseteq \bigcap_{\mu < \omega_1} [E_\theta^\mu]^{*\Sigma(s)}. \quad \dots(i)$$

Next, suppose $x \in [E_\theta^\mu]^{*\Sigma(s)}$ for all $\mu < \omega_1$. For each $p \in \mathbb{N}$, $\{(E_p^\mu)^x : \mu < \omega_1\}$ is a decreasing sequence of sets with Baire property. Hence by the countable chain condition, $\exists \beta(p) < \omega_1$ such that

$$(\forall \rho > \beta(p)) [(E_p^{\beta(p)} - E_p^\rho)^x \text{ is meager}].$$

Choose ρ_0 such that $\beta(p) < \rho_0 < \omega_1$, for all p . Then

$(\forall p) [(E_p^{\rho_0} - E_p^{\rho_0+1})^x \text{ is meager}]$ and hence $(T^{\rho_0})^x$ is meager. Since $x \in [E_\theta^{\rho_0}]^{*\Sigma(s)}$, $(E_\theta^{\rho_0})^x$ is comeager in $\Sigma(s)$. Therefore, $(E_\theta^{\rho_0})^x - (T^{\rho_0})^x$ is comeager in $\Sigma(s)$ and so $x \in [E_\theta^{\rho_0} - T^{\rho_0}]^{*\Sigma(s)}$.

Thus,

$$\bigcap_{\mu < \omega_1} [E_\theta^\mu]^{*\Sigma(s)} \subseteq \bigcup_{\mu < \omega_1} [E_\theta^\mu - T^\mu]^{*\Sigma(s)}. \quad \dots(ii)$$

Finally, since $(E_\theta^\mu - T^\mu) \subseteq E$ for each μ ,

$$\bigcup_{\mu < \omega_1} [E_\theta^\mu - T^\mu]^{*\Sigma(s)} \subseteq E^{*\Sigma(s)}. \quad \dots(iii)$$

The result now follows from (i) - (iii).

6.3 Transfer Theorem : Let \mathbb{I} and \mathbb{J} be two positive analytical operations such that \mathbb{I} preserves the Baire property and \mathbb{J} is normal and subsumes both (countable) \cup and \cap . Suppose, moreover, that there

are (recursive) functions f and g such that for any normal family $\{E_p : p \in \omega\}$ of subsets of $\omega^\omega \times \omega^\omega$ with $E = \bigcap (\{E_p : p \in \omega\})$,

$$E^{*\Sigma}(s) = \bigcup (\{E_{f(p)}^{*\Sigma}(s * g(p)) : p \in \omega\}).$$

Then for any normal family $\{F_p : p \in \omega\}$ of subsets of $\omega^\omega \times \omega^\omega$,

(a) $F = \bigcap^0 (\{F_p : p \in \omega\})$ implies that

$$F^{*\Sigma}(s) = \bigcup^0 (\{F_{\gamma(p)}^{*\Sigma}(s \hat{\delta}(p)) : p \in \omega\}),$$

for suitable (recursive) functions γ and δ (independent of the family $\{F_p\}$).

(b) $F = \bigcap \bigcap^0 (\{F_p : p \in \omega\})$ implies

$$F^{*\Sigma}(s) = \bigcup \bigcup^0 (\{F_{\alpha(p)}^{*\Sigma}(s \hat{\beta}(p)) : p \in \omega\}),$$

where

$$\alpha(s) = \begin{cases} \langle f(n), \gamma(m) \rangle & \text{if } s = \langle n, m \rangle, \\ 0 & \text{otherwise;} \end{cases}$$

$$\beta(s) = \begin{cases} g(n) \hat{\delta}(m) & \text{if } s = \langle n, m \rangle, \\ 0 & \text{otherwise.} \end{cases}$$

(c) $F = \bigcap^* (\{F_p : p \in \omega\})$ implies

$$F^{*\Sigma}(s) = \bigcup^* (\{F_{\tilde{f}(p)}^{*\Sigma}(s \hat{\tilde{g}}(p)) : p \in \omega\}),$$

where

$$\tilde{f}(s) = \begin{cases} \langle \alpha(n_0), \dots, \alpha(n_{k-1}) \rangle & \text{if } s = \langle n_0, \dots, n_{k-1} \rangle, \\ 0 & \text{otherwise;} \end{cases}$$

$$\tilde{g}(s) = \begin{cases} \beta(n_0) \hat{\beta}(n_1) \dots \hat{\beta}(n_{k-1}) & \text{if } s = \langle n_0, \dots, n_{k-1} \rangle, \\ 0 & \text{otherwise.} \end{cases}$$

Proof: Let N and M be the canonical bases for \mathbb{I} and \mathbb{J} , respectively.

(a) Set $G = F^0$ and $G_p = F_p^0$ for each p .

Then $G = \mathbb{I}(\{G_p : p \in \omega\})$; and since \mathbb{I} preserves the Baire property, each G^x has the Baire property. Therefore,

G^x is nonmeager in $\Sigma(s)$

$$\leftrightarrow (\exists u \in \text{Seq}) [G^x \text{ is comeager in } G(s \hat{\cup} u)]$$

$$\leftrightarrow (\exists u \in \text{Seq})(\exists n \in \mathbb{N})(\forall m \in \mathbb{N}) [G_{p(m)}^x \text{ is comeager in } \Sigma(s \hat{\cup} \gamma(m))] \text{ by hypothesis}$$

$$\leftrightarrow (\exists u \in \text{Seq})(\exists n \in \mathbb{N})(\forall m \in \mathbb{N})(\forall v \in \text{Seq}) [G_{p(m)}^x \text{ is nonmeager in } \Sigma(s \hat{\cup} \gamma(m) \hat{\cup} v)]$$

$$\leftrightarrow (\exists n \in \mathbb{N})(\forall m \in \mathbb{N}) [G_{p(m)}^x \text{ is nonmeager in } \Sigma(s \hat{\cup} \delta(m))].$$

for some (recursive) functions γ and δ , as $\mathbb{I}_M \geq \cup, \cap$ and is normal.

Hence,

$$F^x \text{ is comeager in } \Sigma(s) \leftrightarrow G^x \text{ is meager in } \Sigma(s)$$

$$\leftrightarrow (\forall n \in \mathbb{N})(\exists m \in \mathbb{N}) [G_{p(m)}^x \text{ is meager in } \Sigma(s \hat{\cup} \delta(m))]$$

$$\leftrightarrow (\forall n \in \mathbb{N})(\exists m \in \mathbb{N}) [F_{p(m)}^x \text{ is comeager in } \Sigma(s \hat{\cup} \delta(m))]$$

$$\leftrightarrow (\exists n \in \mathbb{N})(\forall m \in \mathbb{N}) [F_{p(m)}^x \text{ is comeager in } \Sigma(s \hat{\cup} \delta(m))].$$

This proves (a).

To prove (b), observe that

F^x is comeager in $\Sigma(s)$

$$\leftrightarrow [\mathbb{I}(\{\mathbb{I}^0(\{F_{\langle p, q \rangle} : q \in \omega\}) : p \in \omega\})]^x \text{ is comeager in } \Sigma(s)$$

$\leftrightarrow (\exists \eta \in M)(\forall n \in \eta) [(\mathbb{I}^\circ (\{F_{\langle f(n), q \rangle} : q \in \omega\})^X \text{ is comeager in } \Sigma(\hat{e}^g(n))]$ (by hypothesis),

$\leftrightarrow (\exists \eta \in M)(\forall n \in \eta)(\exists \xi \in M^\circ)(\forall m \in \xi) [F_{\langle f(n), \gamma(m) \rangle}^X \text{ is comeager in } \Sigma(\hat{e}^g(n)\hat{\delta}(m))]$ (by (a)).

Setting $\alpha(\langle n, m \rangle) = \langle f(n), \gamma(m) \rangle$, $\beta(\langle n, m \rangle) = g(n)\hat{\delta}(m)$, we get (b),

(c) Since $F = \mathbb{I}^* (\{F_p : p \in \omega\})$, the canonical inductive operator Γ is given by

$$p \in \Gamma_z(A) \leftrightarrow z \notin F_p \vee (\forall \eta \in N)(\exists n \in \eta)(\exists \xi \in N)(\forall m \in \xi) [p * \langle \langle n, m \rangle \rangle \in A] ; z \in \omega^\omega \times \omega^\omega.$$

Define $z \in F_p^\mu \leftrightarrow p \notin \Gamma_z^\mu$. Then by Lemma 6.2,

$$F^* \Sigma(s) = \bigcap_{\mu < \omega_1} [F_s^\mu]^* \Sigma(s). \quad \dots(i)$$

Define a set relation Δ operative on ω as follows :

$$t \in \Delta_x(A) \leftrightarrow F_{(t)_0}^X \text{ is not comeager in } \Sigma((t)_1) \vee (\forall \eta \in M)(\exists n \in \eta)(\exists \xi \in M)(\forall m \in \xi) [\langle (t)_0 \hat{\langle} \langle f(n), \gamma(m) \rangle \rangle, (t)_1 \hat{\langle} g(n)\hat{\delta}(m) \rangle \in A] .$$

We shall show that for any $s \in \text{Seq}$

$$x \notin F^* \Sigma(s) \leftrightarrow \langle s, s \rangle \in \Delta_x^\omega. \quad \dots(ii)$$

To prove (ii) we shall show by induction on μ that for any $s, t \in \text{Seq}$,

$$x \notin [F_t^\mu]^* \Sigma(s) \leftrightarrow \langle t, s \rangle \in \Delta_x^\mu.$$

This is clearly true for $\mu = 0$, so assume $\mu > 0$. Now,

$$x \notin [F_t^\mu]^{\ast\Sigma(s)} \iff (F_t^\mu)^x \text{ is not comeager in } \Sigma(s)$$

$$\iff F_t^x \text{ is not comeager in } \Sigma(s) \vee \left[\bigcap_{\lambda < \mu} F_t^\lambda \right]^x \text{ is not comeager in } \Sigma(s)$$

is not comeager in $\Sigma(s)$

$$\iff F_t^x \text{ is not comeager in } \Sigma(s) \vee (\forall \eta \in M)(\exists n \in \eta)(\exists \xi \in M)(\forall m \in \xi)$$

$$\left[\left(\bigcap_{\lambda < \mu} F_t^\lambda \right)^x \text{ is not comeager in } \Sigma(s \hat{g}(n) \hat{\delta}(m)) \right]$$

(by (b))

$$\iff F_t^x \text{ is not comeager in } \Sigma(s) \vee (\forall \eta \in M)(\exists n \in \eta)(\exists \xi \in M)(\forall m \in \xi)$$

$$(\exists \lambda < \mu) \left[(F_t^\lambda)^x \text{ is not comeager in } \Sigma(s \hat{g}(n) \hat{\delta}(m)) \right]$$

$$\iff F_t^x \text{ is not comeager in } \Sigma(s) \vee (\forall \eta \in M)(\exists n \in \eta)(\exists \xi \in M)(\forall m \in \xi)$$

$$(\exists \lambda < \mu) \left[\langle t^{\hat{g}(n)} \hat{\delta}(m) \rangle \in \Delta_x^\lambda \right] \text{ by the}$$

induction hypothesis

$$\iff \langle t, s \rangle \in \Delta_x^\mu.$$

Hence, putting $t = s$ and using (i), we obtain (ii). Therefore,

$$F^x \text{ is comeager in } \Sigma(s) \iff x \in F^{\ast\Sigma(s)} \iff \langle s, s \rangle \notin \Delta_x^\mu \text{ by (ii)}$$

$$\iff (\exists \eta_0 \in M)(\forall n_0 \in \eta_0)(\forall \xi_0 \in M)(\exists m_0 \in \xi_0) \dots$$

$$\dots (\forall k) \left[F^x \text{ is comeager in } \left[\langle f(n_0), \gamma(m_0) \rangle, \dots, \langle f(n_{k-1}), \gamma(m_{k-1}) \rangle \right] \right]$$

$$\Sigma(s \hat{g}(n_0) \hat{\delta}(m_0) \dots \hat{g}(n_{k-1}) \hat{\delta}(m_{k-1}))$$

Define \tilde{f} and \tilde{g} as follows :

$$\tilde{f}(s) = \begin{cases} \langle \alpha(n_0), \dots, \alpha(n_{k-1}) \rangle & \text{if } s = \langle n_0, \dots, n_{k-1} \rangle, \\ 0 & \text{otherwise} \end{cases}$$

$$g(s) = \begin{cases} \beta(n_0) \hat{\beta}(n_1) \dots \hat{\beta}(n_{k-1}) & \text{if } s = \langle n_0, \dots, n_{k-1} \rangle, \\ 0 & \text{otherwise.} \end{cases}$$

The result now follows immediately.

Remark : It is clear from the proof of 6.3(c) that, under the hypothesis of the theorem, for any normal family $\{F_p\}$ of subsets of $\omega^\omega \times \omega^\omega$ with $F = \mathbb{R} \tilde{\mathbb{I}}_N(\{F_p\})$,

$$F^{*\Sigma}(s) = \mathbb{R} \tilde{\mathbb{I}}_M(\{F_{f'(p)}^{*\Sigma}(\hat{s}g'(p)) : p \in \omega\}),$$

where

$$f'(\langle n_0, \dots, n_{k-1} \rangle) = \langle f(n_0), \dots, f(n_{k-1}) \rangle,$$

$$g'(\langle n_0, \dots, n_{k-1} \rangle) = g(n_0) \hat{g}(n_1) \dots \hat{g}(n_{k-1}).$$

Here, normality of $\tilde{\Psi} = \tilde{\mathbb{I}}_M$ is not required.

We shall now apply the Transfer Theorem to deduce Vaught's Formula for E^* (cf [5] Theorem 1.6] and the Category Formula of Burgess [13]).

6.4 Theorem (Vaught) : Assume $E = \mathcal{A}(\{E_n : n \in \omega\})$, with each $E_n \subseteq \omega^\omega \times \omega^\omega$ normal. Then $x \in E^{*\Sigma}(s)$ if and only if

$$(\forall u_0 \in \text{seq})(\exists v_0 \in \text{seq})(\exists k_0)(\forall u_1 \in \text{seq})(\exists v_1 \in \text{seq})(\exists k_1) \dots$$

$$\dots (\forall i) \left[E^x_{\langle k_0, \dots, k_{i-1} \rangle} \text{ is comeager in } \Sigma(\hat{s}u_0 \hat{v}_0 \dots \hat{u}_{i-1} \hat{v}_{i-1}) \right].$$

Proof : Let $N = \{\{n\} : n \in \omega\}$ so that $\bar{\mathcal{I}}_N = U$. Suppose

$F = \bar{\mathcal{I}}_N(\{F_n : n \in \omega\}) = \bigcup_{n \in \omega} F_n$, where each F_n is a normal subset $\omega^\omega \times \omega^\omega$. Then,

$$F^X \text{ is comeager in } \Sigma(s) \iff (\forall u \in \text{Seq}) [F^X \text{ is nonmeager in } \Sigma(s^{\wedge}u)]$$

$$\iff (\forall u \in \text{Seq})(\exists k) [F_k^X \text{ is nonmeager in } \Sigma(s^{\wedge}u)]$$

$$\iff (\forall u \in \text{Seq})(\exists k)(\exists v \in \text{Seq}) [F_k^X \text{ is comeager in } \Sigma(s^{\wedge}v)].$$

Let $\bar{\mathcal{I}}_M$ be a δ -s operation such that for any family $\{A_n\}$,

$$\bar{\mathcal{I}}_M(\{A_n : n \in \omega\}) = \bigcap_{u \in \text{Seq}} \bigcup_{v \in \text{Seq}} \bigcup_{k \in \omega} A_{\langle u, \langle k, v \rangle \rangle}.$$

Hence,

F^X is comeager in $\Sigma(s)$

$$\iff (\exists \eta \in M)(\forall n \in \eta) [F_{(n)_{1,0}}^X \text{ is comeager in } \Sigma(s^{\wedge}(n)_{0,1}^{\wedge}(n)_{1,1})]$$

$$\iff (\exists \eta \in M)(\forall n \in \eta) [F_{f(n)}^X \text{ is comeager in } \Sigma(s^{\wedge}g(n))],$$

where $f(n) = (n)_{1,0}$ and $g(n) = (n)_{0,1}^{\wedge}(n)_{1,1}$. Since

$$E = \bar{\mathcal{A}}(\{E_n : n \in \omega\}) = R\bar{\mathcal{I}}_M(\{E_n\}),$$

the Transfer Theorem immediately gives

$$E^* \Sigma(s) = R\bar{\mathcal{I}}_M(\{E_{\bar{f}(p)}^* \Sigma(s^{\wedge}g(p)) : p \in \omega\}),$$

where $\bar{g}(\langle n_0, \dots, n_{k-1} \rangle) = g(n_0)g(n_1)^{\wedge} \dots^{\wedge} g(n_{k-1})$

and $\bar{f}(\langle n_0, \dots, n_{k-1} \rangle) = \langle f(n_0), \dots, f(n_{k-1}) \rangle$. Therefore,

E^X is comeager in $\Sigma(s)$

$$\iff (\exists \eta_0 \in M)(\forall n_0 \in \eta_0)(\exists \eta_1 \in M)(\forall n_1 \in \eta_1) \dots$$

$$\dots (\forall i) [E_{\langle \bar{f}(n_0), \dots, \bar{f}(n_{i-1}) \rangle}^X \text{ is comeager in } \Sigma(s^{\wedge}g(n_0)^{\wedge} \dots^{\wedge} g(n_{i-1}))]$$

$$\begin{aligned} &\leftrightarrow (\forall u_0 \in \text{Seq})(\exists v_0 \in \text{Seq})(\exists k_0)(\forall u_1 \in \text{Seq})(\exists v_1 \in \text{Seq})(\exists k_1) \dots \dots \\ &\dots (\forall i) \left[E^x_{\langle k_0, \dots, k_{i-1} \rangle} \text{ is comeager in } \Sigma(\hat{u}_0 \hat{v}_0 \dots \hat{u}_{i-1} \hat{v}_{i-1}) \right]. \end{aligned}$$

This completes the proof.

6.5 Definition: Define an operation \mathcal{V} as follows.

$$x \in \mathcal{V}(\{E_n : n \in \omega\})$$

$$\begin{aligned} &\leftrightarrow (\forall u_0 \in \text{Seq})(\exists v_0 \in \text{Seq})(\exists k_0)(\forall u_1 \in \text{Seq})(\exists v_1 \in \text{Seq})(\exists k_1) \dots \dots \\ &\dots (\forall i) \left[x \in E_{\langle \langle u_0, \langle k_0, v_0 \rangle \rangle, \dots, \langle u_{i-1}, \langle k_{i-1}, v_{i-1} \rangle \rangle} \right]. \end{aligned}$$

Clearly \mathcal{V} is positive analytical and $\mathcal{V} \prec \mathcal{A}$. Call \mathcal{V} the Vaught operation.

Define a sequence of positive analytical operations $\{S_\rho : \rho < \omega_1\}$ by the induction

$$S_0 = \mathcal{V}, \quad S_{\rho+1} = S_\rho^*.$$

If λ is limit, choose $\rho_i \uparrow \lambda$ and set for any family $\{E_n\}$,

$$\Psi(\{E_n\}) = \bigcap_{i \in \omega} \bigcap_{M_{\rho_i} M_{\rho_i}^0} (\{E_{\langle i, m \rangle} : m \in \omega\}),$$

where M_{ρ_i} is the canonical base for S_{ρ_i} . Then define $S_\lambda = \Psi^*$. It is easy to check by induction that for each ρ , $R_\rho \prec S_\rho$ and hence

$$\underline{R}^\rho = \sum_{\prec 1}^{R_\rho} = \sum_{\prec 1}^{S_\rho}.$$

We shall now deduce the Category Formula of Burgess by showing that if E is in \underline{R}^ρ , the E^* is computed by S_ρ .

6.6 Theorem : Let $E \subseteq \omega^\omega \times \omega^\omega$ be a set in \mathcal{R}^ρ . Then $E^{*\Sigma(s)} = \{x : E^x \text{ is comeager in } \Sigma(s)\}$ is also in \mathcal{R}^ρ .

Proof : We shall prove the theorem by induction. Let N_ρ and M_ρ denote the canonical bases for R_ρ and S_ρ , respectively.

Assume that for all $\lambda < \rho$ there are functions f_λ and g_λ such that if

$$F = R_\lambda (\{F_n : n \in \omega\}),$$

with each $F_n \subseteq \omega^\omega \times \omega^\omega$ normal, then

$$F^{*\Sigma(s)} = S_\lambda (\{F_{f_\lambda(n)}^{*\Sigma(\hat{s}g_\lambda(n))} : n \in \omega\}).$$

We shall then show that if $E = R_\rho(\{E_n\})$, then $E^{*\Sigma(s)}$ is similarly computed by S_ρ . The result then follows by observing that for each clopen E_n , E_n^* is also clopen.

Case 1 : $\rho = \lambda + 1$. In this case, $E = R_\lambda^*(\{E_n : n \in \omega\})$. Hence by the Transfer Theorem,

E^x is comeager in $\Sigma(s) \iff (\exists n_0 \in M_\lambda)(\forall n_0 \in \eta_0)(\forall \xi_0 \in M_0)(\exists m_0 \in \xi_0) \dots$

$$\dots (\forall i) \left[E_{\tilde{f}_\lambda}^x (\langle \langle n_0, m_0 \rangle, \dots, \langle n_{i-1}, m_{i-1} \rangle \rangle) \text{ is comeager in } \Sigma(\hat{s} \tilde{g}_\lambda (\langle \langle n_0, m_0 \rangle, \dots, \langle n_{i-1}, m_{i-1} \rangle \rangle)) \right]$$

where \tilde{f}_λ and \tilde{g}_λ are related to f_λ and g_λ as in 6.3. Setting

$f_\rho = \tilde{f}_\lambda$ and $g_\rho = \tilde{g}_\lambda$ we get

$$E^{*\Sigma(s)} = S_\rho (\{E_{f_\rho(n)}^{*\Sigma(\hat{s}g_\rho(n))} : n \in \omega\}).$$

Case 2 : ρ is limit. Choose a sequence $\rho_i \uparrow \rho$ and for any normal family $\{H_n : n \in \omega\}$ set

$$H = \bigcap_{i \in \omega} \bigcap_{N_{\rho_i}^0} \bigcap_{N_{\rho_i}^0} (\{H_{\langle i, n \rangle} : n \in \omega\}).$$

Then,

H^x is comeager in $\Sigma(s)$

$$\iff (\forall i) \left[\left(\bigcap_{N_{\rho_i}^0} \bigcap_{N_{\rho_i}^0} (\{H_{\langle i, m \rangle} : m \in \omega\}) \right)^x \text{ is comeager in } \Sigma(s) \right]$$

$$\iff (\forall i) (\exists \eta \in M_{\rho_i}) (\forall n \in \eta) (\forall \xi \in M_{\rho_i}) (\exists m \in \xi) \left[H^x_{\langle i, \alpha_i(\langle n, m \rangle) \rangle} \text{ is comeager in } \Sigma(s^{\beta_i}(\langle n, m \rangle)) \right] \text{ (by Theorem 6.3)}$$

where α_i and β_i are obtained as in 6.3(b). Therefore,

$$H^{*\Sigma(s)} = \bigcap_{f(m)} \{H_{f(m)}^{*\Sigma(s^{\hat{g}(m)})} : m \in \omega\},$$

where \bigcap is the operation in 6.5 and

$$f(\langle i, n \rangle) = \langle i, \alpha_i(n) \rangle, \quad g(\langle i, n \rangle) = \beta_i(n).$$

Since $E = \bigcap^* (\{E_n\})$, by applying the Transfer Theorem again we obtain the result as in Case 1.

As the R -sets have the Baire property the next result follows immediately.

6.7 Corollary : If $E \in \underline{cR}^\rho$, then $E^{*\Sigma(s)}$ is also in \underline{cR}^ρ .

By 6.6 and by repeatedly applying the Transfer Theorem to every level of the hierarchy of \underline{BR}^ρ -sets one obtains

6.8 Corollary : If $E \in \mathcal{BR}_{\mu}^{\rho}(\mu, \rho < \omega_1)$, then $E^{*\Sigma(s)}$ is also in $\mathcal{BR}_{\mu}^{\rho}$. In particular, putting $\rho = 0$, one has for any C-set E, E^* is also a C-set.

6.9 We will now establish the measure-theoretic counterparts of Theorem 6.3 and its corollaries.

Throughout the rest of this section $G \subseteq \omega^{\omega} \times \omega^{\omega}$ will be a fixed good Σ_1^0 universal set which is ω -universal for Σ_1^0 subsets of ω^{ω} , and μ will denote the Lebesgue measure on ω^{ω} as defined in [24] (or any other normalized, complete, regular measure with a reasonable definability condition). We shall also fix a recursive function b such that

$$G(n, x) \ \& \ G(m, x) \leftrightarrow G(b(n, m), x).$$

(We shall also denote by E^x the vertical section of E at x and think of the rationals coded in some recursive way by elements of ω).

Call a set $A \subseteq X \times Y$ normal measurable if A^x is measurable for each x .

For any normal measurable set $A \subseteq \omega^{\omega} \times \omega^{\omega}$ and any $n \in \omega, r \geq 0$ rational, we define

$$A^{*[n, r]} = \{x : \mu(A^x \cap G^n) \geq r\}.$$

Call $s \in \text{Seq}$ relevant if $\text{lh}(s) = 3$ & $(s)_0$ is rational & $(s)_2 \in \text{Seq}$.

Let s and t be two relevant sequences. Then t is an extension of s if

$$G^{(t)_1} \subseteq G^{(s)_1} \ \& \ (s)_2 \subseteq (t)_2 \ \& \ \text{lh}((t)_2) = \text{lh}((s)_2) + 1.$$

For any rational r , define

$$N_{\theta}^{(r)} = \left\{ \eta \subseteq \omega : \eta \text{ is a finite set of relevant sequences} \right. \\ \left. \begin{aligned} &\& \sum_{s \in \eta} (s)_0 > r \ \& \ (\forall s \in \eta) (\text{lh}(s)_2 = 1) \\ &\& \forall s, t \in \eta (s \neq t \rightarrow G^{(s)}_1 \cap G^{(t)}_1 = \emptyset) \end{aligned} \right\}.$$

If s_1, s_2, \dots, s_k are relevant sequences such that s_{i+1} is an extension of s_i and $\text{lh}((s_i)_2) = i$, then define

$$N_{\langle s_1, \dots, s_k \rangle} = \left\{ \eta \subseteq \omega : \eta \text{ is a finite set of extensions of } s_k \text{ such} \right. \\ \left. \text{that } \sum_{t \in \eta} (t)_0 > (s_k)_0 \ \& \ \forall t, u \in \eta (t \neq u \rightarrow G^{(t)}_1 \cap G^{(u)}_1 = \emptyset) \right\}.$$

We will now formulate and write down less detailed proofs of the measure analogues of the results we have proved for category.

Notice that the next result due to Burgess [13] and implicit in Kechris [24] gives us the Measure Formula at the base level.

6.10 Lemma : Let $\{E_n : n \in \omega\}$ be a family of normal measurable subsets of $\omega^\omega \times \omega^\omega$ and suppose

$$E = \bigcap \{E_n : n \in \omega\}.$$

Then for any rational $r \geq 0$,

$$\mu(E^X) > r \iff (\exists \eta_1 \in N_{\theta}^{(r)}) (\forall s_1 \in \eta_1) (\exists \eta_2 \in N_{\langle s_1 \rangle}) (\forall s_2 \in \eta_2) \dots \\ \dots (\forall k) \left[\mu(E_{(s_k)_2}^X \cap G^{(s_k)}_1) > (s_k)_0 \right]. \quad \square$$

To prove Lemma 6.10 one needs the following observation due to Burgess and may be regarded as the Measure Formula for countable union.

6.11 Lemma : Let $\{E_n : n \in \omega\}$ be a family of measurable subsets of ω^ω and let $E = \bigcup_{n \geq 0} E_n$. Plainly E is measurable and moreover for any rational $r \geq 0$, $n \in \omega$

$$\mu(E \cap G^n) > r \iff (\exists \eta \in N_{\langle s^* \rangle}) (\forall s \in \eta) [\mu(E_{(s)_2} \cap G^{(s)_1}) > (s)_0],$$

where $s^* = \langle r, n, \varepsilon \rangle$ and $N_{\langle s^* \rangle}$ is defined as in 6.9.

Proof : Suppose for some $\eta \in N_{\langle s^* \rangle}$ we have

$$(\forall s \in \eta) [\mu(E_{(s)_2} \cap G^{(s)_1}) > (s)_0].$$

Recall the definition of $N_{\langle s^* \rangle}$. Then observe that

$$\begin{aligned} \mu\left(\bigcup_{i \geq 0} (E_i \cap G^n)\right) &\geq \mu\left(\bigcup_{s \in \eta} (E_{(s)_2} \cap G^{(s)_1})\right) \\ &= \sum_{s \in \eta} \mu(E_{(s)_2} \cap G^{(s)_1}) > \sum_{s \in \eta} (s)_0 > r. \end{aligned}$$

To prove the reverse implication, assume without loss of generality that $G^n = \omega^\omega$. Since $\mu\left(\bigcup_{i \geq 0} E_i\right) > r$, we can find a $p \in \omega$ and $\varepsilon > 0$ such that

$$\mu\left(\bigcup_{i < p} E_i\right) > r + p\varepsilon.$$

Since E_i is measurable, it is easy to get a recursive (in fact, a finite union of basic open sets) B_i such that $\mu(E_i \Delta B_i) < \varepsilon$.

Set

$$B'_i = B_i - \bigcup_{j < i} B_j.$$

Clearly each B'_i is recursive and moreover,

$$\bigcup_{i < p} E_i - \bigcup_{i < p} E_i \cap B'_i \subseteq \bigcup_{i < p} (E_i \Delta B'_i).$$

Hence,

$$\mu\left(\bigcup_{i < p} E_i\right) - \mu\left(\bigcup_{i < p} E_i \cap B'_i\right) \leq p \varepsilon.$$

Consequently,

$$\sum_{i < p} \mu(E_i \cap B'_i) \geq \mu\left(\bigcup_{i < p} E_i\right) - p \varepsilon > r.$$

Let i_1, \dots, i_k be an enumeration of all those $i < p$ for which $\mu(E_i \cap B'_i) > 0$. Get rationals r_1, \dots, r_k such that

$$\mu(E_{i_j} \cap B'_{i_j}) > r_j \quad \text{and} \quad \sum_{j=1}^k r_j > r.$$

Get n_1, \dots, n_k such that $G^{n_j} = B'_{i_j}$.

Put $\eta = \{ \langle r_j, n_j, \langle i_j \rangle \rangle : 1 \leq j \leq k \}$.

Clearly $\eta \in N_{\langle s^* \rangle}$ and for each $s \in \eta$

$$\mu(F_{(s)_2} \cap G^{(s)_1}) > (s)_0. \quad \square$$

Lemma 6.11 yields Lemma 6.10 and the proof of the next theorem (6.12) will give us an idea of obtaining 6.10 from 6.11. (For details see [13]).

With notation as in 6.10, now define

$$x \in A_{\langle s_1, \dots, s_k \rangle} \leftrightarrow \mu(E_{(s_k)_2}^x \cap G^{(s_k)_1}) > (s_k)_0$$

and $x \in A_e \leftrightarrow \mu(E_e^x \cap G^{n^*}) > r,$

where $G^{n^*} = \omega^\omega$. Then the above equivalence (6.10) reduces to

$$\begin{aligned} \mu(E^x) > r \leftrightarrow & (\exists \eta_1 \in N_e^{(r)}) (\forall s_1 \in \eta_1) (\exists \eta_2 \in N_{\langle s_1 \rangle}) (\forall s_2 \in \eta_2) \dots \dots \\ & \dots (\forall k) [x \in A_{\langle s_1, \dots, s_{k-1} \rangle}]. \end{aligned}$$

The right side of the above equivalence clearly defines a p.a.o. whose canonical base we will denote by $M^{(r)}$. Thus we have

$$\begin{aligned} \mu(E^x) > r \leftrightarrow & (\exists \eta \in M^{(r)}) (\forall m \in \eta) [x \in A_m] \dots (9) \\ \leftrightarrow & (\exists \eta \in M^{(r)}) (\forall m \in \eta) [\mu(E_{f(m)}^x \cap G^{g(m)}) > h(r, m)], \end{aligned}$$

where f, g, h are (recursive) functions such that

$$f(e) = e, f(\langle s_1, \dots, s_k \rangle) = (s_k)_2;$$

$$g(e) = n^*, g(\langle s_1, \dots, s_k \rangle) = (s_k)_1;$$

$$h(r, e) = r, h(r, \langle s_1, \dots, s_k \rangle) = (s_k)_0.$$

This gives us a motivation to formulate the measure-theoretic counterpart of the Transfer Theorem 6.3 as follows.

6.12 Theorem: Let \mathbb{I} be a measurability preserving p.a.o.. Suppose for each $r \geq 0$, there is an operation $\Psi^{(r)} = \mathbb{I}_{M^{(r)}} \geq \cap, \cup$ such that whenever

$$F = \mathbb{I}(\{F_n : n \in \omega\}), \text{ with each } F_n \subseteq \omega^\omega \times \omega^\omega \text{ normal measurable,}$$

then

$$\mu(F^X) \geq r \leftrightarrow (\exists \eta \in M^{(r)})(\forall m \in \eta) \left[\mu(F_{f(m)}^X \cap G^{g(m)}) \geq h(r, m) \right],$$

where f, g, h are (recursive) functions from ω into ω . Then,

whenever

$$E = \mathbb{R}\overline{\mathbb{I}}(\{E_n : n \in \omega\}),$$

with $\{E_n\}$ a regular family of normal measurable sets, we have

$$\mu(E^X \cap G^n) \geq r$$

$$\leftrightarrow (\exists \eta_0 \in M^{(r)})(\forall m_0 \in \eta_0)(\exists \eta_1 \in M^{(\tilde{h}(r, \langle m_0 \rangle)})(\forall m_1 \in \eta_1) \dots$$

$$\dots (\forall k) \left[\mu(E^X_{\langle f(m_0), \dots, f(m_{k-1}) \rangle} \cap G^{\tilde{b}(\langle m_0, \dots, m_{k-1} \rangle, n)}) \geq \tilde{h}(r, \langle m_0, \dots, m_{k-1} \rangle) \right],$$

where \tilde{b} and \tilde{h} are defined by recursion as follows :

$$\tilde{b}(e, n) = n,$$

$$\tilde{b}(\langle m_0, \dots, m_k \rangle, n) = b(\tilde{b}(\langle m_0, \dots, m_{k-1} \rangle, n), g(m_k)),$$

$$\tilde{b}(\lambda, n) = 0 \quad \text{otherwise.}$$

$$\tilde{h}(r, e) = r,$$

$$\tilde{h}(r, \langle m_0, \dots, m_k \rangle) = h(\tilde{h}(r, \langle m_0, \dots, m_{k-1} \rangle), m_k),$$

$$\tilde{h}(r, \lambda) = 0 \quad \text{otherwise.}$$

Proof : Let N be the canonical base for $\overline{\mathbb{I}}$ and let Γ be the inductive operator associated with N and $\{E_n : n \in \omega\}$. Define E_s^μ and T^μ as in 2.6. Then

$$E = \bigcap_{\rho < \omega_1} E_\theta^\rho = \bigcup_{\rho < \omega_1} (E_\theta^\rho - T^\rho).$$

Using this decomposition one can show as before that

$$E^*[n,r] = \bigcap_{\rho < \omega_1} (E_\theta^\rho)^*[n,r] \quad \dots(i)$$

Now define a set relation ψ as follows

$$\langle s, n, r \rangle \in \psi_x(A) \iff \mu(E_s^x \cap G^n) < r \\ \vee (\forall \eta \in M^{(r)})(\exists m \in \eta) [\langle s * \langle f(m) \rangle, b(g(m), n), h(r, m) \rangle \in A].$$

We shall show by induction on ρ that for every $s \in \text{Seq}$

$$x \notin (E_s^\rho)^*[n,r] \iff \langle s, n, r \rangle \in \psi_x^\rho.$$

First observe that since $\{E_n : n \in \omega\}$ is regular,

$$E_s^\rho = E_s \cap \bigcap_N (\{ \bigcap_{\lambda < \rho} E_{s * \langle m \rangle}^\lambda : m \in \omega \}) \\ = \bigcap_N (\{ \bigcap_{\lambda < \rho} E_{s * \langle m \rangle}^\lambda : m \in \omega \}). \quad \dots(ii)$$

Hence,

$$x \notin (E_s^\rho)^*[n,r] \iff \mu((E_s^\rho)^x \cap G^n) < r \\ \iff (\forall \eta \in M^{(r)})(\exists m \in \eta) [\mu(\bigcap_{\lambda < \rho} E_{s * \langle f(m) \rangle}^\lambda \cap G^{b(g(m), n)}) < h(r, m)] \\ \text{by (ii) and hypothesis of the theorem} \\ \iff (\forall \eta \in M^{(r)})(\exists m \in \eta)(\exists \lambda < \rho) [\mu(E_{s * \langle f(m) \rangle}^\lambda \cap G^{b(g(m), n)}) < h(r, m)]$$

$$\Leftrightarrow (\forall \eta \in M^{(r)})(\exists m \in \eta)(\exists \lambda < \rho) \left[x \notin (E_{s^* \langle f(m) \rangle}^\lambda)^* [b(g(m), n), h(r, m)] \right]$$

$$\Leftrightarrow (\forall \eta \in M^{(r)})(\exists m \in \eta)(\exists \lambda < \rho) \left[\langle s^* \langle f(m) \rangle, b(g(m), n), h(r, m) \rangle \in \psi_X^\lambda \right]$$

by induction hypothesis

$$\Leftrightarrow \langle s, n, r \rangle \in \psi_X^\rho.$$

Thus by (i) we have

$$x \notin E^*[n, r] \Leftrightarrow \langle s, n, r \rangle \in \psi_X^\omega.$$

Consequently,

$$\begin{aligned} \mu(E^X \cap G^n) \geq r &\Leftrightarrow \langle s, n, r \rangle \notin \psi_X^\omega \\ &\Leftrightarrow (\exists \eta_0 \in M^{(r)})(\forall m_0 \in \eta_0)(\exists \eta_1 \in M^{(\tilde{h}(r, \langle m_0 \rangle)})(\forall m_1 \in \eta_1) \dots \\ &\dots (\forall k) \left[\mu(E_{\langle f(m_0), \dots, f(m_{k-1}) \rangle}^X) \cap G^{\tilde{b}(\langle m_0, \dots, m_{k-1} \rangle, n)} \right. \\ &\quad \left. \geq \tilde{h}(r, \langle m_0, \dots, m_{k-1} \rangle) \right]. \quad \square \end{aligned}$$

An immediate consequence of Theorem 6.12 is the following

6.13 Theorem : For every $\rho < \omega_1$ and $r \geq 0$, the sets

$$\{x : \mu(E^X) > r\} \text{ and } \{x : \mu(E^X) \geq r\}$$

are in \mathcal{R}^ρ whenever $E \subseteq X \times Y$ is in \mathcal{R}^ρ .

The proof of the above theorem is exactly along the lines of the proof of Theorem 6.6. One shows by induction that for each $\rho < \omega_1$ and every rational $r \geq 0$ there are functions $f_\rho, f'_\rho, g_\rho, g'_\rho, h_\rho, h'_\rho$ and operations $\Phi_{M_\rho^r}, \Psi_{M_\rho^r} \rightarrow R_\rho$ such that whenever

$F = R_\rho(\{F_n : n \in \omega\})$, $F_n \subseteq X \times Y$ being normal measurable

$$\mu(F^X) > r \iff (\exists \eta \in M_\rho^r)(\forall n \in \eta) \left\{ \mu(F_{f_\rho(n)}^X \cap G^{g_\rho(n)}) > h_\rho(r, n) \right\}$$

$$\text{and } \mu(F^X) \geq r \iff (\exists \eta \in M_\rho^{r'})(\forall n \in \eta) \left\{ \mu(F_{f_\rho(n)}^X \cap G^{g_\rho(n)}) \geq h_\rho(r, n) \right\}.$$

By 6.13 and repeated applications of 6.12 we have

6.14 Corollary : If $E \in \underline{BR}_\lambda^\rho(\rho, \lambda < \omega_1)$, then for every rational $r \geq 0$, the sets

$$\{x : \mu(E^X) > r\} \text{ and } \{x : \mu(E^X) \geq r\}$$

are also in $\underline{BR}_\lambda^\rho$. \square

7. Approximating R-sets in Product Spaces : (In this section spaces of type 1 will be denoted by X, Y , etc. and E^X will denote the vertical section of E at x).

R-sets in $X \times Y$ are in general complicated sets and cannot be related to any reasonable product σ -field. For instance, as shown by B.V. Rao [41], the universal Σ_1^1 set U does not belong to $\mathcal{L} \otimes \mathcal{L}$, the product of σ -fields of Lebesgue measurable sets. However, V.V. Srivatsa [46] has shown that C-sets in $X \times Y$ can be approximated, in the sense of measure and category, by sets in product σ -fields. More precisely, if $A \subseteq X \times Y$ is a C-set, then there are B and C in $\mathcal{C}(X) \otimes \mathcal{B}(Y)$ ($\mathcal{C}(X)$ is the family of C-sets in X and $\mathcal{B}(Y)$ the Borel σ -field on Y), such that $B \subseteq A \subseteq C$ and $C^x - B^x$ is meager (null) for each $x \in X$. The problem of approximating sets in product spaces was

first studied by V.V. Srivatsa and can be viewed as an alternative and, perhaps, easier method for proving various selection theorems (of [46]). Srivatsa's method, however, does not seem to generalize to R-sets. In this section we shall extend Srivatsa's theorem to R-sets. We shall show that R-sets in product spaces can be approximated as above, in the sense of category, by sets in product σ -fields. (The precise statement is given in Theorem 7.10). This incidentally will give the selection theorem of Burgess (of [13]).

We shall prove our result through a series of lemmas.

7.1 Lemma : Let \mathcal{F} be a σ -field closed under operation \mathcal{A} and let $\mathcal{B}_{\omega^\omega}$ denote the Borel σ -field on ω^ω . Let $E \subseteq \omega^\omega \times \omega^\omega$. Suppose there is a family $\{A_n\}$ such that

$$(a) \quad A = \mathcal{A}(\{A_n\})$$

$$(b) \quad A \subseteq E$$

$$(c) \quad \text{Each } A_n \in \mathcal{F} \otimes \mathcal{B}_{\omega^\omega}$$

$$(d) \quad A^X \text{ is comeager whenever } E^X \text{ is comeager.}$$

Then there is a set $B \in \mathcal{F} \otimes \mathcal{B}_{\omega^\omega}$ such that $B \subseteq E$ and B^X is comeager whenever E^X is comeager.

Proof : Let $I : \omega^\omega \rightarrow \omega^\omega$ be the characteristic function of a generator for a countably generated sub- σ -field \mathcal{F}_0 of \mathcal{F} such that each $A_n \in \mathcal{F}_0 \otimes \mathcal{B}_{\omega^\omega}$. Let $I(\omega^\omega) = D$. Then, as is well-known, I is a bimeasurable function between \mathcal{F}_0 and \mathcal{B}_D , the Borel σ -field on D .

Set

$$A'_n = \{(I(\alpha), \beta) : (\alpha, \beta) \in A_n\},$$

$$A' = \{(I(\alpha), \beta) : (\alpha, \beta) \in A\}.$$

Clearly, $A'_n \in \mathcal{B}_D \otimes \mathcal{B}_{\omega^\omega}$ and $A' = \mathcal{A}(\{A'_n\})$.

Hence A' is an analytic set in $D \times \omega^\omega$. Get an analytic set $C \subseteq \omega^\omega \times \omega^\omega$ such that

$$A' = C \cap (D \times \omega^\omega).$$

Then by 1.6 of [46], get $B' \in \mathcal{A} \otimes \mathcal{B}_{\omega^\omega}$, where \mathcal{A} is the analytic σ -field on ω^ω , such that $B' \subseteq C$ and B'^X is comeager whenever C^X is comeager. Put

$$B = \{(\alpha, \beta) : (I(\alpha), \beta) \in B'\}.$$

Since \mathcal{F} is closed under operation \mathcal{A} , clearly $B \in \mathcal{F} \otimes \mathcal{B}_{\omega^\omega}$ and moreover, $B \subseteq E$ and B^X is comeager whenever E^X is comeager.

We have adapted the proof of the Category Formula [13] in the next lemma.

7.2 Lemma : Suppose we have

$$\begin{aligned} & (\forall s_0) (\exists t_0) (\forall s_1) (\exists t_1) \dots \left\{ (\forall a_0) (\exists b_0) (\forall a_1) (\exists b_1) \dots P(\alpha, \beta) \right\} \\ \longleftrightarrow & (\forall s_0) (\forall a_0) (\exists t_0) (\exists b_0) (\forall s_1) (\forall a_1) (\exists t_1) (\exists b_1) \dots P(\alpha, \beta), \end{aligned}$$

where $\alpha = (a_0, b_0, a_1, b_1, \dots)$; $a_i, b_i \in \omega$

and $\beta = s_0 * t_0 * s_1 * t_1 \dots$; $s_i, t_i \in \omega^{<\omega}$.

If \exists wins the second game with strategy σ , then he can modify σ to a

winning strategy σ^* such that, to every complete play $s_0, a_0, t_0, b_0, \dots$ consistent with σ^* , there corresponds a complete play $s'_0, t'_0, s'_1, t'_1, \dots$ consistent with a winning strategy for \exists in the first game such that

$$s_0 * t_0 * s_1 * t_1 \dots = s'_0 * t'_0 * s'_1 * t'_1 \dots$$

Proof. Let $w_{-1}, w_0, w_1, w_2, \dots$ be an enumeration of $\omega^{<\omega}$ such that w_{-1} is the empty sequence and if w_i is an initial segment of w_j , then $i < j$. For $s \in \omega^{<\omega}$, its code, denoted by $|s|$, is its position in the enumeration.

Suppose \exists wins the second game with strategy σ . We shall now construct a winning strategy τ for player II in the first (Banach-Mazur) game. τ will be defined (by induction) in such a way that every partial play $s_0, t_0, \dots, s_{n-1}, t_{n-1}$ consistent with τ corresponds to a partial play $a_0, u_0, b_0, v_0, \dots, a_{m-1}, u_{m-1}, b_{m-1}, v_{m-1}$ consistent with σ , such that $|a_0, \dots, a_{m-1}| = n-1$ and $s_0 * t_0 * \dots * t_{n-1} = u_0 * v_0 * \dots * u_{m-1} * v_{m-1}$.

Suppose $s_0, t_0, \dots, s_{n-1}, t_{n-1}$ have been defined consistent with τ . We now define $\tau(s_0, t_0, \dots, s_{n-1}, t_{n-1}, s)$. Let n be the code of (a_0, \dots, a_{m-1}, a) and let $|a_0, \dots, a_{m-1}| = n'$. Clearly, $n' < n$. Let the partial play (consistent with σ) corresponding to $(s_0, t_0, \dots, s_{n'}, t_{n'})$ be $(a_0, u_0, \dots, a_{m-1}, u_{m-1}, b_{m-1}, v_{m-1})$ such that $s_0 * t_0 * \dots * s_{n'} * t_{n'} = u_0 * v_0 * \dots * u_{m-1} * v_{m-1}$. Put $u = s_{n'+1} * t_{n'+1} * \dots * s_{n-1} * t_{n-1} * s$.

Let

$$\sigma(a_0, u_0, b_0, v_0, \dots, b_{m-1}, v_{m-1}, a, u) = (b, v).$$

Then

$$\tau(s_0, t_0, \dots, s_{n-1}, t_{n-1}, s) = v,$$

and the partial play associated with

$$s_0, t_0, \dots, s_{n-1}, t_{n-1}, s, v$$

is

$$a_0, u_0; b_0, v_0; \dots; b_{n-1}, v_{n-1}; a, u; b, v.$$

We shall now show that τ is a winning strategy for player II in the Banach-Mazur game. Let

$$s_0, t_0, s_1, t_1, s_2, t_2, \dots$$

be a complete play consistent with τ . We shall have to show that

$$(\forall a_0)(\exists b_0)(\forall a_1)(\exists b_1) \dots P(\alpha, \beta),$$

where $\alpha = (a_0, b_0, a_1, b_1, \dots)$ and

$$\beta = s_0 * t_0 * s_1 * t_1 * \dots$$

So let ψ play a_0 . Suppose $|a_0| = n_0$. By definition of τ , the partial play (consistent with σ) corresponding to $s_0, t_0, \dots, s_{n_0}, t_{n_0}$ is

$$a_0, s_0 * t_0 * \dots * s_{n_0}; b_0, t_{n_0}$$

for some $b_0 \in \omega$. \exists replies with b_0 . Next suppose ψ plays a_1 and let $|a_0, a_1| = n_1$. By definition of τ , the partial play (consistent with σ) corresponding to

$$s_0, t_0, \dots, s_{n_0}, t_{n_0}, \dots, s_{n_1}, t_{n_1}$$

is

$$a_0, \underbrace{s_0 * t_0 * \dots * s_{n_0}}_{u_0}; b_0, t_{n_0} = v_0; a_1, \underbrace{s_{n_0+1} * t_{n_0+1} * \dots * s_{n_1}}_{u_1}; b_1, t_{n_1} = v_1.$$

\exists replies with b_1 and the play proceeds as described. Since the play

$$a_0, u_0; b_0, v_0; a_1, u_1; b_1, v_1; \dots$$

is consistent with σ , we have

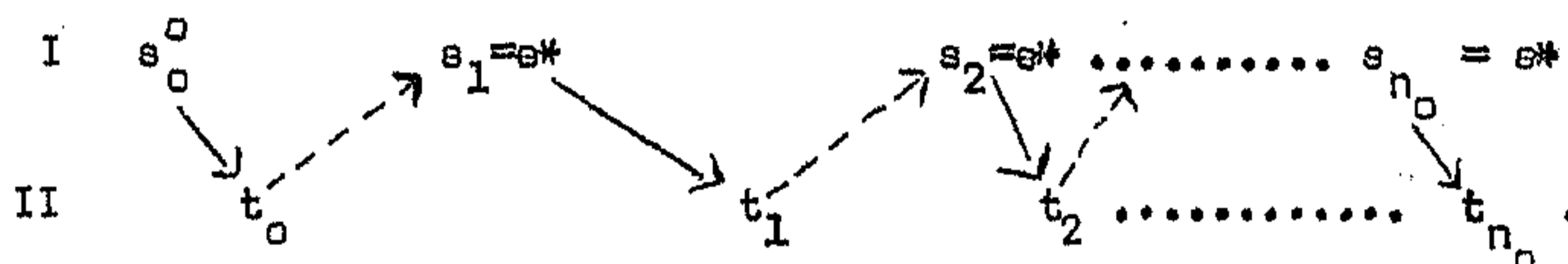
$$P(\alpha, u_0 * v_0 * u_1 * v_1 * \dots).$$

But then $\beta = u_0 * v_0 * u_1 * v_1 * \dots$

Hence $P(\alpha, \beta)$. Consequently, τ is a winning strategy for II in the Banach-Mazur game.

We shall now modify σ to σ^* such that any complete play consistent with σ^* corresponds to a complete play consistent with τ .

Definition of σ^* : Let ψ play a_0, s_0^0 and suppose $|a_0| = n_0$. Simulate the following partial play in the Banach-Mazur game, where II plays with strategy τ , σ^* being a fixed sequence:



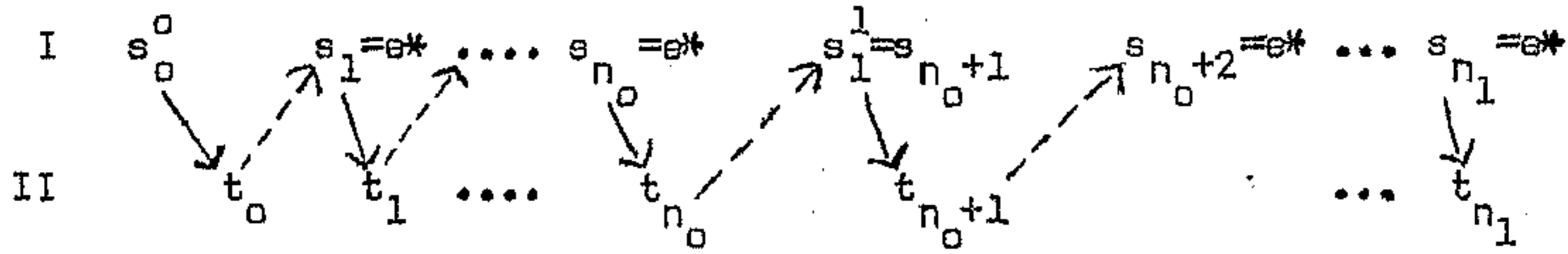
According to the definition of τ ,

$$\sigma(a_0, s_0^0 * t_0 * s_1 * t_1 * \dots * s_{n_0}) = (b_0, t_{n_0}),$$

for some $b_0 \in \omega$.

Define $\sigma^*(a_0, s_0^0) = (b_0, t_0 * s_1 * t_1 * \dots * s_{n_0} * t_{n_0})$.

Next suppose ψ plays a_1, s_1^1 and let $|a_0, a_1| = n_1$. Consider the following partial play consistent with τ



By definition of τ ,

$$\sigma(a_0, s_0^0 * t_0 * \dots * s_{n_0}^0; b_0, t_{n_0}; a_1, s_{n_0+1}^1 * t_{n_0+1} * \dots * s_{n_1}^1) = (b_1, t_{n_1}), b_1 \in \omega.$$

$$\begin{aligned} \text{Define } \sigma^*(a_0, s_0^0; b_0, t_0 * s_1^0 * \dots * s_{n_0}^0 * t_{n_0}; a_1, s_1^1) \\ = (b_1, t_{n_0+1} * s_{n_0+2}^1 * \dots * s_{n_1}^1 * t_{n_1}) \end{aligned}$$

and so on. We shall show that σ^* is a winning strategy for \exists in the second game. First observe that any complete play

$$a_0, u_0; b_0, v_0; a_1, u_1; b_1, v_1; \dots$$

consistent with σ^* corresponds to a complete play

$$s_0, t_0, s_1, t_1, \dots$$

consistent with τ such that

$$s_0 * t_0 * s_1 * t_1 * \dots = u_0 * v_0 * u_1 * v_1 * \dots = \beta, \text{ say.}$$

To prove that σ^* is a winning strategy, we have to show $P(\alpha, \beta)$,

where $\alpha = (a_0, b_0, a_1, b_1, \dots)$. Next observe that any play consistent with σ^* is of the form

$$(\forall) \text{ I } a_0, s_0^0 \qquad a_1, s_1^1 * s_{n_0+1}^1 \dots$$

$$(\exists) \text{ II } b_0, t_0 * s_1^0 * t_1 * \dots * s_{n_0}^0 * t_{n_0} \qquad b_1, t_{n_0+1} * \dots * s_{n_1}^1 * t_{n_1} \dots$$

such that

$$\begin{aligned}
 (\forall) \text{ I } & a_0, s_0^0 * t_0^0 * \dots * s_{n_0} & a_1, s_{n_0+1} * t_{n_0+1} * \dots * s_{n_1} & \dots \dots \dots \\
 (\exists) \text{ II } & & b_0, t_{n_0} & b_1, t_{n_1} \dots \dots
 \end{aligned}$$

is consistent with σ . Consequently we have

$$P(\alpha, s_0^0 * t_0^0 * \dots * s_{n_0} * t_{n_0} * s_{n_0+1} * t_{n_0+1} * \dots)$$

and hence $P(\alpha, \beta)$.

Remark. The definition of σ^* is highly constructive and one can show that the map $\sigma \rightarrow \sigma^*$ is Borel.

The proof of Lemma 7.2 immediately gives us the following

7.3 Corollary. If the second game in 7.2 is determined then

$$\begin{aligned}
 (a) & (\forall s_0)(\exists t_0)(\forall s_1)(\exists t_1)\dots \{ (\forall a_0)(\exists b_0)(\forall a_1)(\exists b_1)\dots P(\alpha, \beta) \} \\
 \iff & (\forall s_0)(\forall a_0)(\exists t_0)(\exists b_0)(\forall s_1)(\forall a_1)(\exists t_1)(\exists b_1)\dots P(\alpha, \beta) , \\
 (b) & (\exists s_0)(\forall t_0)(\exists s_1)(\forall t_1)\dots \{ (\exists a_0)(\forall b_0)(\exists a_1)(\forall b_1) \dots \neg P(\alpha, \beta) \} \\
 \iff & (\exists s_0)(\exists a_0)(\forall t_0)(\forall b_0)\dots \dots \dots \neg P(\alpha, \beta) ,
 \end{aligned}$$

and the conclusion of Lemma 7.2 also holds here.

Remark. (The proofs of) Lemma 7.2 and Corollary 7.3 yield a constructive proof of a particular case of the Game Formula of Kechris [26]. In fact, 7.2 and 7.3 can be formulated in the set up of Kechris.

Lemma 7.4 For $\rho < \omega_1$, let $R_\rho = \mathbb{I}_X = \mathbb{R} \mathbb{Q}_{NN}^\rho$. As in the proof of 6.6 there are functions f and g such that for any normal family $\{E_n\}$ with $E = R_\rho(\{E_n\})$,

E^X is comeager

$\longleftrightarrow (\exists \eta \in \overleftarrow{K}) (\forall n \in \eta) [E_{\overleftarrow{g}(n)}^X \text{ is comeager in } \Sigma(\overleftarrow{g}(n))]]$

$\longleftrightarrow (\exists \eta_0 \in M) (\forall n_0 \in \eta_0) \dots (\forall k) [E_{\langle \alpha(\langle n, m \rangle), \dots, \alpha(\langle n_{k-1}, m_{k-1} \rangle) \rangle}^X$

is comeager in $\Sigma(\beta(\langle n_0, m_0 \rangle) * \dots * \beta(\langle n_{k-1}, m_{k-1} \rangle))]$.

where $\overleftarrow{g}, \overleftarrow{g}$ and α, β are related as in 6.3(c) and $\overline{Q}_M \hookrightarrow \overline{Q}_N$ and $\overleftarrow{K} = \mathbb{R} M M^0$.

Then, one can choose \overleftarrow{K} such that for any $\eta \in \overleftarrow{K}$, $\bigcup_{n \in \eta} \Sigma(\overleftarrow{g}(n))$ is dense in ω^ω .

In fact, \overleftarrow{K} may be taken to be the canonical base for S_ρ with the property that for any $\eta \in \overleftarrow{K}$ and $s \in \text{Seq}$ there is an $n \in \eta$ such that $\overleftarrow{g}(n)$ extends s .

Proof. We shall prove this by induction on ρ . For $\rho = 0$, $\overline{Q}_{\overleftarrow{K}}$ is the Vaught operation \mathcal{V} . Taking \overleftarrow{K} to be the canonical base for \mathcal{V} it is easy to check the assertion of the above theorem.

So assume $\rho > 0$ and the assertion holds for all $\lambda < \rho$.

Case 1. $\rho = \lambda + 1$.

Then $\overline{Q}_N = R_\lambda$. Let $\eta \in \overleftarrow{K} = \mathbb{R} M M^0$, where M has been chosen to satisfy the assertion of the theorem for the operation R_λ .

Fix a basic clopen set $\Sigma(t)$. We have to show that there is $n \in \eta$ such that $\overleftarrow{g}(n)$ is consistent with t (in fact, extends t).

Since $\eta \in \mathbb{R} M M^0$,

$(\exists \eta_0 \in M)(\forall n_0 \in \eta_0)(\forall \xi_0 \in M)(\exists n_0 \in \xi_0) \dots (\forall k) [\langle \langle n_0, m_0 \rangle \dots \langle n_{k-1}, m_{k-1} \rangle \rangle \in \eta]$.

Fix $\eta_0 \in M$ such that for each $n_0 \in \eta_0$,

$(\forall \xi_0 \in M)(\exists m_0 \in \xi_0)(\exists \eta_1 \in M)(\forall n_1 \in \eta_1)(\forall \xi_1 \in M)(\exists m_1 \in \xi_1) \dots$

$\dots (\forall k) [\langle \langle n_0, m_0 \rangle, \dots, \langle n_{k-1}, m_{k-1} \rangle \rangle \in \eta]$.

By induction hypothesis, there is $n'_0 \in \eta_0$ such that $g(n'_0)$ extends t .

Now pick a $\xi_0 \in M$ such that for some $m'_0 \in \xi_0$,

$(\exists \eta_1 \in M)(\forall n_1 \in \eta_1)(\forall \xi_1 \in M)(\exists m_1 \in \xi_1) \dots$

$\dots (\forall k) [\langle \langle n'_0, m'_0 \rangle, \langle n_1, m_1 \rangle, \dots, \langle n_{k-1}, m_{k-1} \rangle \rangle \in \eta]$.

Clearly $n = \langle \langle n'_0, m'_0 \rangle \rangle \in \eta$ and $\tilde{g}(n) = g(n'_0) * \delta(n'_0)$ (cf. 6.3).

Thus $\tilde{g}(n)$ extends t .

Case 2. ρ is limit.

Let $\rho_i \uparrow \rho$. Then $R_\rho = \tilde{\mathcal{G}}_N^* = \mathbb{R} \tilde{\mathcal{G}}_{NN^0}$, where $\tilde{\mathcal{G}}_N$ is the operation given by

$$\tilde{\mathcal{G}}_N(\{F_n\}) = \bigcap_{i \in \omega} \tilde{\mathcal{G}}_{N\rho_i N^0}(\{F_{\langle i, m \rangle} : m \in \omega\})$$

and $\tilde{\mathcal{G}}_{N\rho_i} = R_{\rho_i}$.

In this case, if $E = \tilde{\mathcal{G}}_N(\{E_n\})$, then E^X is comeager

$\iff (\exists \eta \in \tilde{\mathcal{N}})(\forall n \in \eta) [E^X_{\tilde{f}(n)}$ is comeager in $\Sigma(\tilde{g}(n))]$

$\iff (\forall i)(\exists \eta \in M_{\rho_i})(\forall n \in \eta)(\exists \xi \in M_{\rho_i}^0)(\forall m \in \xi) [E^X_{\langle i, \alpha_i \langle n, m \rangle \rangle}$ is comeager in $\Sigma(\beta_i(\langle n, m \rangle))]$ (cf. 6.6, Case 2).

In view of Case 1, it suffices to show that for any $\eta \in \tilde{N}$ and for any $\Sigma(t)$, there is $n \in \eta$ such that $\tilde{g}(n)$ extends t .

Now, \tilde{N} may be chosen such that

$$\eta \in \tilde{N} \iff (\forall i) (\exists \eta' \in M_{\rho_i}^0) (\forall n \in \eta') (\exists \xi' \in M_{\rho_i}^0) (\forall m \in \xi') [\langle i, \langle n, m \rangle \rangle \in \eta],$$

where each $M_{\rho_i}^0$ has been chosen to satisfy the assertion of the theorem.

Let $\eta \in \tilde{N}$ and fix any $i' \in \omega$. Get $\eta' \in M_{\rho_{i'}}^0$ such that for each $n \in \eta'$, $(\exists \xi' \in M_{\rho_{i'}}^0) (\forall m \in \xi') [\langle i', \langle n, m \rangle \rangle \in \eta]$.

Since $\eta' \in M_{\rho_{i'}}^0$, by induction hypothesis, there is $n' \in \eta'$ such that $g_{i'}(n')$ extends t . Get $\xi' \in M_{\rho_{i'}}^0$ and $m' \in \xi'$ such that $\langle i', \langle n', m' \rangle \rangle \in \eta$. Take $n = \langle i', \langle n', m' \rangle \rangle$. Then

$$\tilde{g}(n) = \beta_{i'}(\langle n', m' \rangle) = g_{i'}(n') * \delta_{i'}(m') \quad [\text{cf. 6.6}]$$

and hence $\tilde{g}(n)$ extends t .

Lemma 7.5: Suppose $E \subseteq \omega^\omega$ be a set in \mathcal{R}^ρ , $\rho < \omega_1$ and let

$R_\rho = \mathcal{R}_{\mathbb{I}_N}$. Assume $E = R_\rho(\{E_n\})$, where each $\{E_n\}$ is clopen. Then,

E is comeager $\iff (\exists \eta_0 \in M) (\forall n_0 \in \eta_0) (\exists \eta_1 \in M) (\forall n_1 \in \eta_1) \dots$

$\dots (\forall k) [E_{\langle f(n_0), \dots, f(n_{k-1}) \rangle}$ is comeager in $\Sigma(g(n_0) * \dots * g(n_{k-1}))]$

where f and g are some suitable functions and $\mathbb{I}_\eta \leftrightarrow \mathbb{I}_N$. Moreover,

f and g can be chosen such that, with any winning strategy σ for

\exists one can associate a winning strategy σ^* such that for any run

$\eta_0, n_0, \eta_1, n_1, \dots$ consistent with σ^* , the sequence $\beta = g(n_0) * g(n_1) * \dots$

is in E .

Proof: The first assertion follows immediately from the proof of 6.6.

Moreover, by 7.4, M may be chosen such that for any $\eta \in M$ and $s \in \text{Seq}$ there is $n \in \eta$ such that $g(n)$ extends s . Therefore, we have

$$(\exists \eta_0 \in M)(\forall n_0 \in \eta_0) \dots (\forall k) [E_{\langle f(n_0), \dots, f(n_{k-1}) \rangle} \text{ is comeager in } \Sigma(g(n_0) * g(n_1) * \dots * g(n_{k-1}))]$$

$$\longleftrightarrow (\exists \eta_0 \in M)(\forall n_0 \in \eta_0) \dots (\forall k) [\beta \in E_{\langle f(n_0), \dots, f(n_{k-1}) \rangle}]$$

where $\beta = g(n_0) * g(n_1) * \dots$

Hence,

E is comeager

$$\longleftrightarrow (\exists \eta_0 \in M)(\forall n_0 \in \eta_0) (\exists \eta_1 \in M)(\forall n_1 \in \eta_1) \dots \dots (\forall k) [g(n_0) * g(n_1) * \dots \in E_{\langle f(n_0), \dots, f(n_{k-1}) \rangle}]. \quad (1)$$

We shall prove the second assertion only for $\rho = 1$, since for higher levels the proof involves similar ideas, although notationally cumbersome. (For a complete proof see Appendix).

First observe that if $F = \mathcal{A}(\{F_n\})$, with each F_n clopen, then by the Category Formula [13], we have

F is meager

$$\longleftrightarrow (\forall k_0) (\forall s_0) (\exists t_0) (\forall k_1) (\forall s_1) (\exists t_1) \dots (\exists i) [F_{\langle k_0, \dots, k_{i-1} \rangle} \text{ is meager in } \Sigma(s_0 * t_0 * \dots * s_{i-1} * t_{i-1})].$$

Therefore, whenever $G = \mathcal{I}(\{G_n\})$, where $\mathcal{I} = \mathcal{A}^0$,

G is comeager

$$\begin{aligned}
 \longleftrightarrow & (w_{k_0^0})(w_{s_0^0})(\exists t_0^0)(w_{k_1^0})(w_{s_1^0})(\exists t_1^0) \dots (\exists i_0) \\
 & (w_{k_0^1})(w_{s_0^1})(\exists t_0^1)(w_{k_1^1})(w_{s_1^1})(\exists t_1^1) \dots (\exists i_1) \\
 & \dots \\
 & \dots (w_j) \left[s_0^0 * t_0^0 \dots * s_{i_0-1}^0 * t_{i_0-1}^0 * s_0^1 * t_0^1 \dots * s_{i_1-1}^1 * t_{i_1-1}^1 \dots \right. \\
 & \left. \in E \langle \langle k_0^0, \dots, k_{i_0-1}^0 \rangle, \dots, \langle k_0^{j-1}, \dots, k_{i_{j-1}-1}^{j-1} \rangle \rangle \right].
 \end{aligned} \tag{iii}$$

As in [13; §11], we shall define a game G' of length ω such that \exists wins G' iff \exists wins the second game in (iii). The game G' is as follows. Though its total length is ω we think of it as consisting of potentially infinite sequence of subgames each consisting of potentially infinite sequence of rounds. If in any play of G' the j -th subgame actually goes through infinitely many rounds, then the $(j+1)$ th subgame never gets started. This keeps the total length within bounds.

In the j -th subgame the two players play the rounds of that subgame. The λ -th such round opens with \exists signalling (by a choice of 0 or 1) either a challenge or a pass. If he challenges, the whole j -th subgame ends at once and the players proceed to the $(j+1)$ th subgame; in this case we record $u_j = \langle k_0^j, \dots, k_{\lambda-1}^j \rangle$ and $v_j = s_0^j * t_0^j * \dots * s_{\lambda-1}^j * t_{\lambda-1}^j$ formed from the moves. If \exists passes, \forall chooses $k_{\lambda}^j \in \omega$ and a sequence number s_{λ}^j and \exists replies with $t_{\lambda}^j \in \text{Seq}$; then the players proceed to the $(\lambda+1)$ -th round.

If some j -th subgame goes on forever because \exists fails to challenge on any round, \exists forfeits the game. If this provision does not apply, then a sequence $(u_0, v_0; u_1, v_1; \dots)$ will have been generated. \exists wins iff for all j

$$v_0 * v_1 * v_2 * \dots = \beta \in E_{\langle u_0, u_1, \dots, u_{j-1} \rangle}$$

Thus the game G' is essentially of the form

$$(\forall a_0) (\forall a_0) (\exists t_0) (\exists b_0) (\forall a_1) (\forall a_1) (\exists t_1) (\exists b_1) \dots [M(\alpha, \beta)],$$

where the condition M is Borel (in fact G_δ , cf. [13; §11]). We shall now make a few observations. First observe that each subgame in the second game of (iii) is a Banach-Mazur game interlaced with operation \mathcal{S} (which may be looked upon as a game!). Secondly the condition within [] is the same as the condition in the first game. Consequently, any reduction of the second game of (iii) involves the interlacing of a Banach-Mazur game with an ordinary game of length ω which is obtained by the corresponding reduction of the operation R_1 . Therefore, reducing the R_1 operation within { } in the first game of (iii) as above, we observe that the equivalence (iii) is equivalent to

$$\begin{aligned} & (\forall a_0) (\exists t_0) (\forall a_1) (\exists t_1) \dots \left\{ (\forall a_0) (\exists b_0) \dots M(\alpha, \beta) \right\} \\ & \longleftrightarrow (\forall a_0) (\forall a_0) (\exists t_0) (\exists b_0) \dots M(\alpha, \beta). \end{aligned} \tag{iv}$$

Now, if \exists wins the second game in (iii), he also wins G' , the second game in (iv). By 7.2 \exists wins with a strategy σ such that for any

complete play $s_0, a_0, t_0, b_0, \dots$ consistent with σ , $\beta = s_0 * t_0 * \dots$ corresponds to a complete play of the Banach-Mazur game in (iv).

Moreover, the strategy σ gives rise to a strategy σ^* in the second game of (iii) such that, if $s_0^0, t_0^0, k_0^0, \dots, i_0^0; s_0^1, t_0^1, k_0^1, \dots, i_1^1, \dots$ is a complete run consistent with σ^* , then the sequence

$\beta = s_0^0 * t_0^0 * \dots * s_{i_0^0-1}^0 * t_{i_0^0-1}^0 * s_{i_0^0}^1 * t_{i_0^0}^1 * \dots$ is also produced by a complete run of the game G' , when \exists plays with σ . Consequently, β satisfies the condition in $\{ \}$ in (iv) and hence $\beta \in E$.

Remark: For sets E at the higher level we have an equivalence similarly to (iii). On the left side of the equivalence we have a Banach-Mazur game followed by the corresponding R -operation, and on the right side we have Banach-Mazur games "interlaced" with the operation (regarded as a game played on ω) with dummy moves, if necessary. As in the proof above, the second game can be reduced to a game of length ω and the proof proceeds exactly as above (see Appendix).

7.6 On the Δ -transform: We shall now look at the dual of the $*$ -transform and obtain some analogous results. The proofs being very similar to what we have already done, we shall omit them.

For a set $E \subseteq X \times Y$ and $U \subseteq Y$, put $E^{\Delta U} = \{x \in X : E^x \text{ is non-meager in } U\}$.

Note that Δ is the dual of the $*$ -transform in the sense that $E^{\Delta U} = [(E^c)^* U]^c$. The Δ -transform behaves very much like the $*$ -transform and we have the counterpart of 6.2 :

With notation as in 6.2, we have

$$E^{\Delta} \Sigma(s) = \bigcap_{\mu < \omega_1} [E_{\sigma}^{\mu}]^{\Delta} \Sigma(s) = \bigcup_{\mu < \omega_1} [E_{\sigma}^{\mu} - T^{\mu}]^{\Delta} \Sigma(s).$$

This decomposition of the set E^{Δ} suggests, as in the case of E^* , that $(E^{\Delta})^{\square}$ may be obtained as a fixed point of an inductive operator. Indeed we have the following counterpart of 6.3.

Theorem: Let \bar{Q}_N and \bar{Q}_M be two δ -e operations such that \bar{Q}_N preserves the Baire property and \bar{Q}_M is normal and subsumes both (countable) \cup and \cap . Suppose there are functions f_1 and g_1 such that for any normal family $\{E_p\}$ of subsets of $\omega^{\omega} \times \omega^{\omega}$ with $E = \bar{Q}_N(\{E_p\})$,

$$E^{\Delta} \Sigma(s) = \bar{Q}_M(\{E_{f_1(p)}^{\Sigma(s * g_1(p))} : p \in \omega\}).$$

Then for any regular normal family $\{F_p\}$, if $F = \bar{Q}_N(\{F_p\})$ then

F^X is non-meager in $\Sigma(s)$

$\iff (\exists \eta_0 \in M^1)(\forall n_0 \in \eta_0)(\exists \eta_1 \in M^1)(\forall n_1 \in \eta_1) \dots$

$\dots (\forall k) [F_{\langle f^1(n_0), \dots, f^1(n_{k-1}) \rangle}^X$ is non-meager in $\Sigma(s * g^1(n_0) * \dots * g^1(n_{k-1}))]$,

where f^1 and g^1 are suitable functions (independent of $\{F_p\}$) and

$$\bar{Q}_M = \bar{Q}_N \cup \cap.$$

Proof: First observe that for any regular family $\{F_p\}$, $F_s^{\mu} \subseteq F_t$ if $t \subseteq s$. As in 6.3(c) we define a set relation operative on ω as follows.

$t \in \Gamma_x(A) \iff F_x^x(t)$ is meager in $\Sigma((t)_1) \vee$

$(\forall u \in M)(\exists n \in \mathbb{N})(\forall v) [\langle (t)_0 * \langle f_1(n) \rangle, (t)_1 * g_1(n) * u * v \rangle \in A]$.

To obtain the result, one then shows that $x \in F^{\Delta \Sigma(e)} \iff \langle e, s \rangle \notin \Gamma_x^\infty$. \square

As in 6.6, this theorem immediately implies that for a set $E \in \mathcal{R}^0$, if

$E = \mathbb{R} \overline{\bigcap}_N \{E_n\}$, then

$$x \in E^\Delta \iff (\exists \eta_0 \in M) (\forall n_0 \in \eta_0) (\exists \eta_1 \in M) (\forall n_1 \in \eta_1) \dots \dots (\forall k) [E_{\langle f^i(n_0), \dots, f^i(n_{k-1}) \rangle}^x \text{ is non-meager in } \Sigma(g^i(n_0) * \dots * g^i(n_{k-1}))] \quad (9)$$

for suitable M, f^i, g^i . Moreover, one can choose M such that

$$\eta \in M^0 \& s \in \text{Seq} \implies (\exists n \in \eta) (g^i(n) \text{ extends } s).$$

Finally observe that by (9) we have for $E \subseteq \omega^\omega$

$$E \text{ is meager} \iff (\exists \eta_0 \in M^0) (\forall n_0 \in \eta_0) (\exists \eta_1 \in M^0) (\forall n_1 \in \eta_1) \dots \dots (\exists k) [g^i(n_0) * g^i(n_1) * \dots \notin E_{\langle f^i(n_0), \dots, f^i(n_{k-1}) \rangle}] \quad (v)$$

Arguing as in 7.5 and invoking 7.3(b), we may show that \exists can win the game (v) (if he does so) by a strategy σ^{**} such that if $\eta_0, n_0, \eta_1, n_1, \dots$ is a complete play consistent with σ^{**} , then $\beta = g(n_0) * g(n_1) * \dots$ is not in E .

The next theorem is the first step towards our approximation theorem.

7.7 Theorem: Let $E \subseteq \omega^\omega \times \omega^\omega$ be a set in \mathcal{R}^0 ($\rho > 0$). Then there is a set $B \in \mathcal{BR}_0^\rho \otimes \mathcal{B}_\omega$ such that $B \subseteq E$ and B^x is comeager on E^* .

Proof: Let $R_0 = \mathbb{R}_{\mathbb{Q}_N}$ and assume $E = \mathbb{R}_{\mathbb{Q}_N}(\{E_n\})$, with each E_n clopen. As in 6.6, there are functions f and g such that

$$x \in E^* \iff (\exists \eta_0 \in M)(\forall n_0 \in \eta_0)(\exists \eta_1 \in M)(\forall n_1 \in \eta_1) \dots \\ \dots (\forall k) [E_{\langle f(n_0), \dots, f(n_{k-1}) \rangle}^x \text{ is comeager in } \Sigma(g(n_0)*\dots*g(n_{k-1}))] \quad (i)$$

where M is such that $\mathbb{Q}_M \hookrightarrow \mathbb{Q}_N$. Further, by 3.8, there is a $\mathbb{B}R_0^0$ -measurable function $x \mapsto \tau_x$ such that whenever $x \in E^*$, τ_x is a winning strategy for the player \exists in the game (i). Moreover, by 7.5 and the Remark following 7.2, we may assume that the function $x \mapsto \tau_x$ is such that for any complete play $\eta_0, n_0, \eta_1, n_1, \dots$ agreeing with τ_x the sequence $\beta = g(n_0)*g(n_1)*g(n_2)*\dots$ is in E^x , and that the base M has been chosen as in 7.4.

Define for each $k \geq 1$,

$$W_k(s, x) \iff \text{Seq}(s) \ \& \ \text{lh}(s) = k \ \& \ (\forall i < \text{lh}(s)) [(s)_i \in \tau_x(s \upharpoonright i)]$$

i.e., W_k^x consists of the codes of the first k possible moves of \forall when the existential player plays according to the strategy τ_x . Clearly,

$W_k \in \mathbb{B}R_0^0$. Define,

$$C(x, y) \iff (\exists \alpha)(\forall k) [W_k(\bar{\alpha}(k), x) \ \& \ y \in \Sigma(g(\alpha(0))*\dots*g(\alpha(k-1)))] \ \& \ x \in E^*.$$

Plainly, C is the result of operation A on sets in $\mathbb{B}R_0^0 \otimes \mathbb{B}_{\omega}^0$. We shall show that

$$(a) \ C \subseteq E$$

and (b) C^x is comeager whenever E^x is comeager.

If $(x, y) \in C$ then $x \in E^*$ and \exists wins the game (i). Further, there

is a sequence $\{n_k : k \in \omega\}$ such that n_0, n_1, n_2, \dots are the moves of \forall when \exists plays according to τ_x and $y = g(n_0) * g(n_1) * \dots$. Consequently, $y \in E^X$. To prove (b), fix an x such that E^X is comeager. Then

$$C^X = \left\{ y : (\exists \alpha)(\forall k) \left[\omega_k(\bar{\alpha}(k), x) \ \& \ y \in \Sigma(g(\alpha(0)) * \dots * g(\alpha(k-1))) \right] \right\}$$

Define,

$$A_{\langle n_0, \dots, n_{k-1} \rangle} = \Sigma(g(n_0) * g(n_1) * \dots * g(n_{k-1})), \text{ if } \omega_k(\langle n_0, \dots, n_{k-1} \rangle, x) \\ = \varnothing, \text{ otherwise.}$$

Then, $C^X = \bigcap \{A_s\}$. Hence to show C^X is comeager, it is enough to show (by Vaught's formula) that

$$(\forall s_0)(\exists k_0)(\exists t_0)(\forall s_1)(\exists k_1)(\exists t_1) \dots (\forall i) \left[A_{\langle k_0, \dots, k_{i-1} \rangle} \text{ is} \right. \\ \left. \text{comeager in } \Sigma(s_0 * t_0 * \dots * s_{i-1} * t_{i-1}) \right].$$

So let \forall play s_0 . By 7.4, pick $k_0 \in \tau_x(\langle \rangle)$ such that $g(k_0)$ extends s_0 . Then \exists replies with k_0 and a sequence number t_0 such that $s_0 * t_0 = g(k_0)$. Next when \forall plays s_1 , \exists plays $k_1 \in \tau_x(\langle k_0 \rangle)$ such that $g(k_1)$ extends s_1 and then plays t_1 such that $s_1 * t_1 = g(k_1)$, and so on. By adopting this strategy, \exists ensures that for each i ,

$$A_{\langle k_0, \dots, k_{i-1} \rangle} \text{ is comeager in } \Sigma(s_0 * t_0 * \dots * s_{i-1} * t_{i-1})$$

Thus, C^X is comeager. An application of 7.1 gives the required set B .

The next lemma is the counterpart of 7.7.

Lemma 7.8: Let $E \subseteq \omega^\omega \times \omega^\omega$ be a set in \mathcal{R}^p .

Let $E^\# = \{x : E^X \text{ is meager}\}$. Then $E^\# \in \mathcal{C}\mathcal{R}^p$ and there is a set

$C \in \mathcal{B}(\mathcal{R}^D) \otimes \mathcal{B}_{\omega\omega}$ (in fact, in $\sigma(\mathcal{R}^D) \otimes \mathcal{B}_{\omega\omega}$) such that $E \subseteq C$ and C^X is meager on $E^\#$.

Proof: Plainly, $E^\# \in \sigma(\mathcal{R}^D)$, by 6.7.

Let f, g, N, M, E_n be as in 7.6 with $\{E_n\}$ regular. Then, by (v) of 7.6,

E^X is meager

$$\begin{aligned} \longleftrightarrow (\exists n_0 \in M^0) (\forall n_0 \in \eta_0) (\exists n_1 \in M^0) (\forall n_1 \in \eta_1) \dots \\ \dots (\exists k) [g(n_0) * g(n_1) * \dots * \notin E^X_{\langle f(n_0), \dots, f(n_{k-1}) \rangle}] \end{aligned} \tag{i}$$

Further, by the Remark following Theorem 3.8 there is a $\sigma(\mathcal{R}^D)$ -measurable function $x \rightarrow \tau_x$ such that τ_x is a winning strategy for \exists in the above game whenever $x \in E$. Moreover, τ_x can be chosen as in 7.6. Call $u = \langle n_0, \dots, n_{k-1} \rangle$ good with respect to τ_x and β if β extends $g(n_0) * \dots * g(n_{k-1})$ and n_0, \dots, n_{k-1} are the first k possible moves of \forall when \exists plays with strategy τ_x . Define

$$\begin{aligned} T(u, x, \beta) \longleftrightarrow \text{seq}(u) \ \& \ (\forall i < h(u)) [(u)_i \in \tau_x(u|_i)] \\ \& \ \beta \in \Sigma((u)_0 * g((u)_1) * \dots * g((u)_{h(u)-1})) \ \& \\ (\forall n \in \tau_x(u)) [\beta \notin \Sigma((u)_0 * g((u)_1) * g((u)_2) * \dots * g((u)_{h(u)-1}) * g(n))] \end{aligned}$$

In other words, for each x, β , the section $T_{x, \beta}$ consists of all maximal good sequences with respect to τ_x and β . Clearly, $T \in \sigma(\mathcal{R}^D) \otimes \mathcal{B}_{\omega\omega}$ and moreover $T^{u, x}$ is closed nowhere-dense for each good sequence u and x . Now define,

$$C'(x, \beta) \longleftrightarrow x \in E^\# \ \& \ (\exists u) T(u, x, \beta)$$

Set

$$C = C' \cup (\omega^\omega - E^\#) \times \omega^\omega.$$

Since $E^\# \in \sigma(\mathcal{R}^p)$, $C \in \sigma(\mathcal{R}^p) \otimes \mathcal{B}$; and to conclude the proof we shall show that $E \subseteq C$. So let $(x, \beta) \in E$. If $x \notin E^\#$, then we are done. So assume $x \in E^\#$. Therefore, \exists wins the game (i) with strategy τ_x . Now, suppose when \exists plays the game (i) according to τ_x , \forall is able to play n_0, n_1, n_2, \dots such that for each k , $\langle n_0, \dots, n_{k-1} \rangle$ is a good sequence with respect to τ_x and β . Then

$$\beta = g(n_0) * g(n_1) * \dots$$

and hence, by 7.6, $\beta \notin E^X$. Consequently, since we have $\beta \in E^X$, \forall is able to play n_0, n_1, \dots 'consistent' with β only upto a finite stage. Let u be the code of the maximal sequence. Then $T(\alpha, x, \beta)$ and hence $(x, \beta) \in C$.

The following gives the approximation at the first level.

7.9 Lemma. Let $E \subseteq \omega^\omega \times \omega^\omega$ be a set in \mathcal{R}^p ($p > 0$). Then there are sets B, C in $\mathcal{B}\mathcal{R}_0^p \otimes \mathcal{B}_{\omega^\omega}$ such that $B \subseteq E \subseteq C$ and $C^X - B^X$ is meager for each x .

Proof. Let $T_s = E^{*\Sigma(s)}$; $s \in \text{Seq}$.

As E^X satisfies the Baire property for each x ,

$$\bigcup_s T_s = \{x : E^X \text{ is non-meager}\}.$$

Since each $\Sigma(s)$ is a (recursive) homeomorphic copy of ω^ω , we can apply 7.7 to get B_s as in the Lemma. Put

$$B = \bigcup_s (B_s \cap (\tau_s \times \omega^\omega)).$$

Clearly, $B \subseteq E$, $B \in \mathcal{BR}_0^\rho \otimes \mathcal{B}_{\omega^\omega}$ and for each x , $E^x - B^x$ is meager.

To get the set C , work with E^c and argue as above using 7.8.

7.10 Approximation Theorem: Let A be a \mathcal{R}_μ^ρ subset of $\omega^\omega \times \omega^\omega$ ($\rho > 0$, $\rho, \mu < \omega_1$). Then there are sets B and C in $\mathcal{BR}_\mu^\rho \otimes \mathcal{B}_{\omega^\omega}$ such that $B \subseteq A \subseteq C$ and $C^x - B^x$ is meager for each x .

Moreover, if $A \in \mathcal{BR}_\mu^\rho$, then one can find B, C in $\mathcal{BR}_\mu^\rho \otimes \mathcal{B}_{\omega^\omega}$ with the above properties.

Proof: We shall prove this by induction on μ . When $\mu = 0$, this is just Lemma 7.9. So suppose $\mu > 0$ and that the result is true for all $\lambda < \mu$. Let $A \in \mathcal{R}_\mu^\rho$ and let $R_\rho = \mathcal{R}_{\mathbb{N}}^\rho$. Then A is the result of operation R_ρ on a family $\{A_n\}$ of sets from $\mathcal{C}[\bigcup_{\lambda < \mu} \mathcal{R}_\lambda^\rho]$. By induction hypothesis, for each n we have B_n and C_n both in $\bigcup_{\lambda < \mu} \mathcal{BR}_\lambda^\rho \otimes \mathcal{B}_{\omega^\omega}$ such that $B_n \subseteq A_n \subseteq C_n$ and $C_n^x - B_n^x$ is meager for each x . Let

$$B' = R_\rho(\{B_n\}) \text{ and } C' = R_\rho(\{C_n\}).$$

Then, $B' \subseteq C'$ and for each x , $C'^x - B'^x \subseteq \bigcup_n (C_n^x - B_n^x)$, and hence meager. To complete the proof it is enough to get $B \subseteq B'$ and $C \supseteq C'$ such that B and C are in $\mathcal{BR}_\mu^\rho \otimes \mathcal{B}_{\omega^\omega}$, $B'^x - B^x$ is meager and $C^x - C'^x$ is meager for each x . We will show how to obtain B ; C can be obtained similarly.

Let \mathcal{F} be the σ -field generated by $\bigcup_{\lambda < \mu} \mathcal{BR}_\lambda^\rho$. Then

$B' = R_\rho(\{B_n\})$ with each B_n in $\mathcal{F} \otimes \mathcal{B}_{\omega^\omega}$. Thus, one can obtain a countably generated sub- σ -field \mathcal{G} of \mathcal{F} such that each $B_n \in \mathcal{G} \otimes \mathcal{B}_{\omega^\omega}$. Fix a countable generator of \mathcal{G} and let $f : (\omega^\omega, \mathcal{G}) \rightarrow \omega^\omega$ be its characteristic function. Put $M = f(\omega^\omega)$. Then $(\omega^\omega, \mathcal{G})$ and (M, \mathcal{B}_M) are Borel isomorphic. For each n , let $B'_n = \{(f(x), y) : (x, y) \in B_n\}$. Then $B'_n \in \mathcal{B}_M \times \omega^\omega$, for each n . Let $B'' = \{(f(x), y) : (x, y) \in B'\}$. Then $B'' = R_\rho(\{B'_n\})$. Hence, there is a set $E \subseteq \omega^\omega \times \omega^\omega$ in \mathcal{R}^p such that $E \cap (M \times \omega^\omega) = B''$. Apply 7.9 and get $E' \subseteq E$ such that $E' \in \mathcal{BR}_0^p \otimes \mathcal{B}_{\omega^\omega}$ and $E'^x - (E')^x$ is meager for each x . Let $\tilde{f} : \omega^\omega \times \omega^\omega \rightarrow \omega^\omega \times \omega^\omega$ be the map

$$\tilde{f}(x, y) = (f(x), y).$$

Put $B = (\tilde{f})^{-1}(E')$. Note that since \tilde{f} is a bimeasurable map of $(\omega^\omega, \mathcal{G})$ and (M, \mathcal{B}_M) ,

$$f^{-1}(\mathcal{R}_\mu^p \upharpoonright \omega^\omega) \subseteq \{A \subseteq \omega^\omega : A \text{ is the result of operation } R_\rho \text{ on sets in } \mathcal{H} \subseteq \mathcal{R}_\mu^p\}.$$

Consequently, $f^{-1}(\mathcal{BR}_\mu^p \upharpoonright \omega^\omega) \subseteq \mathcal{BR}_\mu^p$. Thus $B \in \mathcal{BR}_\mu^p \otimes \mathcal{B}_{\omega^\omega}$ and clearly, $B'^x - B^x$ is meager for each x .

The second assertion follows from the first by observing that the class of all sets for which the result holds is closed under $\mathbb{I}_{\mathbb{N}}$ and complementation.

The next proposition is quite well-known and follows from the Von-Neumann selection theorem.

Proposition: Let (T, \mathcal{M}) be a measurable space, \mathcal{M} being a σ -field closed under operation \wedge , and let Y be a polish space. Let $B \in \mathcal{M} \otimes \mathcal{B}_Y$ have non-empty vertical sections. Then B has a \mathcal{M} -measurable selection.

As a consequence of the approximation theorem we have the following.

7.11 Theorem: Suppose $A \subseteq \omega^\omega \times \omega^\omega$ is a \mathcal{BR}_μ^ρ set ($\rho > 0$) such that A^x is non-meager for each x . Then A has a \mathcal{BR}_μ^ρ -measurable selection.

Proof: By 7.10, get $B \in \mathcal{BR}_\mu^\rho \otimes \mathcal{B}_{\omega^\omega}$ such that $B \subseteq A$ and $A^x - B^x$ is meager for each x . B^x is then non-empty for each x , and the above proposition yields the result.

The next selection theorem is due to Burgess [13].

7.12 Theorem: Let $F: \omega^\omega \rightarrow \omega^\omega$ be a multifunction such that $F(x)$ is non-meager in its closure $c\ell(F(x))$. If F is \mathcal{BR}^ρ -measurable and its graph $\text{Gr}(F) \in \mathcal{BR}^\rho$, then F has a \mathcal{BR}^ρ -measurable selection.

Proof: Define G by $G(x) = c\ell(F(x))$. Then G is a closed-valued, \mathcal{BR}^ρ -measurable multifunction. Hence there is a map $g: \omega^\omega \times \omega^\omega \rightarrow \omega^\omega$ such that g is $\mathcal{BR}^\rho \otimes \mathcal{B}_{\omega^\omega}$ -measurable and $g(x, \cdot)$ is continuous, open and onto $G(x)$, for each x (cf. [47]). Define $G' \subseteq \omega^\omega \times \omega^\omega$ by

$$G' = \{(x, y) : g(x, y) \in F(x)\}.$$

As $\text{Gr}(F) \in \mathcal{BR}^\rho$ and g is $\mathcal{BR}^\rho \otimes \mathcal{B}_{\omega^\omega}$ -measurable, G' is in \mathcal{BR}^ρ .

Also, as the inverse image of a non-meager set under a continuous open map is non-meager, G' has non-meager sections. By 7.11, G' has a \mathcal{BR}^ρ -measurable selection g' . Then $f(x) = g(x, g'(x))$ is a \mathcal{BR}^ρ -measurable selection for F .

PART II

THE FINITE LEVELS OF THE HIERARCHY OF EFFECTIVE
R-SETS AND THE EFFECTIVE C-SETS

1. Until now we have only considered the boldface pointclasses

Σ_1^0, Π_1^0 etc. and the classical R-sets and have studied their various

properties. We have, however, formulated and proved many of our results

so as to carry them over to the effective case without much difficulty.

We first recall the definition of the effective pointclasses. For any

p.a.o. Φ ,

$$\Sigma_1^\Phi = \left\{ \Phi(\{E_n : n \in \omega\}) : \{E_n\} \text{ is a recursive family} \right\},$$

where a family $\{E_n\}$ of subsets of a space of type 0 or 1 is recursive
iff the relation $x \in E_n$ is recursive;

$$\Pi_1^\Phi = \left\{ \neg E : E \in \Sigma_1^\Phi \right\}$$

and $\Pi_1^\Phi \cap \Sigma_1^\Phi = \Delta_1^\Phi.$

For each n , the n th level of the hierarchy of the effective R-sets

is defined to be $\Sigma_1^{R_n}$ and will be denoted by \mathcal{R}_n . We denote by ${}^c\mathcal{R}_n$

the class of complements of \mathcal{R}_n -sets. Plainly, $\mathcal{R}_0 = \Sigma_1^1$ and by results

of Aczel, Hinman, Solovay ${}^c\mathcal{R}_1 = \Sigma_1^1\text{-ind} = \mathcal{O}\Sigma_2^0$. The effective R-sets

have been studied by Hinman in [20] and [21]. However, Hinman was

more interested in the various aspects of the interplay between the

operator \mathcal{R} and hierarchies of recursion theory and considered various

hierarchies in $\mathcal{R}_n \cap {}^c\mathcal{R}_n$ (see [21]). He has also shown that ${}^c\mathcal{R}_{n+1}$

is precisely those sets which are semi-recursive in $\mathbb{R}_n^\#$ ([21]). Here, we shall be more interested in the "set-theoretic" aspects much in the spirit of (Chapter 6 of) Moschovakis [38]. For instance, we obtain the prewellordering property of $c\mathcal{R}_n$ by adapting Blackwell's proof of the reduction property for Π_1^1 ([6]), thus showing that Blackwell's method can be extended to more general set-theoretic operations.

(Blackwell's idea was further generalized to obtain a wide range of structural results by Moschovakis and others; see e.g. [38], [36]). Of course, the prewellordering property of $c\mathcal{R}_n$ follows from Hinman's characterization. We prove it using Blackwell's idea not only for technical reasons but also to bring out structural similarities between these pointclasses and Π_1^1 . (Subtle differences however exist as observed by Hinman. In his seminal paper [20], he proved that Suslin-Kleene-type theorems do not hold for operations more powerful than operation A_1).

We now define the following sequence of operations.

$$Q_0 = A_1, \quad Q_{n+1} = \mathbb{R}Q_n^0.$$

Then by [21], $Q_n \leftrightarrow R_n$ and so

$$R_n = \left\{ Q_n(\{E_n\}) : \{E_n\} \text{ is recursive} \right\}.$$

Thus, for convenience, we shall work with Q_n and shall make no distinction between Q_n and R_n .

It is not very hard to see that $\{R_n : n \in \omega\}$ forms a strict hierarchy. In fact, by the usual universal set argument one can show that

for every n ,

$$\mathcal{R}_n \subsetneq \mathcal{R}_{n+1} \cap \circ \mathcal{R}_{n+1}$$

It is well-known that $\circ \mathcal{R}_0 = \Pi_1^1$ and $\circ \mathcal{R}_1 = \mathcal{O} \Sigma_2^0$ are Spector pointclasses (cf [38]). Our first result shows that this is true in general.

1.1 Theorem : For any p.a.o. $\Phi \geq \cup, \cap$; $\Pi_1^{\Phi^*}$ is a Spector pointclass.

Proof : It is easy to see that $\Pi_1^{\Phi^*}$ is an adequate pointclass and is ω -parametrized. To see that $\Pi_1^{\Phi^*}$ has the substitution property, let $P \subseteq Y$ be a $\Pi_1^{\Phi^*}$ -set and $f: X \rightarrow Y$ a partial function which is $\Pi_1^{\Phi^*}$ -recursive. Let M be the canonical base for Φ^* . Let $\{F_n\}$ be a recursive family such that

$$P(y) \leftrightarrow (\forall n \in M)(\exists n \in \eta) [y \in F_n]. \quad \dots(i)$$

Since f is partial $\Pi_1^{\Phi^*}$ -recursive. There is a set $P' \in \Pi_1^{\Phi^*}$ such that for all s ,

$$f(x) \downarrow \rightarrow [f(x) \in N(Y; s) \leftrightarrow P'(x, s)] \quad \dots(ii)$$

Since $\{F_n\}$ is recursive (and hence Σ_1^0) there is a Σ_1^0 set $U \subseteq \omega \times \omega$ such that

$$y \in F_n \leftrightarrow (\exists p)(y \in N(Y; p) \ \& \ U(p, n)). \quad \dots(iii)$$

Define $P^*(x) \leftrightarrow (\forall n \in M)(\exists n \in \eta)(\exists p) [P'(x, p) \ \& \ U(p, n)]$.

By Theorem I.2.17 it is not hard to see that $P^* \in \Pi_1^{\Phi^*}$. We claim that

for every x ,

$$f(x) \downarrow \rightarrow [P^*(x) \leftrightarrow P(f(x))].$$

So assume $f(x) \downarrow$. Then

$$\begin{aligned} P(f(x)) &\leftrightarrow (\forall n \in M)(\exists n \in \eta)[f(x) \in F_n] && \text{by (i)} \\ &\leftrightarrow (\forall n \in M)(\exists n \in \eta)(\exists p)\{f(x) \in N(\gamma; p) \ \& \ U(p, n)\} && \text{by (iii)} \\ &\leftrightarrow (\forall n \in M)(\exists n \in \eta)(\exists p)\{p'(x, p) \ \& \ U(p, n)\} && \text{by (ii)} \\ &\leftrightarrow P^*(x). \end{aligned}$$

Thus $\Pi_1^{\Phi^*}$ has the substitution property. It remains to show that $\Pi_1^{\Phi^*}$ is normed. To see this let N be the canonical base for $\Phi \Phi^0$ and $F \subseteq X$ be in $\Pi_1^{\Phi^*}$. Suppose

$$\uparrow F = \Phi^*(\{E_n : n \in \omega\}),$$

where $\{E_n : n \in \omega\}$ is a regular recursive family, i.e.

$$E_s = X, \text{ and } \forall s, t \in \text{Seq} [s \subseteq t \rightarrow E_t \subseteq E_s].$$

For $x, y \in X$, we define a game $G(x, y)$ as follows. The game is played between \forall (I) and \exists (II). First \forall plays ξ_0 from N , \exists replies with $n_0 \in \xi_0$ and an $\eta_0 \in N$; \forall then plays an $m_0 \in \eta_0$ and $\xi_1 \in N$. The game then proceeds as described above. We say that \exists wins if for some k , $x \notin E_{\langle n_0, \dots, n_k \rangle}$ and $y \in E_{\langle m_0, \dots, m_{k-1} \rangle}$.

We define $x \leq y$ to hold iff $x, y \in F$ and \exists has a winning strategy in the game $G(x, y)$. We shall show that \leq is a prewellordering. (cf [36, 23]).

1^o. We claim that there is no infinite sequence x_i in F such that $x_i \neq x_{i+1}$ for all i .

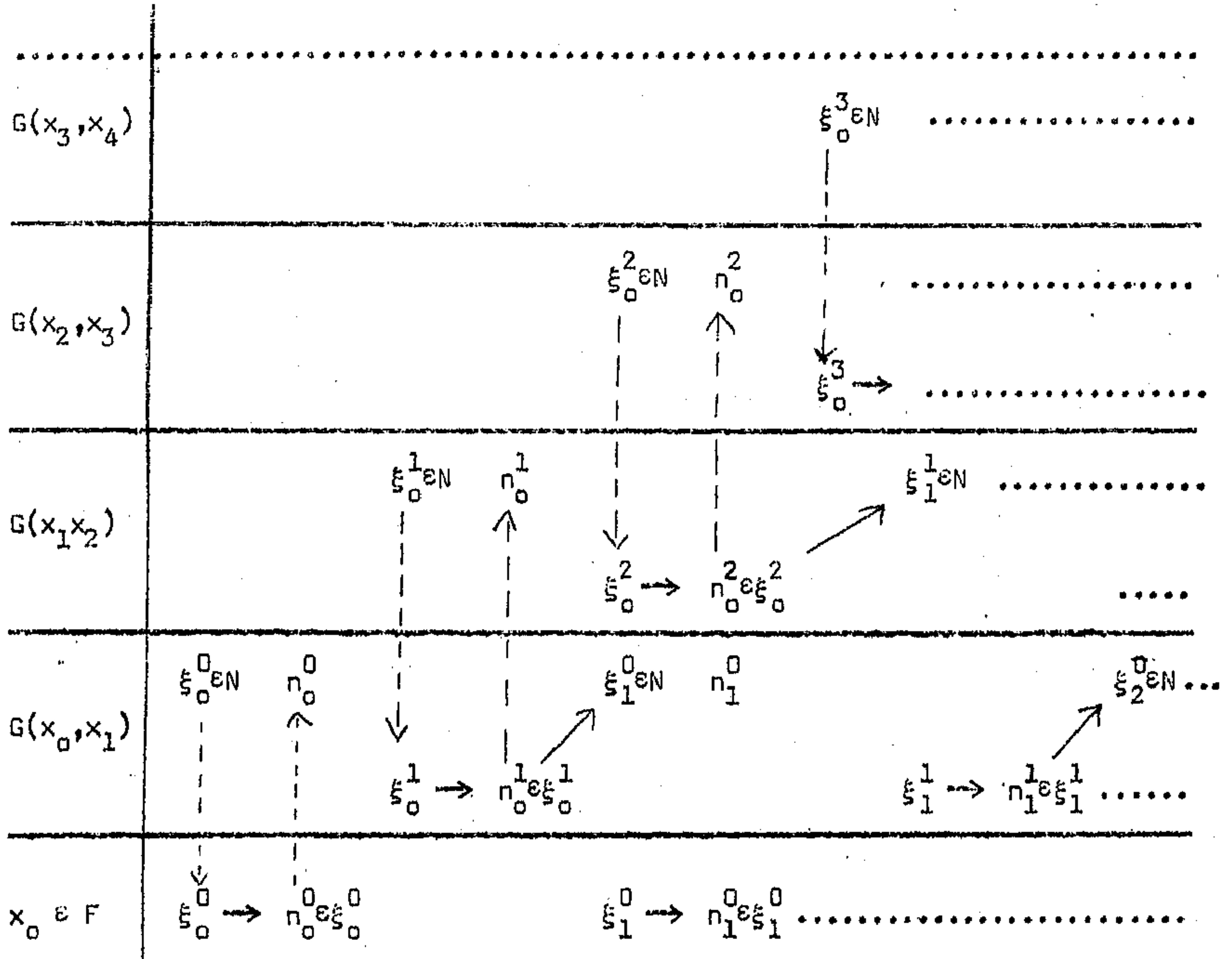
Suppose such a sequence exists. Since $x_0 \in F$ we have

$$(\forall \xi_0 \in N)(\exists n_0 \in \xi_0)(\forall \xi_1 \in N)(\exists n_1 \in \xi_1) \dots (\exists k) [x_0 \notin E_{\langle n_0, \dots, n_{k-1} \rangle}] \dots (iv)$$

Fix a winning strategy σ for \exists in this game. Now notice that for each i , II does not win the game $G(x_i, x_{i+1})$ and since the game is determined I wins. Thus for each i ,

$$(\exists \xi_0 \in N)(\forall n_0 \in \xi_0)(\forall \eta_0 \in N)(\exists m_0 \in \eta_0) \dots \dots (\forall k) [x_i \in E_{\langle n_0, \dots, n_k \rangle} \vee x_{i+1} \notin E_{\langle m_0, \dots, m_{k-1} \rangle}] \dots (v)$$

Let σ_1 be a winning strategy for \exists in (v). We now simulate the following sequence of games in which dotted lines indicate simulated moves and solid arrows show responses as dictated by the fixed winning strategies.



From the bottom game we have, for some $k' \in \omega$, $x_0 \notin E_{\langle n_0^0, n_1^0, \dots, n_{k'-1}^0 \rangle}$.

Since player I plays according to his winning strategy σ_0 in the game $G(x_0, x_1)$ we have $(\forall k) \left[x_0 \in E_{\langle n_0^0, \dots, n_k^0 \rangle} \vee x_1 \notin E_{\langle n_0^1, \dots, n_{k-1}^1 \rangle} \right]$.

Thus, putting $k = k' - 1$, we get

$$x_1 \notin E_{\langle n_0^1, \dots, n_{k'-2}^1 \rangle}.$$

At the next stage we shall obtain

$$x_2 \notin E_{\langle n_0^2, \dots, n_{k'-3}^2 \rangle}.$$

Proceeding thus we see that at stage k' , $x_{k'} \notin E_\emptyset$ and thus obtain a contradiction.

2^o. By 1^o we clearly have $x \preceq x$, for all x . Furthermore, we cannot have $\neg(x \preceq x' \vee x' \preceq x)$. Finally, if $x_1 \preceq x_2$ & $x_2 \preceq x_3$, then combining the two strategies gives $x_1 \preceq x_3$. Thus \preceq is a prewellordering of F . Let ψ be the norm it induces on F .

3^o. Note that \exists has a winning strategy in the game $G(x,y)$ iff $x \in F$ and either $y \notin F$ or $\psi(x) \leq \psi(y)$. Thus

$$x \leq_{\psi}^* y \iff \exists \text{ wins } G(x,y) \\ \iff (\forall \xi_0 \in \mathbb{N})(\exists \eta_0 \in \xi_0)(\exists \eta_0 \in \mathbb{N})(\forall m_0 \in \eta_0) \dots \\ \dots \exists k \left[x \notin E_{\langle n_0, \dots, n_k \rangle} \text{ \& } y \in E_{\langle m_0, \dots, m_{k-1} \rangle} \right].$$

The last expression clearly defines a set in \prod_1^* .

4^o. We now claim that

$$x <_{\psi}^* y \iff (\exists \eta_0 \in \mathbb{N})(\forall m_0 \in \eta_0)(\forall \xi_0 \in \mathbb{N})(\exists n_0 \in \xi_0) \dots \\ \dots (\exists k) \left[x \notin E_{\langle n_0, \dots, n_{k-1} \rangle} \text{ \& } y \in E_{\langle m_0, \dots, m_{k-1} \rangle} \right] \dots (vi)$$

To see this suppose \exists wins the above game. Then clearly $x \in F$. If $y \notin F$ then $x <_{\psi}^* y$. So assume $y \in F$. If $\neg(x <_{\psi}^* y)$ then we must have $\psi(y) \leq \psi(x)$ and hence \exists wins $G(y,x)$, i.e.

$$\dots (\forall \eta_0 \in \mathbb{N})(\exists m_0 \in \eta_0)(\exists \xi_0 \in \mathbb{N})(\forall n_0 \in \xi_0) \dots \\ \dots (\exists k) \left[y \notin E_{\langle m_0, \dots, m_k \rangle} \text{ \& } x \in E_{\langle n_0, \dots, n_{k-1} \rangle} \right].$$

Combining the two games we can simulate moves such that at the end of the plays we shall obtain k_1 and k_2 satisfying

$$x \notin E_{\langle n_0, \dots, n_{k_1-1} \rangle} \text{ and } y \in E_{\langle m_0, \dots, m_{k_1-1} \rangle}$$

$$x \in E_{\langle n_0, \dots, n_{k_2-1} \rangle} \text{ and } y \notin E_{\langle m_0, \dots, m_{k_2-1} \rangle}.$$

By regularity of $\{E_n\}$ this gives a contradiction. We therefore have $x \prec_{\psi}^* y$.

Conversely suppose $x \prec_{\psi}^* y$. If $x \in F$ and $y \notin F$ then plainly \exists wins (vi). So assume $x \in F$ and $y \in F$ and $\psi(x) < \psi(y)$. Then we have $\neg(y \leq x)$. Thus

$$(\exists \eta_0 \in \mathbb{N})(\forall m_0 \in \eta_0)(\forall \xi_0 \in \mathbb{N})(\exists n_0 \in \xi_0) \dots$$

$$\dots (\forall k) \left[y \in E_{\langle m_0, \dots, m_k \rangle} \vee x \notin E_{\langle n_0, \dots, n_{k-1} \rangle} \right]. \dots \text{(vii)}$$

If \exists does not win (vi) then we have

$$(\forall \eta_0 \in \mathbb{N})(\exists m_0 \in \eta_0)(\exists \xi_0 \in \mathbb{N})(\forall n_0 \in \xi_0) \dots$$

$$\dots (\forall k) \left[x \in E_{\langle n_0, \dots, n_{k-1} \rangle} \vee y \notin E_{\langle m_0, \dots, m_{k-1} \rangle} \right] \dots \text{(viii)}$$

From (vii), (viii) and the game witnessing $y \in F$ we can play a series of games as in 1^0 and arrive at a contradiction. This shows that \exists must win the game (vi). Thus $\prec_{\psi}^* \in \Pi_1^{\Phi^*}$. This completes the proof.

1.2 Corollary: For each n , $\circ \mathcal{R}_n$ is a Spector pointclass. \square

The effective versions of 1.6.6, 1.6.7 and 1.6.12 give us the following. The proofs follow the same pattern and are therefore omitted.

1.3 Theorem : Let $\Gamma = \mathcal{R}_n$ or $c\mathcal{R}_n$. Then for any $A \in \Gamma$ and any s, n, r the following sets are also in Γ (uniformly).

$$(a) A^{*\Sigma(s)} = \{x : A^x \text{ is comeager in } \Sigma(s)\}$$

$$(b) A^{*[n,r]} = \{x : \mu(A^x \cap G^n) \geq r\}$$

$$(c) A^{*(n,r)} \stackrel{\text{def}}{=} \{x : \mu(A^x \cap G^n) > r\}. \quad \square$$

The next result is an effective analogue of Theorem 1.3.8 and can be viewed as the counterpart of the third periodicity theorem of Moschovakis [38].

1.4 Theorem : Suppose $F \in \Pi_1^{\mathbb{R}\Phi_N}$ and assume that

$$F(x) \leftrightarrow (\forall \eta_0 \in N)(\exists n_0 \in \eta_0)(\forall \eta_1 \in N)(\exists n_1 \in \eta_1) \dots \dots \dots \\ \dots \dots (\exists k) \left[\neg E(x, \langle n_0, \dots, n_{k-1} \rangle) \right] \quad \dots (i)$$

where E is a regular recursive relation. Assume, moreover, that Φ_N is normal.

Then, there is a $\Pi_1^{\mathbb{R}\Phi_N}$ -recursive partial function τ such that $\tau_x = \tau(x, \cdot)$ is a winning strategy for \exists in the above game, whenever $x \in F$.

Proof : Let Γ be the canonical inductive operator associated with N and E so that

$$s \in \Gamma_x(A) \leftrightarrow \neg E(x, s) \vee (\forall \eta \in N)(\exists n \in \eta) [s * \langle n \rangle \in A].$$

Plainly, $F(x) \leftrightarrow e \in \Gamma_x^\omega$.

Define $\beta(s, x) = \begin{cases} \text{least } \rho \text{ such that } s \in \Gamma_x^\rho, \text{ if such a } \rho \text{ exists} \\ \omega_1 \text{ otherwise.} \end{cases}$

Clearly, $F(x) \leftrightarrow \beta(e, x) < \omega_1$.

Define $F^s(x) \leftrightarrow s \in \Gamma_x^\omega$.

This relation in x and s is in $\Pi_1^{\mathbb{R}\overline{\mathbb{Q}}_N}$. Now let $x \in F$, so \exists wins the above game. If $\eta_0, \eta_1, \dots, \eta_{i-1}; n_0, n_1, \dots, n_{i-1}$ are the first $2i$ relevant moves of \forall and \exists , observe that \exists is in a winning position i.e. he goes on to win the game (if he has not yet done so) iff $\langle n_0, \dots, n_{i-1} \rangle \in \Gamma_x^\omega$ i.e. iff $\beta(\langle n_0, \dots, n_{i-1} \rangle, x) < \omega_1$. Hence, corresponding to any subsequent move $\eta \in N$ made by \forall , \exists has to play an $n \in \eta$ such that $\beta(\langle n_0, \dots, n_{i-1}, n \rangle, x) < \beta(\langle n_0, \dots, n_{i-1} \rangle, x)$ holds. Playing thus \exists ensures that he is in a winning position at each stage and ultimately arrives at a stage k when $\beta(\langle n_0, \dots, n_{k-1} \rangle, x) = 0$ i.e. $\exists \in E(x, \langle n_0, \dots, n_{k-1} \rangle)$ and, therefore, wins the game (i). We shall show that \exists can ensure this in a $\Delta_1^{\mathbb{R}\overline{\mathbb{Q}}_N}(x)$ -recursive way. Put

$$\begin{aligned} & \text{Min}(x, \langle n_0, \dots, n_{k-1} \rangle, \langle n_0, \dots, n_{k-1} \rangle) \\ \leftrightarrow & (\forall i < k)(\eta_i \in N \ \& \ n_i \in \eta_i) \ \& \ x \in F^{\langle n_0, \dots, n_{k-1} \rangle} \ \& \\ & (\forall v) \{ (\text{Seq}(v) \ \& \ \text{lh}(v) = k \ \& \ (\forall i < k-1)(v)_i = n_i) \ \& \ (v)_{k-1} \in \eta_{k-1} \} \\ & \rightarrow \beta(x, \langle n_0, \dots, n_{k-1} \rangle) \leq \beta(x, v) \}. \end{aligned}$$

The proof of Theorem 1.1 shows that the relation

$$\beta(x,u) < \omega_1 \ \& \ \beta(x,u) \leq \beta(x,v)$$

is in $\prod_1^{\mathbb{R}\bar{\mathbb{Q}}_N}$. It follows that the relation Min is in $\prod_1^{\mathbb{R}\bar{\mathbb{Q}}_N}$.

Call $\langle a_0, \dots, a_k \rangle$ best with respect to x and η_0, \dots, η_k if it is "minimal" with respect to x and η_0, \dots, η_k and if, moreover, there is no $b < a_k$ in η_k such that $\langle a_0, \dots, a_{k-1}, b \rangle$ is minimal.

Thus

$$\begin{aligned} & \text{Best}(x, \langle \eta_0, \dots, \eta_k \rangle, \langle a_0, \dots, a_k \rangle) \\ \iff & \text{Min}(x, \langle \eta_0, \dots, \eta_k \rangle, \langle a_0, \dots, a_k \rangle) \ \& \\ & (\forall b < a_k) \left[b \in \eta_k \rightarrow \beta(x, \langle a_0, \dots, a_k \rangle) < \beta(x, \langle a_0, \dots, a_{k-1}, b \rangle) \right] \end{aligned}$$

The relation Best is also in $\prod_1^{\mathbb{R}\bar{\mathbb{Q}}_N}$. Finally, τ is defined as follows

$$\begin{aligned} & \tau(x; \eta_0, \eta_0, \eta_1, \eta_1, \dots, \eta_{k-1}, \eta_{k-1}, \eta_k) \downarrow \\ \iff & F(x) \ \& \ x \in F^s \ \& \ \beta(x,s) > 0. \end{aligned}$$

In this case, put

$$\begin{aligned} & \tau(x; \eta_0, \eta_0, \dots, \eta_{k-1}, \eta_{k-1}, \eta_k) = a_k \\ \iff & (\exists a_0) \dots (\exists a_{k-1}) \left[(\forall j < k) \text{Best}(x, \langle \eta_0, \dots, \eta_j \rangle, \langle a_0, \dots, a_j \rangle) \right. \\ & \left. \ \& \ \text{Best}(x, \langle \eta_0, \dots, \eta_k \rangle, \langle a_0, \dots, a_k \rangle) \right]. \end{aligned}$$

It is quite routine to check that if $x \in F$, τ_x is a winning strategy for \exists in the game (i). \square

The natural question that arises now is whether the Scale-Transfer Theorem of Moschovakis [38] can be generalized to the type of games that we have dealing with so far. We do not know whether this can be done. However, we can establish the scale property for $\circ\mathcal{R}_n$ by observing the following facts. First note that by extending the notion of presentability to any pointclass, one can obtain an effective version of Theorem I.4.6. Then observe ^{that} the result of Burgess and Kaniewski (Lemma I.5.1) can be effectivized using the canonical scales on \prod_1^1 sets and the effective analogue of the theorem of Von-Neumann-Yankov. Using these facts and the effective version of Corollary I:3.5 (which follows easily from the proof of the norm property in Theorem 1.1) one can effectivize Theorem I:5.3. Without going into any further details we announce

1.5 Theorem : For every n , $\circ\mathcal{R}_n$ has the scale property. \square

The above results show that for each n , $\circ\mathcal{R}_n$ is a "nice" pointclass with pleasant structural properties and is therefore an interesting class to study.

2. The effective C-sets : In this section we shall consider the problem of effectivizing the $\mathcal{BR}^0 = \mathcal{C}$ -sets of Selivanovskij in a rather unorthodox fashion. The effective C-sets have already been constructed by Hinman in [20] where he has given a fairly general and satisfactory method of effectivizing $\nabla(\mathcal{Q})$ on ω for operations \mathcal{Q} more powerful than \cup . Hinman gets the effective hierarchy by generalizing Addison's construction of "effective" Borel hierarchy which, roughly, consists

of alternating applications of r.e. union and complementation. Let us briefly discuss Hinman's method. We shall confine ourselves to operation \mathcal{A} for simplicity. Briefly, Hinman's method consists of assigning to each set C as it is generated an index or code $i(C)$ and at each stage of the inductive definition applying operation \mathcal{A} to those families $\{F_n\}$ of subsets of ω for which the function $n \rightarrow i(F_n)$ is recursive in some set previously generated. The class of sets thus obtained is exactly the class of sets recursive in \mathbb{E}_1 , the type-2 object associated with operation \mathcal{A} . This generalizes the result for the "effective" Borel sets or HYP, since HYP is precisely the class of sets recursive in ${}^2\mathbb{E}$, the Kleene's type-2 object associated with \cup . More specifically, let I be the smallest subset of ω such that for all $a \in I$, there exist $P_a \subseteq \omega$ which satisfy the following conditions: for all a, b, c

(i) if c is an index of a partial recursive function on ω , then $\langle 7, c \rangle \in I$ and $P_{\langle 7, c \rangle} = \text{Dm}\{c\}$;

(ii) if $b \in I$ and for all k , $\{a\}(k, P_b) \in I$, then $\langle a, b \rangle \in I$ and $P_{\langle a, b \rangle} = \mathcal{A}(\{ \uparrow P_{\{a\}}(k, P_b) : k \in \omega \})$.

This is an inductive definition whose closure ordinal is $\omega_1[\mathbb{E}_1]$, the least ordinal not the order-type of a well-ordering of ω recursive in \mathbb{E}_1 (cf [22]). Put

$$\nabla(\mathcal{A}) = \{P_a : a \in I\}.$$

Then by [20; Theorem 2.10],

$$\nabla(\mathcal{A}) = {}_1\text{sc}(\mathbb{E}_1), \text{ the class of subsets of } \omega \text{ recursive in } \mathbb{E}_1.$$

This, in short, is the effective C-sets obtained by Hinman by appropriately generalizing Addison's construction.

Now observe that the C-sets can be obtained by different methods. For example, as shown in [13], the C-sets are obtained by applying the game quantifier \mathcal{G} to \mathcal{A}_2^0 sets. In fact, by applying \mathcal{G} to the λ th level of the "difference hierarchy" of \mathcal{A}_2^0 given by alternated series of decreasing closed sets (of [28, § 34 VI]) one obtains the λ th level of the hierarchy of C-sets of Selivanovskij. This naturally raises the following question. Can one obtain the effective C-sets by applying \mathcal{G} to "effective" alternated series of Π_1^0 sets of appropriate lengths? We feel that this should be possible. In this section we shall confine ourselves to another, albeit unorthodox, method of obtaining the hierarchy of the classical C-sets and then obtain an effective analogue. But first we need the following definitions.

2.1 Definition : $S_1(\alpha) \stackrel{\text{def}}{\longleftrightarrow} \alpha$ codes some tree on ω
 $\longleftrightarrow (\forall s) [\alpha(s) = 0 \rightarrow (\text{Seq}(s) \& (\forall t)(\text{Seq}(t) \& t \subseteq s \rightarrow \alpha(t) = 0))]$.

$S_2(\alpha, s) \stackrel{\text{def}}{\longleftrightarrow} s$ is terminal for the tree coded by α
 $\longleftrightarrow S_1(\alpha) \& \alpha(s) = 0 \& (\forall n) (\alpha(s * \langle n \rangle) \neq 0)$.

$WF(\alpha) \stackrel{\text{def}}{\longleftrightarrow} \alpha$ codes a wellfounded (wff) tree on ω
 $\longleftrightarrow S_1(\alpha) \& (\forall \beta) (\exists i) [\alpha(\bar{\beta}(i)) \neq 0]$
 $\longleftrightarrow S_1(\alpha) \& (\forall \beta) (\exists i) S_2(\alpha, \bar{\beta}(i))$.

Plainly S_1 and S_2 are Π_1^0 while WF is Π_1^1 .

If T is a wff tree on ω then $|T|$ denotes its length. If $T \neq \emptyset$, then $|T| = \rho(e)$, where ρ is the rank function (of [38; 2D]). If α codes a wff tree T then put $\|\alpha\| = |T|$.

Let X be a space of type 0 or 1 and let $G \subseteq \omega^\omega \times X$ be a good universal for $\Sigma_1^0 \upharpoonright X$. Define

$$H(\alpha, x) \leftrightarrow G(\alpha^*, x) \ \& \ \alpha(0) = 0 \vee \neg G(\alpha^*, x) \ \& \ \alpha(0) \neq 0,$$

where $\alpha^* = \lambda n. \alpha(n+1)$.

Now, define $D^X \subseteq \omega^\omega \times \omega^\omega \times X$ (or D when X is clear from the context) as follows.

$$D(\gamma, \delta, x) \leftrightarrow WF(\gamma) \ \& \ (\exists \alpha_0)(\forall n_0)(\forall \alpha_1)(\exists n_1)(\exists \alpha_2)(\forall n_2) \dots \\ \dots (\exists k) \left\{ s_2(\gamma, \underbrace{\langle \bar{\alpha}_0(n_0), \dots, \bar{\alpha}_{k-1}(n_{k-1}) \rangle}_t) \ \& \ H((\delta)_t, x) \right\}$$

Notice that the above equivalence can also be written as

$$D(\gamma, \delta, x) \leftrightarrow WF(\gamma) \ \& \ (\exists \alpha_0)(\forall n_0)(\forall \alpha_1)(\exists n_1) \dots \\ \dots (\forall k) \left\{ s_2(\gamma, \underbrace{\langle \bar{\alpha}_0(n_0), \dots, \bar{\alpha}_{k-1}(n_{k-1}) \rangle}_t) \rightarrow H((\delta)_t, x) \right\} (*)$$

This is because for any wff tree T and any sequence $\{m_i\}$ there is a unique i such that $\langle m_0, \dots, m_{i-1} \rangle$ is terminal for T . From the above two equivalences it follows that D is in $\Delta_1^{\mathcal{R}_1} = \mathcal{C}\mathcal{R}_1 \cap \mathcal{R}_1$.

The set D above (for spaces X of type 1) was constructed by Burgess and Lockhart in [14] where they show that D is universal for

the C-sets. (In fact they prove that D is a Borel-programmable set and so the class of C-sets is strictly included in the class of Borel-programmable sets of Blackwell (see [14; § 14])).

A set $A \subseteq \mathfrak{X}$, a space of type-1, is at the λ th level ($\lambda < \omega_1$) of the hierarchy of C-sets iff there exist $\gamma \in WF$, $\delta \in \omega^\omega$, $\|\gamma\| \leq \lambda$ and $A = D_{\gamma, \delta}$.

The effectivization is achieved by taking sections of D by wff trees and reals of increasing complexity (obtained à la Hinman). We first show that D is recursive in \mathbb{E}_1 .

2.2 Lemma: For any space \mathfrak{X} (of type 0 or 1), $D^{\mathfrak{X}}$ is recursive in \mathbb{E}_1 .

Proof: We shall omit the superscript. We first fix a recursive function $(\delta, s) \rightarrow \delta^s$ from $\omega^\omega \times \omega \rightarrow \omega^\omega$ such that $(\delta^s)_t = (\delta)_{s * t}$. If γ codes a tree T then γ_s is a code for $T_s = \{t : s * t \in T\}$ such that the function $(\gamma, s) \rightarrow \gamma_s$ is recursive. Now define a partial function φ as follows.

$$\begin{aligned} \varphi(\alpha, \gamma, \delta, x) &\stackrel{1}{=} 0 && \text{if } WF(\gamma) \ \&\ \|\gamma\| = 0 \ \&\ G((\delta)_\alpha, x) \\ &\stackrel{2}{=} \mathbb{E}_1(\lambda k. \mathbb{E}_1^0(\lambda k'. \{0\}(\gamma_{\langle k, k' \rangle}, \delta^{\langle k, k' \rangle}, x))) && \\ &&& \text{if } WF(\gamma) \ \&\ \|\gamma\| > 0 \\ &\stackrel{3}{=} 1 && \text{otherwise;} \end{aligned}$$

$$\text{where } \mathbb{E}_1^0(\alpha) = \begin{cases} 0 & \text{if } (\forall \beta)(\exists m)(\alpha(\bar{\beta}(m)) = 0) \\ 1 & \text{if } (\exists \beta)(\forall m)(\alpha(\bar{\beta}(m)) \neq 0). \end{cases}$$

Clearly φ is partial recursive in \mathbb{E}_1 . By the recursion theorem there is an index c^* such that

$$\varphi(c^*, \gamma, \delta, x) \stackrel{?}{=} \{c^*\}^{\mathbb{E}_1}(\gamma, \delta, x).$$

It is easy to check by induction on $\|\gamma\|$ that $\{c^*\}^{\mathbb{E}_1}$ is the characteristic function of D . Thus D is recursive in \mathbb{E}_1 . \square

Since D^{ω^ω} is recursive in \mathbb{E}_1 , any effective hierarchy we wish to construct over ω^ω will not exhaust ${}^2\text{sc}(\mathbb{E}_1)$ as can be shown by diagonalizing. We therefore confine ourselves to ω and from now on denote D^ω by D .

2.3 Definition: Let \tilde{I} be the smallest subset of ω such that for every $a \in \tilde{I}$ there exist $\text{tr}(a) \in \omega^\omega$ coding a wff tree and a real $r(a)$ satisfying the following conditions.

For all a, b, b', c, k, s, n, m

(i) if c is the index of a total recursive function on ω , then $\langle 17, c \rangle \in \tilde{I}$ and $\text{tr}(\langle 17, c \rangle)$ is the canonical code (i.e. the characteristic function) for the tree $\{e\}$ and

$$r(\langle 17, c \rangle)(\langle e, n \rangle) = \{c\}(n)$$

$$r(\langle 17, c \rangle)(m) = 0 \quad \text{otherwise ;}$$

(ii) if $b, b' \in \tilde{I}$ and for all k , $\{a\}(k, D_{\text{tr}(b), r(b')}) \in \tilde{I}$, then $\langle a, b, b' \rangle \in \tilde{I}$ and

$$\text{tr}(\langle a, b, b' \rangle)(e) = 0$$

$$\text{tr}(\langle a, b, b' \rangle)(ks) = \begin{cases} 0 & \text{if } s = e \vee [\text{Seq}(s) \ \& \\ & \text{tr}(\{a\}(k, D_{\text{tr}(b), \text{r}(b')})(s) = 0)] \\ 1 & \text{otherwise} \end{cases}$$

$$\text{r}(\langle a, b, b' \rangle)(\langle ks, n \rangle) = \text{r}(\{a\}(k, D_{\text{tr}(b), \text{r}(b')})(\langle s, n \rangle),$$

otherwise $\text{r}(\langle a, b, b' \rangle)(m) = 1$.

Now put

$$\mathcal{K} = \left\{ D_{\text{tr}(a), \text{r}(a')} : a, a' \in \tilde{\mathbb{I}} \ \& \ |a| = |a'| \right\}.$$

Obviously, there is a natural way of assigning ordinals to each $a \in \tilde{\mathbb{I}}$ (viz $|a|$) so that one can define a natural hierarchy whose scope is \mathcal{K} . We shall first show that it is a hierarchy in ${}_{1}\text{sc}(\mathbb{E}_1)$.

2.4 Lemma : There are primitive recursive functions σ_1 and σ_2 such that for all $a \in \tilde{\mathbb{I}}$,

$$\text{tr}(a) = \lambda n. \left\{ \sigma_1(a) \right\}^{\mathbb{E}_1}(n)$$

$$\text{and } \text{r}(a) = \lambda n. \left\{ \sigma_2(a) \right\}^{\mathbb{E}_1}(n).$$

Proof : Define two partial functions ψ_1, ψ_2 as follows. For any

$e_1, e_2, a, b, b', c, n, k, s, u$

$$(i) \quad \psi_1(e_1, e_2, \langle 17, c \rangle, n) = \begin{cases} 0 & \text{if } n = e \\ 1 & \text{otherwise.} \end{cases}$$

(ii) $\psi_1(e_1, e_2, \langle a, b, b' \rangle, t) = 0$ if $t = e \vee \text{Seq}(t) \ \& \ \text{lh}(t) = 1$

$$\psi_1(e_1, e_2, \langle a, b, b' \rangle, ks) \stackrel{\text{IE}_1}{\simeq} \left\{ e_1 \right\}^{\text{IE}_1} (e_2, \{a\} (k, D \lambda n. \left\{ e_1 \right\}^{\text{IE}_1} (e_2, b, n), \lambda n. \left\{ e_2 \right\}^{\text{IE}_1} (e_1, b', n)), s)$$

$\psi_1(e_1, e_2, \langle a, b, b' \rangle, n) = 1$ if $n \notin \text{Seq}$

(iii) $\psi_1(e_1, e_2, u, n)$ is undefined otherwise.

(i') $\psi_2(e_2, e_1, \langle 17, 0 \rangle, \langle e, n \rangle) \stackrel{\text{IE}_1}{\simeq} \{0\}(n)$

(ii') $\psi_2(e_2, e_1, \langle a, b, b' \rangle, \langle ks, n \rangle)$

$$\stackrel{\text{IE}_1}{\simeq} \left\{ e_2 \right\}^{\text{IE}_1} (e_1, \{a\} (k, D \lambda n. \left\{ e_1 \right\}^{\text{IE}_1} (e_2, b, n), \lambda n. \left\{ e_2 \right\}^{\text{IE}_1} (e_1, b', n)), \langle e, n \rangle)$$

$\psi_2(e_2, e_1, \langle a, b, b' \rangle, m) = 1$ if $m \notin \text{Seq} \vee (m)_0 \notin \text{Seq}$

(iii') $\psi_2(e_2, e_1, u, n)$ is undefined otherwise.

Clearly, ψ_1, ψ_2 are partial recursive in IE_1 and so by the recursion theorem there are indices \bar{e}_1, \bar{e}_2 such that

$$\psi_1(\bar{e}_1, e_2, a, n) \stackrel{\text{IE}_1}{\simeq} \left\{ \bar{e}_1 \right\}^{\text{IE}_1} (e_2, a, n)$$

and $\psi_2(\bar{e}_2, e_1, a, n) \stackrel{\text{IE}_1}{\simeq} \left\{ \bar{e}_2 \right\}^{\text{IE}_1} (e_1, a, n).$

Let σ_1 and σ_2 be primitive recursive functions such that

$$\{\sigma_1(a)\}^{\text{IE}_1}(n) \stackrel{\text{IE}_1}{\simeq} \left\{ \bar{e}_1 \right\}^{\text{IE}_1} (\bar{e}_2, a, n)$$

$$\{\sigma_2(a)\}^{\text{IE}_1}(n) \stackrel{\text{IE}_1}{\simeq} \left\{ \bar{e}_2 \right\}^{\text{IE}_1} (\bar{e}_1, a, n).$$

It is routine to check by induction on \overline{I} that σ_1 and σ_2 have the required properties. \parallel

2.5 Theorem : $\mathcal{K} = {}_1\text{sc}(\mathbb{E}_1)$.

Proof : One part of the inclusion follows from Lemmas 2.2 and 2.4. For the other inclusion we shall show

$$\nabla(A) = \{p_a : a \in I\} \subseteq \mathcal{K}.$$

In fact, we shall define partial recursive functions ρ_1 and ρ_2 such that for every $a \in I$, $\rho_1(a), \rho_2(a) \in \overline{I}$; $|\rho_1(a)| = |\rho_2(a)|$ and

$$p_a = D_{\text{tr}(\rho_1(a)), \text{r}(\rho_2(a))}. \quad \dots(1)$$

By the usual argument via the recursion theorem one can get a primitive recursive function π such that for every $a \in \overline{I}$ and $t \in \text{Seq}$, $\pi(a) \in \overline{I}$,

$$\uparrow H_{(\text{r}(a))_t} = H_{(\text{r}(\pi(a)))_t} \quad \dots(2)$$

$$\text{and } |a| = |\pi(a)|.$$

We now define two partial functions ψ_1 and ψ_2 by cases as follows.

$$(i) \quad \begin{aligned} \psi_1(a_1, a_2, \langle 7, c \rangle) &= \langle 17, c \rangle \\ \psi_2(a_2, a_1, \langle 7, c \rangle) &= \langle 17, c' \rangle, \end{aligned}$$

where $\text{Dm } \{c\} = G_{\lambda n. \{c''\}}(n)$

$$\text{and } \{c'\}(n) = \begin{cases} 0 & \text{if } n = 0 \\ \{c''\}(n) & \text{if } n > 0. \end{cases}$$

(ii) If $u = \langle a, b \rangle$, then let τ_1 and τ_2 be primitive recursive functions such that for any v, t, A ,

$$\{e_1\}(e_2, \{v\}(t, A)) \stackrel{\omega}{\sim} \{\tau_1(e_1, e_2, v)\}(t, A) \quad \dots(3)$$

and $\pi(\{e_2\}(e_1, \{v\}(t, A))) \stackrel{\omega}{\sim} \{\tau_2(e_1, e_2, v)\}(t, A) \quad \dots(4)$

Set $\psi_1(e_1, e_2, u) \stackrel{\omega}{\sim} \langle \tau_1(e_1, e_2, a), \{e_1\}(e_2, b), \{e_2\}(e_1, b) \rangle$

$$\psi_2(e_2, e_1, u) \stackrel{\omega}{\sim} \langle \tau_2(e_1, e_2, a), \{e_1\}(e_2, b), \{e_2\}(e_1, b) \rangle.$$

By the recursion theorem there are indices \bar{e}_1 and \bar{e}_2 such that

$$\psi_1(\bar{e}_1, e_2, u) \stackrel{\omega}{\sim} \{\bar{e}_1\}(e_2, u)$$

and $\psi_2(\bar{e}_2, e_1, u) \stackrel{\omega}{\sim} \{\bar{e}_2\}(e_1, u).$

Set $\rho_1(u) \stackrel{\omega}{\sim} \{\bar{e}_1\}(\bar{e}_2, u)$

$$\rho_2(u) \stackrel{\omega}{\sim} \{\bar{e}_2\}(\bar{e}_1, u).$$

We shall show by induction on I that ρ_1 and ρ_2 have the desired properties. So let $u \in I$. If u is of the form $\langle 7, c \rangle$, then (1) holds easily. So let $u = \langle a, b \rangle$, $a \neq 7$. Then

$$x \in P_u \leftrightarrow (\exists \alpha)(\forall n) \left[x \notin P_{\{a\}(\bar{a}(n), P_b)} \right]$$

$$\leftrightarrow (\exists \alpha)(\forall n) \left[x \notin D_{\text{tr}(\rho_1(\{a\}(\bar{a}(n), D_{\text{tr}(\rho_1(b)), r(\rho_2(b))})), r(\rho_2(\{a\}(\bar{a}(n), D_{\text{tr}(\rho_1(b)), r(\rho_2(b))})))} \right]$$

by induction hypothesis.

$$\begin{aligned} & \leftrightarrow (\exists \alpha) (\forall n) (\forall \alpha_0) (\exists \alpha_1) (\forall n_1) \dots \dots \\ & \dots (\exists k) \left[S_2(\text{tr}(\rho_1(\{a\}(\tilde{\alpha}(n), D_{\text{tr}(\rho_1(b))}, r(\rho_2(b))))) , \underbrace{\langle \tilde{\alpha}_0(n_0), \dots, \tilde{\alpha}_{k-1}(n_{k-1}) \rangle}_t) \right. \\ & \quad \left. \& \text{H}((r(\rho_2(\{a\}(\tilde{\alpha}(n), D_{\text{tr}(\rho_1(b))}, r(\rho_2(b)))))_t, x)) \right], \end{aligned}$$

by (*) of 2.1 and determinacy (of closed games).

$$\begin{aligned} & \leftrightarrow (\exists \alpha_0) (\forall n_0) (\forall \alpha_1) (\exists n_1) (\exists \alpha_2) (\forall n_2) \dots \dots \\ & \dots (\exists k) \left[S_2(\text{tr}(\rho_1(\{a\}(\tilde{\alpha}_0(n_0), D_{\text{tr}(\rho_1(b))}, r(\rho_2(b))))) , \underbrace{\langle \tilde{\alpha}_1(n_1), \dots, \tilde{\alpha}_k(n_k) \rangle}_t) \right. \\ & \quad \left. \& \text{H}((r(\pi(\rho_2(\{a\}(\tilde{\alpha}_0(n_0), D_{\text{tr}(\rho_1(b))}, r(\rho_2(b)))))_t, x)) \right] \text{ by (2)} \end{aligned}$$

$$\begin{aligned} & \leftrightarrow (\exists \alpha_0) (\forall n_0) (\forall \alpha_1) (\exists n_1) (\exists \alpha_2) (\forall n_2) \dots \dots \\ & \dots (\exists k) \left[S_2(\text{tr}(\tau_1(\tilde{\alpha}_1, \tilde{\alpha}_2, a))(\tilde{\alpha}_0(n_0), D_{\text{tr}(\rho_1(b))}, r(\rho_2(b))))_t, x) \right] \& \\ & \quad \left[\text{H}((r(\tau_2(\tilde{\alpha}_1, \tilde{\alpha}_2, a))(\tilde{\alpha}_0(n_0), D_{\text{tr}(\rho_1(b))}, r(\rho_2(b))))_t, x) \right], \text{ by (3) and (4)} \end{aligned}$$

$$\leftrightarrow D(\text{tr}(\langle \tau_1(\tilde{\alpha}_1, \tilde{\alpha}_2, a), \rho_1(b), \rho_2(b) \rangle), r(\langle \tau_2(\tilde{\alpha}_1, \tilde{\alpha}_2, a), \rho_1(b), \rho_2(b) \rangle), x),$$

$$\leftrightarrow D(\text{tr}(\rho_1(u)), r(\rho_2(u)), x), \text{ by definition.}$$

This completes the proof.

Remark : It is clear that the above construction is perfectly general and can be generalized to obtain a hierarchy over ${}_{1 \text{ so}}(F_{\mathbb{Q}})$, for any p.a.o \mathbb{Q} subsuming \cup . \square

PART III

SOME PROPERTIES OF LARGE Π_1^1 SETS AND THE STRUCTURE OF HYPERARITHMETICAL SETS OF AMBIGUOUS BOREL CLASSES

1. Large Π_1^1 sets : In his paper 'Measure and category in effective descriptive set theory' [24], Kechris proved a number of measure-theoretic and category-theoretic results for the analytical classes under Projective Determinacy (PD). Although separate proofs were given for measure and category, Kechris, however, unified (and extended to the higher levels of the analytical hierarchy under PD) the basis theorem of Tanaka-Sacks ([49], [40]) and Hinman-Thomason ([19], [50]) by formulating the problem in terms of σ -ideals of subsets of reals satisfying certain definability conditions. In this section we shall show that, at the first level of the analytical hierarchy, some of these results (see (1) - (5) of Theorem 1.4) are instances of more general results much in the spirit of Kechris' Basis Theorem. Except for the basis theorem, we need only to consider hereditary families \mathcal{F} (i.e. closed under subsets) satisfying a weak definability condition viz., the Π_1^1 computational formula (see Definition 1.3 below). Elements of \mathcal{F} are thought of being small sets and we shall obtain the basis theorem for large Π_1^1 sets when \mathcal{F} is a σ -ideal satisfying the Π_1^1 computational formula. The corresponding results for measure and category are, therefore, particular instances of our results. Furthermore, the basis theorem of Kechris for large Π_1^1 sets is strengthened since we prove it under a much weaker definability assumption.

Our original results and proofs (see [3]) were strengthened and modified by V.V. Srivatsa and we present below mainly the modified version.

We work in $ZF + DC$. Our notation and terminology is as in Moschovakis [38]. Standard results in effective theory used here will be found in [38]. The following definitions would be needed in the sequel.

1.1 Definition : Let $A \subseteq \omega^\omega$ be a Π_1^1 set. The core of A , denoted by $c(A)$, is the union of all Δ_1^1 subsets (possibly empty) of A .

1.2 Definition : For reals α, β we write $\alpha \leq_1 \beta$ iff $\alpha \in \Delta_1^1(\beta)$ and $\alpha =_1 \beta$ iff $\alpha \leq_1 \beta$ & $\beta \leq_1 \alpha$. Clearly $=_1$ is an equivalence relation and both \leq_1 and $=_1$ are Π_1^1 relations.

Let \mathcal{F} be a family of subsets of ω^ω and let Γ be a pointclass.

1.3 Definition : The family \mathcal{F} is said to satisfy the Γ -computational formula if every $B \in \Gamma$, $B \subseteq X \times \omega^\omega$, $X = (\omega^\omega)^k \times \omega^l$, the set

$$B^A \stackrel{\text{def}}{=} \{x \in X, B_x \notin \mathcal{F}\}$$

is also in Γ , where $B_x = \{\beta : (x, \beta) \in B\}$.

Examples : The σ -ideals of meager sets and Lebesgue measure zero sets satisfy both Σ_1^1 and Π_1^1 computational formulas ([24]); the σ -ideal of countable sets satisfy the Σ_1^1 but not the Π_1^1 computational formula, while the σ -ideal of Ramsey null sets (i.e. meager with respect to the Ellentuck topology) satisfies neither the Σ_1^1 nor the Π_1^1 computational formula.

We shall also need the following result whose proof can be found in [25].

Gandy's basis theorem: If $A \subseteq \omega^\omega$ is a non-empty Σ_1^1 set, then A contains a real α with $\omega_1^\alpha = \omega_1^{\text{ck}}$. \square

1.4 Theorem: Let \mathcal{F} be a hereditary family of subsets of reals not containing ω^ω which satisfies the Π_1^1 computational formula. Then we have the following:

(1) Every Π_1^1 set of reals $A \notin \mathcal{F}$ contains a Δ_1^1 subset $B \notin \mathcal{F}$

(2) $\{\alpha \in \omega^\omega; \omega_1^{\text{ck}} < \omega_1^\alpha\} \in \mathcal{F}$.

(3) $\alpha \in \Delta_1^1 \cap \omega^\omega \leftrightarrow \{\beta \in \omega^\omega; \alpha \leq_1 \beta\} \notin \mathcal{F}$

where $\Delta_1^1 \cap \omega^\omega$ is the set of all Δ_1^1 reals.

(4) For every Π_1^1 set $A \subseteq \omega^\omega$, $A - c(A) \in \mathcal{F}$.

If, moreover, \mathcal{F} is a σ -ideal, then

(5) every Π_1^1 set of reals $A \notin \mathcal{F}$ contains a Δ_1^1 real.

Proof (1) This is implicitly (rather dualization of) an effective analogue of Piatkiewicz's result (of [39]).

Fix a Π_1^1 norm φ on A . Suppose A is not Δ_1^1 . Let

$$C = \{\alpha \in A; \exists \beta \in A \text{ such that } \varphi(\beta) < \varphi(\alpha)\}.$$

Since \mathcal{F} satisfies the Π_1^1 computational formula, C is Σ_1^1 and since $A \notin \mathcal{F}$, $C \subseteq A$. Moreover

$$c(\alpha) \& (\varphi(\beta) \leq \varphi(\alpha)) \rightarrow c(\beta).$$

Hence there is $\alpha_0 \in A$ such that $C = \{\alpha : \alpha <_{\varphi}^* \alpha_0\}$. Since $\alpha_0 \notin C$, it follows that

$$C = \{\beta : \beta <_{\varphi}^* \alpha_0\} \notin \mathcal{F}_1.$$

Let B be a Δ_1^1 set such that $C \subseteq B \subseteq A$. Plainly B is the required set.

(2) Let $A = \{\alpha \in \omega^\omega : \omega_1^{\text{ck}} < \omega_1^\alpha\}$. Plainly A is Π_1^1 ([38, 4F]) and if $A \notin \mathcal{F}_1$ then by (1) A includes a non-empty Δ_1^1 set. But this is not possible by Gandy's basis theorem. So $A \in \mathcal{F}_1$.

(3) If α is a Δ_1^1 point then clearly $\{\beta \in \omega^\omega : \alpha \leq_1 \beta\} = \omega^\omega \notin \mathcal{F}_1$. For the converse, first observe that if α is a $\Pi_1^1(\gamma)$ singleton i.e. $\{\alpha\}$ is $\Pi_1^1(\gamma)$, then the set $E[\alpha] = \{\beta \in \omega^\omega : \alpha \leq_1 \beta\}$ is $\Pi_1^1(\gamma)$. If furthermore $E[\alpha] \notin \mathcal{F}_1$, then by the relativized version of (1), there is a $\Delta_1^1(\gamma)$ set $B \subseteq E[\alpha]$ such that $B \notin \mathcal{F}_1$. So $\alpha \leq_1 \beta$ for all $\beta \in B$. But then

$$\alpha(n) = m \iff (\forall \beta) [\beta \in B \rightarrow (\exists \delta \in \Delta_1^1(\beta)) (\delta \in \{\alpha\} \ \& \ \delta(n) = m)],$$

so that α is a $\Delta_1^1(\gamma)$ real. Hence the set

$$G = \{\alpha \in \omega^\omega : \{\beta : \alpha \leq_1 \beta\} \notin \mathcal{F}_1\}$$

contains no non-trivial $\Pi_1^1(\gamma)$ singleton for any γ . Also by hypothesis G is Π_1^1 . We now claim that G is thin i.e. contains no perfect set. If not, then there is a $\Delta_1^1(\gamma)$, one-one function $g : 2^\omega \rightarrow G$ for some γ . But 2^ω contains a non-trivial $\Pi_1^1(\gamma)$ singleton and so G contains a non-trivial $\Pi_1^1(\gamma)$ singleton. But this is not possible. So G is thin and consequently, $G \subseteq C_1$, the largest thin Π_1^1 set (cf [25]). Since G

is closed under \leq_1 and G contains no non-trivial Π_1^1 singleton G must be the set of Δ_1^1 reals (see [25; §1B]). This completes the proof of (3).

(4) Let $A \subseteq \omega^\omega$ be a Π_1^1 set. Let $A' = \{\alpha \in \omega^\omega : \omega_1^{\text{ck}} = \omega_1^\alpha\}$. Fix a recursive function $f : \omega^\omega \rightarrow \omega^\omega$ such that $A(\alpha) \leftrightarrow \text{WO}(f(\alpha))$. Suppose $\alpha \in A'$. Then $f(\alpha) \in \text{WO}$ and $\omega_1^\alpha = \omega_1^{\text{ck}}$. But $f(\alpha)$ is $\Delta_1^1(\alpha)$ and hence $|f(\alpha)| < \omega_1^\alpha = \omega_1^{\text{ck}}$. For each $\xi < \omega_1^{\text{ck}}$, let $A_\xi^* = \{\alpha \in A : |f(\alpha)| \leq \xi\}$. Then $A' \subseteq \bigcup_{\xi < \omega_1^{\text{ck}}} A_\xi^*$. By boundedness, $\bigcup_{\xi < \omega_1^{\text{ck}}} A_\xi^* = c(A)$, the union of all Δ_1^1 subsets of A . Thus

$A - c(A) \subseteq \{\alpha : \omega_1^{\text{ck}} < \omega_1^\alpha\}$ and by (2), $A - c(A) \in \mathcal{F}$.

Assume now that \mathcal{F} is a σ -ideal.

(5) Let $A \subseteq \omega^\omega$ be a Π_1^1 set, $A \notin \mathcal{F}$. By (1) get a Δ_1^1 set $B \subseteq A$ such that $B \notin \mathcal{F}$. Fix a recursive function $f : \omega^\omega \rightarrow \omega^\omega$ and a Π_1^0 set C such that $f|_C$ is one-one and $f(C) = B$. Consider the predicate $Q(s) \leftrightarrow f(C \cap N_s) \notin \mathcal{F}$. Clearly, Q is Π_1^1 . Fix a Π_1^1 norm on Q and define

$P(s, t) \leftrightarrow Q(s) \ \& \ Q(t) \ \& \ s \subseteq t \ \& \ \text{lh}(t) = \text{lh}(s) + 1 \ \&$

$(\forall t') \{s \subseteq t' \ \& \ \text{lh}(t') = \text{lh}(s) + 1$

$\rightarrow (t \leq_\varphi^* t' \ \& \ [t <_\varphi^* t' \vee (t)_{\text{lh}(t)-1} \leq (t')_{\text{lh}(t')-1}])\}$.

Plainly P is Π_1^1 and since \mathcal{F} is a σ -ideal for each $s \in Q$ there is a (unique) $t \in Q$ such that $P(s, t)$. By the Recursion Theorem ([38; 7A.2]) the unique α that P defines is a Δ_1^1 point in C . Hence $f(\alpha)$ is a Δ_1^1 point in A .

Remark 1 : It is clear from the proof of (1) that if Γ is a Spector pointclass and \mathcal{F} a hereditary family of subsets of ω^ω which satisfies the Γ computational formula, then for any Γ set $A \notin \mathcal{F}$ there is a $\Delta (= \Gamma \cap \neg\Gamma)$ set $B \subseteq A$ such that $B \notin \mathcal{F}$.

Remark 2 : An appropriate generalization of (3) to higher levels of the analytical hierarchy has been obtained by Srivatsa.

Remark 3 : The following example of Srivatsa shows that one cannot hope to get the basis theorem for hereditary families.

Example : Let $A \subseteq \omega^\omega$ be a Δ_1^1 set not containing any Δ_1^1 real. Define $I(A) = \{B \subseteq \omega^\omega : A \not\subseteq B\}$. Clearly $I(A)$ is a hereditary family and it is easy to see that it satisfies the Π_1^1 computational formula. Further $A \notin I(A)$ and A contains no Δ_1^1 real.

Remark 4 : A boldface analogue of (1) above yields the following (cf [35]).

Let \mathcal{F} be a σ -ideal which satisfies the Π_1^1 computational formula.

Let $C \subseteq \omega^\omega \times \omega^\omega$ be a Π_1^1 set such that for each α , $C_\alpha = \{\beta : C(\alpha, \beta)\} \notin \mathcal{F}$.

Then there is a Borel measurable function f such that $(\alpha, f(\alpha)) \in C$ for each $\alpha \in \omega^\omega$.

To see this first obtain a Borel set $B \subseteq C$ such that $B_\alpha \notin \mathcal{F}$ for each α , by (dualizing) Piatkiewicz's theorem. Then proceed as in the proof of (5).

We now obtain an analogue of a result of Kechris on the existence of the largest Π_1^1 set in \mathcal{F} where \mathcal{F} is a σ -ideal. Predictably, our proof relies on the techniques developed in [25].

1.5 Theorem : Let the σ -ideal \mathcal{F} satisfy both the Π_1^1 and Σ_1^1 computational formulas. Then there is a largest Π_1^1 set in \mathcal{F} .

Proof : Let $W \subseteq \omega \times \omega^\omega$ be a Π_1^1 set which is ω -universal for Π_1^1 sets of reals. Fix a regular Π_1^1 norm ψ on W . Define

$$Q(n, \alpha) \leftrightarrow W(n, \alpha) \ \& \ \{ \beta : (n, \beta) \leq_{\psi}^* (n, \alpha) \} \in \mathcal{F}.$$

Put $R(\alpha) \leftrightarrow (\exists n) Q(n, \alpha)$. By hypothesis, Q is Π_1^1 and so R is Π_1^1 .

It is easy to check that if $W_n = \{ \alpha : W(n, \alpha) \} \in \mathcal{F}$ then $W_n \subseteq R$. To complete the proof it remains to show that $R \in \mathcal{F}$. For this it suffices to show $Q_n \in \mathcal{F}$ for each n . Suppose there is an $n \in \omega$ such that $Q_n \notin \mathcal{F}$.

Put

$$\alpha \in Q'_n \leftrightarrow \alpha \in Q_n \ \& \ \omega_1^\alpha = \omega_1^{ck}.$$

By Theorem 1, $Q_n \setminus Q'_n \in \mathcal{F}$ and hence $Q'_n \notin \mathcal{F}$. For each ξ , let

$$Q_{n, \xi} = \{ \alpha \in Q_n : \psi(n, \alpha) = \xi \}.$$

Suppose $\alpha \in Q'_n$ and let $\psi(n, \alpha) = \mu$. Then the initial segment of \leq_{ψ}^* determined by (n, α) is a $\Delta_1^1(\alpha)$ prewellordering of $\omega \times \omega^\omega$ of length μ . Hence μ is the length of a $\Delta_1^1(\alpha)$ prewellordering of ω^ω and so $\mu < \omega_1^\alpha = \omega_1^{ck}$ (cf [24]). Thus

$$Q'_n \subseteq \bigcup_{\xi < \omega_1^{ck}} Q_{n, \xi},$$

and, consequently, there is a $\xi < \omega_1^{ck}$ such that $Q_{n, \xi} \notin \mathcal{F}$. But this is clearly a contradiction, since for each ξ , $Q_{n, \xi}$ is easily seen to be in \mathcal{F} . This completes the proof.

Obviously the problem that is of interest now is to find necessary conditions for σ -ideals \mathcal{F} to satisfy the Π_1^1 computational formula.

No satisfactory answer has been obtained. However, we shall show that if \mathcal{F} is generated by an ideal of closed sets which is Π_1^1 coded, then every Σ_1^1 set in \mathcal{F} has the covering property and hence, the Σ_1^1 computational formula holds. The covering theorem incidentally gives alternative proofs of some well-known results like the covering property of Σ_1^1 meager sets, of Σ_1^1 σ -bounded sets and Harrison's theorem on countable Σ_1^1 sets. Louveau ([30], [31]) has already obtained these earlier but our methods are different. Results in the rest of this section were obtained jointly with A. Maitra and V.V. Srivatsa.

1.6 Definition : A family \mathcal{G} of closed subsets of ω^ω will be called an ideal of closed sets if whenever $A \in \mathcal{G}$ and $B \subseteq A$ is closed, then $B \in \mathcal{G}$. \mathcal{F} will be called the σ -ideal generated by \mathcal{G} provided $B \in \mathcal{F}$ iff there is a sequence $\{B_n\}$ in \mathcal{G} such that $B \subseteq \bigcup_n B_n$.

An ideal of closed sets \mathcal{G} is said to be Π_1^1 -coded if the set of codes

$$C = \{ \alpha \in \omega^\omega : \alpha \text{ codes a tree } T \text{ on } \omega \text{ such that } [T] \in \mathcal{G} \}$$

is Π_1^1 . If \mathcal{F} is the σ -ideal generated by a Π_1^1 -coded ideal, then \mathcal{F} is also said to be Π_1^1 -coded.

1.7 Lemma : Let \mathcal{G} be a Π_1^1 coded ideal of closed sets. If $A \in \mathcal{G}$, A is Σ_1^1 , then there is a Δ_1^1 tree T on ω such that $[T] \in \mathcal{G}$ and $A \subseteq [T]$.

Proof : Suppose not, i.e. whenever T is a Δ_1^1 tree such that $A \subseteq [T]$, $[T] \notin \mathcal{G}$. In particular, T' is not Δ_1^1 , where $T' = \{ s \in \text{Seq} : N_s \cap A \neq \emptyset \}$

so that T' is a Σ_1^1 tree and $[T'] \in \mathcal{G}$. Let φ be a Π_1^1 norm on $\text{Seq} \setminus T'$. Define $\psi(t) = \min \{ \varphi(u) : \text{Seq}(u) \ \& \ u \subseteq t \ \& \ \neg T'(u) \}$, $t \in \text{Seq} \setminus T'$. Clearly $u, t \in \text{Seq} \setminus T' \ \& \ u \subseteq t \rightarrow \psi(t) \leq \psi(u)$. It is not hard to check that ψ is a Π_1^1 norm on $\text{Seq} \setminus T'$. Then

$$s \in T' \iff \text{Seq}(s) \ \& \ (\forall \alpha) [(\alpha \text{ codes a tree } \& \dots(i) \\ (\forall t) (\alpha(t) = 0 \rightarrow \neg(t <_{\psi}^* s))] \rightarrow \mathcal{C}(\alpha)].$$

To see this, let $s \in T'$. Now if α codes a tree T'' such that $(\forall t) (t \in T'' \rightarrow \neg(t <_{\psi}^* s))$, then $T'' \subseteq T'$ so that $[T''] \in \mathcal{G}$ and so $\alpha \in \mathcal{C}$. Conversely, suppose $s \notin T' \ \& \ \text{Seq}(s)$. Let

$$S = \{ t \in \text{Seq} : \neg(t <_{\psi}^* s) \}.$$

Clearly S is Δ_1^1 . We define

$$S^* = \{ u \in \text{Seq} : (\exists t)(s(t) \ \& \ u \subseteq t) \}.$$

Then S^* is a Δ_1^1 tree and $T' \subseteq S \subseteq S^*$. Choose $\alpha \in \omega^\omega$ such that $\alpha = 0$ on S^* and $\alpha \neq 0$ off S^* . Then $(\forall t) [\alpha(t) = 0 \rightarrow \neg(t <_{\psi}^* s)]$. Since α codes a Δ_1^1 tree S^* and $A \subseteq [S^*]$, so by assumption $\neg \mathcal{C}(\alpha)$. This proves (i). But (i) is a Π_1^1 definition of T' , so T' is Δ_1^1 , which is a contradiction.

The above lemma can also be formulated in terms of a Spector point-class closed under \forall^{ω^ω} and is of independent interest.

We now prove the covering theorem.

1.8 Theorem: Let \mathcal{G} be as above and \mathcal{I} the σ -ideal generated by \mathcal{G} .

Suppose $A \subseteq \omega^\omega$ is Σ_1^1 and $A \in \mathcal{I}$. Then there is a Δ_1^1 set $B \subseteq \omega \times \omega^\omega$

such that for each n , $B_n = \{\alpha : B(n, \alpha)\} \in \mathcal{B}$ and $(\forall \alpha) [A(\alpha) \rightarrow (\exists n) B(n, \alpha)]$.

Proof : Let $M = \cup \{ [T] : T \text{ is a } \Delta_1^1 \text{ tree \& } [T] \in \mathcal{B} \}$. It is easy to check that M is Π_1^1 .

We claim that if $L \neq \emptyset$ is Σ_1^1 and $L \in \mathcal{K}$, then $L \cap M \neq \emptyset$. To see this, let $f : \omega^\omega \rightarrow \omega^\omega$ be a recursive function and $D \subseteq \omega^\omega$ be a Π_1^0 set such that $f(D) = L$. Now $L \subseteq \bigcup_{n \geq 0} L_n$ where each $L_n \in \mathcal{B}$. By Baire Category theorem, there is $s_0 \in \text{Seq}$ and $n_0 \in \omega$ such that $\emptyset \neq f(D \cap N_{s_0}) \subseteq L_{n_0}$. Now $f(D \cap N_{s_0})$ is Σ_1^1 and so $\text{cl}(f(D \cap N_{s_0}))$ is Σ_1^1 and in \mathcal{B} . Hence by Lemma 1 there is a Δ_1^1 tree T such that $[T] \in \mathcal{B}$ and $f(D \cap N_{s_0}) \subseteq [T]$. It follows that $L \cap M \supseteq L \cap [T] \supseteq f(D \cap N_{s_0}) \neq \emptyset$.

Now if $A \not\subseteq M$, then $L' = A \setminus M$ is Σ_1^1 , nonempty and in \mathcal{K} . Hence $L' \cap M \neq \emptyset$ which is impossible. So $A \subseteq M$. Let

$$R(n, \alpha) \leftrightarrow d(n) \downarrow \& C(d(n)) \& (\forall m) (d(n)(\bar{\alpha}(m)) = 0),$$

where $d : \omega \rightarrow \omega^\omega$ is a partial Π_1^1 recursive function which parametrizes $\Delta_1^1 \cap \omega^\omega$. R is Π_1^1 and $A(\alpha) \rightarrow (\exists n) R(n, \alpha)$. Get a Δ_1^1 function $g : \omega^\omega \rightarrow \omega$ such that $(\forall \alpha \in A) R(g(\alpha), \alpha)$.

$$\text{Let } Q_1(n) \leftrightarrow (\exists \alpha \in A) (g(\alpha) = n)$$

$$Q_2(n) \leftrightarrow d(n) \downarrow \& C(d(n)).$$

So Q_1 is Σ_1^1 , Q_2 is Π_1^1 and $Q_1 \subseteq Q_2$. Get a Δ_1^1 Q such that $Q_1 \subseteq Q \subseteq Q_2$. Define

$$B(n, \alpha) \leftrightarrow Q(n) \& (\forall m) (d(n)(\bar{\alpha}(m)) = 0).$$

Clearly, B is Δ_1^1 and does the job.

The Σ_1^1 computational formula now follows immediately.

1.9 Theorem: Let \mathcal{F} be as above. Then \mathcal{F} satisfies the Σ_1^1 computational formula.

Proof: Let $A \subseteq \mathcal{X} \times \omega^\omega$ be a Σ_1^1 set, where $\mathcal{X} = (\omega^\omega)^k \times \omega^\omega$. Let

$$B = \{x : A_x \in \mathcal{F}\}.$$

It suffices to show that B is Π_1^1 . By a relativization of Theorem 1.8 we have

$$B(x) \leftrightarrow (\forall \beta) [A(x, \beta) \rightarrow (\exists n)(d(n, x) \downarrow \& c(d(n, x)) \\ \& (\forall m)(d(n, x)(\bar{\beta}(m)) = 0))],$$

where d is as in [38; 4D.2]. Thus B is Π_1^1 . \parallel

As an illustration, we shall we shall adopt the above methods to obtain the following separation theorem (see [16] and the next section).

Suppose A and B are two Σ_1^1 sets and A can be separated from B by the difference of two Σ_2^0 sets. Then A can be separated from B by the difference of two $\Delta_1^1 \cap \Sigma_2^0$ sets.

To see this consider the following ideal of closed sets

$$\mathcal{I} = \{ [T] : T \text{ is a tree } \& [T] \cap B \text{ is contained in a } \Sigma_2^0 \\ \text{set disjoint from } A \}.$$

$= \{ [T] : T \text{ is a tree } \& [T] \cap B \text{ is contained in a } \\ \Delta_1^1(T), \Sigma_2^0 \text{ set disjoint from } A \}$, by the separation theorem of Louveau [29].

Then \mathcal{S} is Π_1^1 -coded and if \mathcal{I} is the σ -ideal generated by \mathcal{S} , then $A \in \mathcal{I}$. So by the covering theorem, there is a Δ_1^1 -sequence of sets $F_n = \bigcup_m G_{nm}$ such that $A \subseteq \bigcup_n (F_n = \bigcup_m G_{nm})$ and each F_n, G_{nm} is closed and (uniformly) Δ_1^1 and $G_{nm} \cap A = \varnothing, F_n \cap B \subseteq \bigcup_m G_{nm}$. Then

$$A \subseteq \bigcup_n F_n = \bigcup_n \bigcup_m G_{nm} \text{ and } \left(\bigcup_n F_n = \bigcup_{n,m} G_{nm} \right) \cap B = \varnothing.$$

2. Hyperarithmetical Sets of Ambiguous Borel Classes with Applications to Borel Sets in Product Spaces.

2.0 Introduction: In this part of our dissertation we shall digress a bit and undertake a study of the structure of hyperarithmetical sets which are of ambiguous Borel classes. Thus our object of study is mainly $\Delta_1^1 \cap \Delta_{\xi+1}^0$ ($\xi < \omega_1^{ck}$). Borel sets of ambiguous class $\xi+1$ (i.e. belonging to $\Delta_{\xi+1}^0$) are precisely those sets which can be written as the union of alternating differences of a transfinite decreasing sequence of Π_{ξ}^0 sets. Such a representation gives rise to a natural hierarchy (the so-called difference hierarchy) of $\Delta_{\xi+1}^0$. (cf [28; § 37 IV]). Here we shall show that there is an analogous effective hierarchy of $\Delta_1^1 \cap \Delta_{\xi+1}^0$ which coincides with the difference hierarchy of $\Delta_{\xi+1}^0$ restricted to Δ_1^1 .

Although we believe that the study of such hierarchies is not entirely out of place in the context of the effective \mathbb{R} -sets, the content of this section stems mainly from a result of G. Debs [16] on the structure of Borel sets in product spaces which states the following.

Suppose X and Y be two Polish (analytic) spaces and $B \subseteq X \times Y$ a Borel set such that each vertical section is the difference of two G_δ sets. Then $B = C - D$, where C, D are Borel subsets of $X \times Y$ with G_δ sections. In the previous section, we have obtained an effective analogue by showing that a Δ_1^1 set of reals which is the difference of two \prod_2^0 sets is of the form $C - D$, where C and D are both $\Delta_1^1 \cap \prod_2^0$. This immediately gives Debs' result in product spaces. We shall generalize Debs' result by showing the following¹. Suppose A is a Δ_1^1 set of reals and λ, ξ are fixed recursive ordinals. Suppose there is a decreasing transfinite sequence of \prod_ξ^0 sets $\{G_\mu : \mu < \lambda\}$ such that

$$A = \bigcup_{\substack{\mu < \lambda \\ \mu \text{ even}}} (G_\mu - G_{\mu+1}).$$

Then there is a decreasing " Δ_1^1 -sequence" of $\Delta_1^1 \cap \prod_\xi^0$ sets $\{G_\mu^* : \mu < \lambda\}$ such that

$$A = \bigcup_{\substack{\mu < \lambda \\ \mu \text{ even}}} (G_\mu^* - G_{\mu+1}^*).$$

This completely generalizes Debs' result on Borel sets in product spaces and answers a question raised by us in [3]. (Note that when λ is odd G_λ and G_λ^* above are taken to be the empty set). Now observe that

1. After we have obtained this result we were informed by Prof. Louveau that he has obtained this and much more in his paper "Some properties of the Wadge hierarchy of Borel sets" in [32].

by a result of Burgess and Louveau [29; § 2 Theorem 6] (see Theorem 2.2.8 below) there is an effective analogue of the difference hierarchy in $\Delta_1^1 \cap \Delta_{\xi+1}^0$. Our result quoted above shows that this hierarchy coincides with the hierarchy in $\Delta_1^1 \cap \Delta_{\xi+1}^0$ obtained by restricting the difference hierarchy in $\Delta_{\xi+1}^0$ to Δ_1^1 sets. We shall also show (via a separation theorem for Σ_1^1 sets) that if P and Q are two disjoint Σ_1^1 sets in $\omega^\omega \times \omega^\omega$ such that each vertical section P_α is separated from Q_α by a Δ_ξ^0 set, then there is a Borel set B separating P and Q such that each vertical section B_α is Δ_ξ^0 . This is a result of Burgess for $\xi = 2$; our result extends Burgess' result to every level. Analogues of these results for Δ_1^1 clopen sets are also obtained — results due to J.W. Addison, but never published. We include them for the sake of completeness.

As applications we shall, as is usually the case, obtain structural results for Borel sets in product spaces. These include the generalization of Debs' result mentioned earlier and an unpublished result of Addison and Harrington on the non-existence of Borel sets universal for $\Delta_{\xi+1}^0$.

2.1 Coding pair ; the ξ - and ω -topologies : Most of the material in this subsection is taken from Louveau's article [29]. Here, we fix a recursively presented (r.p.) Polish space X (for definition see [38]; otherwise one may think of X as the space 2^ω or ω^ω or $(\omega^\omega)^k \times \omega^\omega$). Given such a space X , by a coding pair we shall mean a pair $(W(X), C(X))$ (or (W, C) , when X is clear from the context) such that

- (i) W , the set of codes, is a Π_1^1 set of integers,

- (ii) C is a Π_1^1 subset of $\omega \times X$ which is universal for non-empty Δ_1^1 subsets of X
- (iii) $\pi(C) = W$, and
- (iv) the relation $n \in W \ \& \ x \notin C_n$ is Π_1^1 .

Observe that such a coding pair exists (of [38]).

A set $A \subseteq X$ is the Δ_1^1 -union of a sequence $\{B_n : n \in \omega\}$, written $A = \bigcup_1^1 B_n$, if $A = \bigcup_n B_n$ and the set θ defined by

$$\theta(n, x) \leftrightarrow x \in B_n$$

is Δ_1^1 . If Γ is a family of sets, we denote by $\bigcup_1^1 \Gamma$ the family of sets obtained by Δ_1^1 -union performed on sequence of sets from Γ .

The Δ_1^1 -recursive hierarchy $(\Sigma_\xi^*, \pi_\xi^*)_{\xi \in \text{On}}$ is defined by the induction

$$\Sigma_0^* = \pi_0^* = \{N_\emptyset : \emptyset \in \text{Seq}\}$$

$$\Sigma_\xi^* = \bigcup_1^1 (\bigcup_{\eta < \xi} \pi_\eta^*)$$

$$\text{and } \pi_\xi^* = \neg \Sigma_\xi^* = \{A \subseteq X : X - A \in \Sigma_\xi^*\}.$$

Clearly, $\Sigma_\xi^* \subseteq \Sigma_\xi^0 \cap \Delta_1^1$; in fact Louveau proves, for $\xi < \omega_1^{\text{ck}}$,

$$\Sigma_\xi^* = \Sigma_\xi^0 \cap \Delta_1^1 \text{ and } \Delta_1^1 = \bigcup_{\xi < \omega_1^{\text{ck}}} \Sigma_\xi^*.$$

For details regarding this hierarchy we refer to [29].

For each $\xi < \omega_1^{\text{ck}}$, the set W_ξ of ξ -codes is defined by

$$n \in W_\xi \iff n \in W \ \& \ C_n \in \pi_\xi^*$$

One can show that W_ξ and $\bigcup_{\eta < \xi} W_\eta$ are Π_1^1 ([29]).

We now consider a family of topologies as follows.

For each ξ , $1 \leq \xi < \omega_1^{ck}$, the topology T_ξ on X is the topology on X generated by $\Sigma_1^1 \cap (\bigcup_{\eta < \xi} \pi_\eta^0)$. T_ω is the topology on X generated by all Σ_1^1 subsets of X . The topology T_ω will be called the ω -topology or the Harrington topology. Clearly T_1 is the usual topology on X and the topologies $\{T_\xi : \xi < \omega_1^{ck}\}$ are increasing, each T_ξ being coarser than T_ω . Terms and notation like " ξ -open", " ξ -closed", " \overline{A}^ξ " ... refer to the corresponding topological notion in the ξ -topology.

We shall need the following results in the sequel. The proofs can be found in [29].

2.1.1 Lemma (Harrington, Louveau) : The space (X, T_ω) is a Baire space i.e. the countable intersection of ω -dense sets is ω -dense.

Consequently, any open subspace of (X, T_ω) is also a Baire space. \square

2.1.2 Lemma (Louveau) : Let $A \subseteq X$ be a Σ_1^1 set. Then \overline{A}^ξ is π_ξ^0 and Σ_1^1 . \square

For each set $H \subseteq X$, let G_H denote the largest ξ -open set such that $H \cap G_H$ is ω -meager. We set $\overline{H}^\xi = X - G_H$. Clearly \overline{H}^ξ is ξ -closed and $H - \overline{H}^\xi$ is ω -meager. In fact we have

2.1.3 Lemma (Louveau) : If H is π_ξ^0 , then $H \Delta \overline{H}^\xi$ is ω -meager.

Proof : See [29].

For any two subsets A, B of X we write $A \underset{a.e.}{\subset} B$ if $A - B$ is ω -meager. Clearly, if $A \underset{a.e.}{\subset} B$ then $A \underset{a.e.}{\subset} B$, $\underset{a.e.}{\subset}$ is transitive and satisfies many of the properties of \subset .

2.1.4 Lemma: If $A \underset{a.e.}{\subset} X$ is ω -open, and $A \underset{a.e.}{\subset} H$, then $A \underset{a.e.}{\subset} \tilde{H}^\xi$.

Proof: Clearly, $A \underset{a.e.}{\subset} \tilde{H}^\xi$ and hence $A - \tilde{H}^\xi$ is ω -meager. But $A - \tilde{H}^\xi$ is ω -open. So by Lemma 2.1.1, $A - \tilde{H}^\xi = \emptyset$.

2.2 The main results: For any Polish space X , a property \mathcal{P} of subsets of X is called a faithful separation (FS) property if for any disjoint $\sum_1^1 A, E \subset X^2$, if there exist a set $B \subset X^2$ separating A, E such that every vertical section B_x of B has \mathcal{P} , then there exists a Borel set with the same properties. Burgess [9] has shown that the property of being Δ_2^0 is FS. In this section we shall extend this result to every level. We first need the following separation principle.

2.2.1 Lemma: Let $1 \leq \xi < \omega_1^{ck}$. Let P and Q be two \sum_1^1 sets of reals which can be separated by a Δ_ξ^0 set. Then they can be separated by a $\Delta_1^1 \cap \Delta_\xi^0$ set.

Proof: Suppose H is a Δ_ξ^0 set such that $P \subset H$ and $H \cap Q = \emptyset$. Since P is \sum_1^1 , by Lemma 2.1.4, $P \subset \tilde{H}^\xi$. Hence $\bar{P}^\xi \subset \tilde{H}^\xi$. Similarly, $\bar{Q}^\xi \subset \tilde{H}^{0\xi}$. So,

$$\bar{P}^\xi \cap \bar{Q}^\xi \subset \tilde{H}^\xi \cap \tilde{H}^{0\xi} \subset (\tilde{H}^\xi - H) \cup (\tilde{H}^{0\xi} - H^0).$$

Since both H and H^0 are π_ξ^0 , by Lemma 2.1.3, both $\tilde{H}^\xi - H$ and $\tilde{H}^{0\xi} - H^0$ are ω -meager. Hence $\bar{P}^\xi \cap \bar{Q}^\xi$ is ω -meager. By Lemma 2.1.2, $\bar{P}^\xi \cap \bar{Q}^\xi$

is Σ_1^1 and hence by Lemma 2.1.1, $\bar{P}^\xi \cap \bar{Q}^\xi = \emptyset$. Applying the separation theorem of Louveau twice one obtains two disjoint Δ_1^1, π_ξ^0 sets P^* and Q^* such that $\bar{P}^\xi \subseteq P^*$ and $\bar{Q}^\xi \subseteq Q^*$. Now by [29; §2A, Corollary 3] or by effectivizing [28; §30 VII Theorem 2] we get a $\Delta_1^1 \cap \Delta_\xi^0$ set separating P^* from Q^* . This completes the proof.

The relativized version of the above immediately gives the FS property of Δ_ξ^0 .

2.2.2 Theorem: Let P and Q be two disjoint Σ_1^1 subsets of $\omega^\omega \times \omega^\omega$ such that each vertical section $P_\alpha = \{\beta : P(\alpha, \beta)\}$ of P is separated from Q_α by a Δ_ξ^0 set, $1 \leq \xi < \omega_1$. Then there is a Borel set B separating P and Q such that each vertical section B_α is Δ_ξ^0 .

Proof: Fix z such that P and Q are $\Sigma_1^1(z)$ and $\xi < \omega_1^z$. Then for each α , P_α and Q_α are $\Sigma_1^1(\langle \alpha, z \rangle)$ sets which can be separated by a Δ_ξ^0 set. Hence by Lemma 2.2.1 (relativized) there is a $\Delta_1^1(\langle \alpha, z \rangle) \cap \Delta_\xi^0$ set separating P_α and Q_α . As in [29; §2C] fix a pair (W, C) of π_1^1 relations, $W \subseteq X \times \omega^\omega$, $C \subseteq X \times \omega \times \omega^\omega$ such that for each x , $\langle W_x, C_x \rangle$ is a coding pair for $\Delta_1^1(x)$ subsets of ω^ω . Consider the following (of [29] for notation)

$$R(\alpha, m, n) \leftrightarrow m \in W_\xi^{\langle \alpha, z \rangle} \ \& \ n \in W_\xi^{\langle \alpha, z \rangle}$$

$$\& \ C_{\langle \alpha, z \rangle, m} = \neg C_{\langle \alpha, z \rangle, n} \ \& \ P_\alpha \subseteq C_{\langle \alpha, z \rangle, m}$$

$$\& \ C_{\langle \alpha, z \rangle, m} \cap Q_\alpha = \emptyset.$$

One can check that R is $\pi_1^1(z)$ (of [29]) and $(\forall \alpha) (\exists m) (\exists n) R(\alpha, m, n)$.

Hence by Easy Uniformization [38] there exist $\Delta_1^1(z)$ functions f and g such that $(\forall \alpha) R(\alpha, f(\alpha), g(\alpha))$. Now define B as follows.

$$B(\alpha, \beta) \leftrightarrow C(\langle \alpha, z \rangle, f(\alpha), \beta)$$

Clearly, B is $\Delta_1^1(z)$ and hence Borel, B separates P from Q and for each α , B_α is Δ_ξ^0 . This completes the proof.

We shall now prove the structural result mentioned in the introduction. We fix for the rest of this section a recursive function φ such that if $\alpha \in \omega^0$ then $\varphi(\alpha, n)$ codes the strict initial segment of \leq_α with top n , whenever $n \in \text{Field}(\leq_\alpha)$; otherwise it codes the empty relation.

2.2.3 Theorem: Let L and M be two disjoint Σ_1^1 sets of reals, $P \subseteq \omega^\omega$ any Σ_1^1 set and $1 \leq \lambda < \omega_1^{\text{ck}}$, $1 \leq \xi < \omega_1^{\text{ck}}$. Suppose there is a decreasing transfinitary sequence $\{H_\mu : \mu < \lambda\}$ of π_ξ^0 sets such that

- (a) $L \subseteq H_0$
- (b) $H_\mu = \bigcap_{\mu' < \mu} H_{\mu'}$, if $\mu < \lambda$ is limit
- (c) $H_\mu \cap M \subseteq H_{\mu+1}$, if $\mu < \lambda$ is even
- (d) $H_\mu \cap L \subseteq H_{\mu+1}$, if $\mu < \lambda$ is odd
- (e) $(\bigcap_{\mu < \lambda} H_\mu) \cap P = \emptyset$

Then, for each recursive $\alpha \in \omega^0$ with $|\alpha| = \lambda$, there is a π_ξ^* set $H^* \subseteq \omega \times \omega^\omega$ such that

- (0) $H_i^* \supseteq H_j^*$ if $i \leq_\alpha j$

- (1) $L \subseteq H_n^*$ if $n \in \text{Field}(\leq_\alpha)$ and $|\varphi(\alpha, n)| = 0$
- (2) $H_i^* = \bigcap_{j <_\alpha i} H_j^*$ if $|\varphi(\alpha, i)|$ is limit
- (3) $H_i^* \cap M \subseteq H_j^*$ if $|\varphi(\alpha, i)|$ is even and $|\varphi(\alpha, i)| + 1 = |\varphi(\alpha, j)|$
- (4) $H_i^* \cap L \subseteq H_j^*$ if $|\varphi(\alpha, i)|$ is odd and $|\varphi(\alpha, i)| + 1 = |\varphi(\alpha, j)|$
- (5) $(\bigcap_{i \in \text{Field}(\leq_\alpha)} H_i^*) \cap P = \varnothing$.

Proof : We shall prove this by induction on λ . So assume that the result is true for all $\lambda' < \lambda$.

Case 1. $\lambda = \bar{\lambda} + 1$.

If $\bar{\lambda}$ is limit, then observe that the sequence $\{H_\mu : \mu < \bar{\lambda}\}$ satisfies the hypotheses (a) - (e) of the theorem. Hence induction hypothesis applies. So assume $\lambda = \lambda_1 + 2$. Fix a recursive function f and π_1^0 set $A \subseteq \omega^\omega$ such that $f(A) = P$. Let (W, C) and (W', C') be coding pairs which code Δ_1^1 subsets of ω^ω and $\omega \times \omega^\omega$ respectively.

Subcase (i). λ_1 is even.

In this case we define A_1^* as follows.

$$\begin{aligned}
 x \in A_1^* &\leftrightarrow (\exists i) \left\{ i \in \bigcup_{\eta < \xi} W_\eta \ \& \ x \in C_i \ \& \ (\exists m \in \bigcup_{\eta < \xi} W_\eta) (\exists n \in W'_\xi) (\forall \beta) \right. \\
 &\quad \left. [\beta \in A \cap C_i \rightarrow (f(\beta) \in C_m \text{ and} \right. \\
 &\quad (0) : (\forall i) (\forall j) (i \in \text{Field}(\leq_\alpha) \ \& \ i <_\alpha j \rightarrow C'_{n,i} \supseteq C'_{n,j}) \\
 &\quad \& \ (1) : (\forall i) (i \in \text{Field}(\leq_\alpha) \ \& \ |\varphi(\alpha, i)| = 0 \rightarrow L \subseteq C'_{n,i}) \\
 &\quad \& \ (2) : (\forall i) (|\varphi(\alpha, i)| \text{ is limit} \rightarrow C'_{n,i} = \bigcap_{j <_\alpha i} C'_{n,j}) \\
 &\quad \& \ (3) : (\forall i) (\forall j) (i \in \text{Field}(\leq_\alpha) \ \& \ |\varphi(\alpha, i)| \text{ is even} \\
 &\quad \quad \& \ |\varphi(\alpha, i)| + 1 = |\varphi(\alpha, j)| \rightarrow C'_{n,i} \cap M \subseteq C'_{n,j}) \\
 &\quad \& \ (4) : (\forall i) (\forall j) (|\varphi(\alpha, i)| \text{ is odd} \ \& \ |\varphi(\alpha, i)| + 1 = |\varphi(\alpha, j)| \\
 &\quad \quad \rightarrow C'_{n,i} \cap L \subseteq C'_{n,j}) \\
 &\quad \& \ (5') : C'_{n,i^*} \cap C_m \cap M = \varnothing \left. \right\},
 \end{aligned}$$

where $|\varphi(\alpha, i^*)| = \lambda_1$. Clearly, A_1^* is ξ -open and hence ω -open. Moreover, it is not hard to check that A_1^* is also π_1^1 . We now claim that $A \subseteq A_1^*$. If not, then $B = A - A_1^* \neq \varnothing$ and so $f(B) \subseteq f(A) = P$. Hence by (e), $f(B) \subseteq F_{\lambda_1+1}$, where $F_{\lambda_1+1} = \omega^\omega - H_{\lambda_1+1}$ is Σ^0_{ξ} . So we may write $F_{\lambda_1+1} = \bigcup_{k \in \omega} F_{\lambda_1+1}^k$, where each $F_{\lambda_1+1}^k$ is in $\pi^0_{\eta(k)}$ for some $\eta(k) < \xi$. Consequently, $f(B) \subseteq \bigcup_{a \in \theta} \bigcup_k \widetilde{F_{\lambda_1+1}^k}^{\eta(k)}$. Now observe that B is ω -clopen and hence (B, T_ω) is a Baire space. Further $f : (B, T_\omega) \rightarrow (\omega^\omega, T_\omega)$ is continuous and open. Hence $B \subseteq \bigcup_{a \in \theta} \bigcup_k f^{-1}(\widetilde{F_{\lambda_1+1}^k}^{\eta(k)})$. So there is a $k \in \omega$ such that $f^{-1}(\widetilde{F_{\lambda_1+1}^k}^{\eta(k)})$ is not ω -nowhere dense.

Hence there is a non-empty Σ_1^1 set $B' \subseteq B$ such that $f(B') \subseteq \widetilde{F}_{\lambda_1+1}^k{}^{\eta(k)}$.

But $\widetilde{F}_{\lambda_1+1}^k{}^{\eta(k)} \subseteq_{a.e.} F_{\lambda_1+1}^k \subseteq F_{\lambda_1+1}$. Hence $\overline{f(B')}^{\eta(k)} \subseteq_{a.e.} F_{\lambda_1+1}$. This shows

that $\overline{f(B')}^{\eta(k)} \cap H_{\lambda_1+1}$ is ω -meager. Since $\overline{f(B')}^{\eta(k)}$ is ξ -open,

we have

$$\overline{f(B')}^{\eta(k)} \cap \widetilde{H}_{\lambda_1+1}^{\xi} = \varnothing. \quad \dots(1)$$

We now define a decreasing sequence of ξ -closed sets $\{J_\mu : \mu \leq \lambda_1 + 1\}$

satisfying the following

- (i) $L \subseteq J_0$
- (ii) $J_\mu \subseteq \widetilde{H}_\mu^{\xi}$ for each $\mu \leq \lambda_1 + 1$
- (iii) each J_μ is Σ_1^1
- (iv) $J_\mu \cap M \subseteq J_{\mu+1}$ if $\mu < \lambda_1 + 1$ is even
- (v) $J_\mu \cap L \subseteq J_{\mu+1}$ if $\mu < \lambda_1 + 1$ is odd
- (vi) $J_\mu = \bigcap_{\mu' < \mu} J_{\mu'}$ if $\mu < \lambda_1 + 1$ is limit.

Such a sequence is defined by recursion as follow. Fix once and for

all $m^*, n^* \in \text{Field}(\leq_\alpha)$ such that $|\varphi(\alpha, m^*)| = 0$ and $|\varphi(\alpha, n^*)| = \lambda_1 + 1$.

Let G be a good ω -universal set for $\Sigma_1^1 \int \omega \times \omega^\omega$. Define

$$R(p, n, \alpha) \leftrightarrow (n = m^* \ \& \ \alpha \in \overline{L^\xi})$$

$$\vee (|\varphi(\alpha, n)| \text{ is limit } \& \ (\forall j <_\alpha n) G(p, j, \alpha))$$

$$\vee (\exists m \in \text{Field}(\leq_\alpha)) (|\varphi(\alpha, m)| + 1 = |\varphi(\alpha, n)| \ \& \ n <_\alpha m^*)$$

$$\& \ |\varphi(\alpha, n)| \text{ is even } \& \ \alpha \in \overline{G_{p, m} \cap L^\xi}$$

$$\vee (\exists m \in \text{Field}(\leq_\alpha)) (|\varphi(\alpha, m)| + 1 = |\varphi(\alpha, n)| \ \& \ n \leq_\alpha m^*)$$

$$\& \ |\varphi(\alpha, n)| \text{ is odd } \& \ \alpha \in \overline{G_{p, m} \cap M^\xi}.$$

As in the proof of Lemma 2.2.1 (of [29]) one can show that R is Σ_1^1 .

Hence by the Recursion Theorem of Kleene ([38; 3H]) there is a $p^* \in \omega$ such that

$$R(p^*, n, \alpha) \leftrightarrow G(p^*, n, \alpha).$$

Then for $\eta \leq \lambda_1 + 1$, we define J_η by

$$\alpha \in J_\eta \leftrightarrow R(p^*, n, \alpha),$$

where $n \in \text{Field}(\leq_\alpha)$ & $|\varphi(\alpha, n)| = \eta$.

Clearly J_η 's satisfy (i), (iii), (vi) above; (ii), (iv) and (v) are proved by induction. To see (ii), for instance, observe that

$$J_{\eta+1} = \overline{J_\eta \cap L^\xi} \text{ when } \eta+1 \text{ is even.}$$

But $J_\eta \cap L \subseteq \tilde{H}_\eta^\xi \cap L$ by induction hypothesis

$$\underset{\text{a.e.}}{\subseteq} H_\eta \cap L \subseteq H_{\eta+1}.$$

Hence by Lemma 2.1.4, $J_\eta \cap L \subseteq \tilde{H}_{\eta+1}^\xi$.

Consequently, $J_{\eta+1} = \overline{J_{\eta} \cap L^{\xi}} \subseteq \tilde{H}_{\eta+1}^{\xi}$. To see (iv), observe that, when η is even,

$$J_{\eta} \cap M \subseteq \overline{J_{\eta} \cap M}^{\xi} = J_{\eta+1}, \text{ by definition.}$$

The other cases can be proved similarly. Having defined such a sequence we have by (I),

$$\overline{f(B')^{\eta(k)}} \cap J_{\lambda_1+1} = \varnothing$$

So by the separation theorem of Louveau [29], there is a $\pi_{\eta(k)}^*$ set E such that

$$f(B') \subseteq E \text{ and } E \cap J_{\lambda_1+1} = \varnothing.$$

Now $E \cap M \subseteq M - J_{\lambda_1+1} \subseteq \omega^{\omega} - J_{\lambda_1}$ by (iv).

Thus the decreasing sequence of π_{ξ}^0 sets $\{J_{\mu} : \mu \leq \lambda_1\}$ satisfies the hypotheses (a) - (e) of the theorem with P replaced by $E \cap M$. Hence by induction hypothesis, there is an $n \in \omega'_{\xi}$ such that C'_n satisfies (0) - (5) in the statement of the theorem with P replaced by $E \cap M$ and α replaced by $\varphi(\alpha, n^*)$.

Now fix $x \in B'$. Get $m \in \bigcup_{\eta < \xi} \omega_{\eta}$ such that $E = C_m$. Since f is recursive, there is an $i \in \bigcup_{\eta < \xi} \omega_{\eta}$ such that $C_i = f^{-1}(C_m)$. Clearly $x \in C_i$. Suppose $\beta \in A \cap C_i$. Then $f(\beta) \in C_m$ and moreover C'_n satisfies conditions (0) - (5') in the definition of A_1^* . This establishes $x \in A_1^*$, which is a contradiction. Thus $A \subseteq A_1^*$. This means that for each $y \in P$ there exist $m \in \bigcup_{\eta < \xi} \omega_{\eta}$, $n \in \omega'_{\xi}$ such that $y \in C_m$ and C'_n

satisfies conditions (0) - (5') in the definition of A_1^* . Define

$$Q(y,m) \leftrightarrow m \in \bigcup_{\eta < \xi} W_\eta \ \& \ y \in C_m$$

$$\& (\exists n) [n \in W'_\xi \ \& \ C'_n \text{ satisfies (0) - (5')}] .$$

Q is Π_1^1 and for each $y \in P$ there is an m such that $Q(y,m)$ holds.

Hence there is a Δ_1^1 -function f such that $(\forall y \in P) R(y, f(y))$.

Let $U = f[P]$ and $V = \{m \in \bigcup_{\eta < \xi} W_\eta : (\exists n) [n \in W'_\xi \ \& \ C'_n \text{ satisfies (0) - (5')}] \}$.

U is Σ_1^1 , V is Π_1^1 and $U \subseteq V$. Let $U^* \subseteq \omega$ be Δ_1^1 such that $U \subseteq U^* \subseteq V$. Clearly, $P \subseteq \bigcup_{m \in U^*} C_m$ and, moreover, for each $m \in U^*$, $(\exists n) [n \in W'_\xi \ \& \ C'_n \text{ satisfies (0) - (5') in the definition of } A_1^*]$.

By Δ -selection again, there is a Δ_1^1 function g which for each $m \in U^*$ picks an integer $n \in W'_\xi$ such that C'_n satisfies (0) - (5') in the definition of A_1^* . Now define

$$H^*(i,x) \leftrightarrow i \notin \text{Field}(\leq_\alpha) \vee ((i <_\alpha n^*) \ \& \ (\forall m \in U^*) C'(g(m), i, x)) \\ \vee ((i = n^*) \ \& \ (\forall m \in U^*) [\neg C(m, x) \ \& \ C'(g(m), i^*, x)]),$$

where $|\varphi(\alpha, i^*)| = \lambda_1$. A tedious but straightforward computation shows that H^* satisfies the conclusion of the theorem.

Subcase (ii) λ_1 is odd.

In this case we define A_2^* exactly like A_1^* except that we replace in the definition of A_1^* condition (5') by

$$(5'') \quad C'_{n, i^*} \cap C_m \cap L = \varnothing ;$$

and the proof proceeds exactly as above.

Case 2 λ is limit.

In this case we define A_3^* as follows.

$$x \in A_3^* \leftrightarrow (\exists i) \left\{ i \in \bigcup_{\eta < \xi} W_\eta \ \& \ x \in C_i \ \& \right.$$

$$\left. (\exists p \in \text{Field}(\leq_\alpha)) (\exists m \in \bigcup_{\eta < \xi} W_\eta) (\exists n \in W'_\xi) (\forall \beta) [\beta \in A \cap C_i \rightarrow \right.$$

$f(\beta) \in C_m \ \& \ (C'_n \text{ satisfies (0) - (4) in the definition$

of A_1^* with $\varphi(\alpha, p)$ replacing α) \ \&

$$\left. (5''') \left(\bigcap_{i <_{\alpha, p} n, i} C'_{n, i} \cap C_m = \varphi \right) \right\}.$$

As in subcase (i) we can show that $A \subseteq A_3^*$, and further U^* is defined similarly such that for each $m \in U^*$, there exist $p \in \text{Field}(\leq_\alpha)$ and $n \in W'_\xi$ such that C'_n satisfies (0) - (5''') above with α replaced by $\varphi(\alpha, p)$. Get two Δ_1^1 functions f and g which for each $m \in U^*$ select a p and an n as above. Then define

$$x \in H'_i \leftrightarrow i \notin \text{Field}(\leq_\alpha) \vee (\forall m \in U^*) [i \leq_\alpha f(m) \rightarrow C'(g(m), i, x)].$$

Finally define

$$x \in H_i^* \leftrightarrow (\forall j \leq_\alpha i) H'(j, x) \vee [(|\varphi(\alpha, i)| \text{ is limit}) \ \& \ (\forall j <_\alpha i) H'(j, x)].$$

One can now check that H^* is our required set. This completes the proof.

As an immediate consequence of Theorem 2.2.3 we have the following separation result.

2.2.4 Theorem : Let L and M be two disjoint Σ_1^1 sets of reals. Let λ and ξ be two recursive ordinals. Suppose there is a decreasing sequence of π_{ξ}^0 sets $\{H_{\mu} : \mu < \lambda\}$ such that, at limit ordinals μ , $H_{\mu} = \bigcap_{\mu' < \mu} H_{\mu'}$, and the set $B = \bigcup_{\mu < \lambda, \mu \text{ even}} (H_{\mu} - H_{\mu+1})$ separates L from M . Then, whenever α is a recursive code of λ , there is a π_{ξ}^* set $H^* \subseteq \omega \times \omega^{\omega}$ such that

$$(a) \quad H_n^* \subseteq H_m^* \text{ if } m \leq_{\alpha} n.$$

$$(b) \quad H_n^* = \bigcap_{n' <_{\alpha} n} H_{n'}^*, \text{ if } |\varphi(\alpha, n)| \text{ is limit}$$

$$\text{and } (c) \quad B^* = \bigcup \{H_n^* - H_m^* : |\varphi(\alpha, n)| \text{ is even, } n \in \text{Field}(\leq_{\alpha}) \\ \text{and } |\varphi(\alpha, n)| + 1 = |\varphi(\alpha, m)|\}$$

separates L from M (when λ is odd we take $H_{\lambda} = \emptyset$. Similarly for H^*).

Proof : This follows from Theorem 2.2.3 by taking $P = L$ when λ is even and $P = M$ when λ is odd. \square

We now state the main result of this section.

2.2.5 Theorem : Let $A \subseteq \omega^{\omega}$ be a Δ_1^1 set and $\{H_{\mu} : \mu < \lambda\}$ a decreasing sequence of π_{ξ}^0 sets as above such that

$$A = \bigcup_{\mu < \lambda, \mu \text{ even}} (H_{\mu} - H_{\mu+1}).$$

Then for each recursive α coding λ , there is a π_{ξ}^* set $H^* \subseteq \omega \times \omega^{\omega}$ as in Theorem 2.2.4 such that

$$A = \bigcup \{H_n^* - H_m^* : |\varphi(\alpha, n)| \text{ is even, } n \leq_{\alpha} m \text{ \& } |\varphi(\alpha, n)| + 1 = |\varphi(\alpha, m)|\},$$

Proof : Follows from Theorem 2.2.4 by taking $L = A$ and $M = A^0$. \square

To obtain structural properties of Δ_1^1 sets in the ambiguous Borel class $\xi+1$, $\xi < \omega_1^{ck}$, we shall first obtain a structural result for $\Delta_1^1 \cap \Delta_2^0$ sets. This result is due to Burgess and since no published account exists we shall give a proof for the sake of completeness. We need the following result from inductive definability due to Cenzer [15]. (For basic facts about inductive operators we refer to [38]).

2.2.6 Lemma : Let Γ be a monotone Π_1^1 operator on ω^ω and E a Σ_1^1 set of reals such that $E \subseteq \Gamma^\omega$. Then there is a recursive ordinal λ such that $E \subseteq \Gamma^\lambda$.

The above lemma enables us to effectivize [28; §34 VI].

2.2.7 Theorem : Let $E \subseteq \omega^\omega$ be a $\Delta_1^1 \cap \Delta_2^0$ set. Then there is a recursive ordinal λ such that for every recursive code α of λ there is Δ_1^1 , closed set $F \subseteq \omega \times \omega^\omega$ with the following properties

$$(a) F_n \subseteq F_m \text{ if } m <_\alpha n$$

$$(b) F_n = \bigcap_{n' <_\alpha n} F_{n'} \text{ if } |\varphi(\alpha, n)| \text{ is limit}$$

$$\text{and } (c) E = \bigcup \{ F_n - F_m : |\varphi(\alpha, n)| \text{ is even \& } n <_\alpha m \\ \& |\varphi(\alpha, n)| + 1 = |\varphi(\alpha, m)| \}.$$

Proof : Define a transfinite decreasing sequence of closed sets by the induction

$$X_0 = \omega^\omega, \quad X_\mu = \bigcap_{\mu' < \mu} X_{\mu'} \text{ if } \mu \text{ is limit,}$$

$$X_{\mu+1} = \overline{X_\mu \cap E \cap (X_\mu - E)}.$$

Since X_μ is a decreasing sequence of closed sets there is an ordinal $\mu^* < \omega_1$ such that $X_{\mu^*+1} = X_{\mu^*}$. Thus we have

$$X_{\mu^*} = \overline{X_{\mu^*} \cap E} \cap \overline{X_{\mu^*} - E}$$

This implies

$$\overline{X_{\mu^*} \cap E} = X_{\mu^*} = \overline{X_{\mu^*} - E}.$$

Now, X_{μ^*} is closed and hence a Baire space. Moreover, both E and E^c are dense G_δ sets in X_{μ^*} . This is not possible unless $X_{\mu^*} = \emptyset$. Set $Y_\mu = \omega^\omega - X_\mu$. Then $Y_\omega = \bigcup_{\mu} Y_\mu = \omega^\omega$. We shall now define a \prod_1^1 monotone operator Γ on ω^ω whose μ th iterate $\Gamma^\mu = Y_\mu$ if $\mu < \omega$ and $\Gamma^\mu = Y_{\mu+1}$ if $\mu \geq \omega$. The operator Γ is defined as follows, where A varies over subsets of ω^ω .

$$\alpha \in \Gamma(A) \iff (\exists \varepsilon) \left\{ (\alpha \in N_\varepsilon) \ \& \ [(\forall \beta) (\beta \in N_\varepsilon \cap E \rightarrow \beta \in A) \right. \\ \left. \vee (\forall \beta) (\beta \in N_\varepsilon \cap E^c \rightarrow \beta \in A) \right\}.$$

This is a monotone \prod_1^1 operator such that $\Gamma^\omega = Y_\omega = \omega^\omega$. Hence by Lemma 2.2.6, there is a recursive ordinal λ such that $\Gamma^\lambda = \omega^\omega$. This implies that, for a recursive ordinal λ' , $X_{\lambda'} = \emptyset$. Now let $x \in E$ and let $\mu < \lambda'$ be the least ordinal such that $x \notin X_{\mu+1}$. Then $x \in X_\mu - \overline{X_\mu - E}$.

Thus

$$E \subseteq \bigcup_{\mu < \lambda'} (X_\mu - \overline{X_\mu - E})$$

Further $\bigcup_{\mu < \lambda'} (X_\mu - \overline{X_\mu - E}) \subseteq E$. Hence

$$E = \bigcup_{\mu < \lambda'} (X_\mu - \overline{X_\mu - E}).$$

An application of Theorem 2.2.5 (with $\xi = 1$) gives us the result. \square

To get the general result we need a transfer theorem (see [29]).

Transfer Lemma (Kuratowski-Louveau) : For each A in $\Delta_{\xi+1}^0 \cap \Delta_1^1$, $2 \leq \xi < \omega_1^{ck}$, there is a Δ_1^1 closed set $F \subseteq \omega^\omega$, and a Δ_1^1 -recursive function f from F onto ω^ω which is injective and continuous, such that for each open subset G of ω^ω , $f(G) \in \Sigma_\xi^0$, and such that $f^{-1}(A)$ is both Δ_2^0 and Δ_1^1 in ω^ω .

2.2.8 Theorem(Louveau) : For each $E \subseteq \omega^\omega$ in $\Delta_1^1 \cap \Delta_{\xi+1}^0$, $1 \leq \xi < \omega_1^{ck}$, there is a recursive ordinal λ and a $\Delta_1^1 \cap \Pi_\xi^0$ set $H \subseteq \omega \times \omega^\omega$ (corresponding to each recursive code α for λ) such that

- (a) $H_n \subseteq H_m$ if $m \leq_\alpha n$
- (b) $H_n = \bigcap_{n' <_\alpha n} H_{n'}$ if $|\varphi(\alpha, n)|$ is limit, and
- (c) $E = \bigcup \{ H_n - H_m : |\varphi(\alpha, n)| \text{ is even} \ \& \ n <_\alpha m \ \& \ |\varphi(\alpha, n)| + 1 = |\varphi(\alpha, m)| \}$.

Proof : This follows from Theorem 2.2.7 and the Transfer Lemma.

Remark : Theorem 2.2.8 enables us to define an effective hierarchy in $\Delta_1^1 \cap \Delta_{\xi+1}^0$ analogous to the difference hierarchy in $\Delta_{\xi+1}^0$.

2.2.9 Corollary : Suppose $E, E' \subseteq \omega^\omega$ be two Σ_1^1 sets which can be separated by a $\Delta_{\xi+1}^0$ set, $1 \leq \xi < \omega_1^{ck}$. Then there is a recursive ordinal λ and a Π_ξ^* set $H \subseteq \omega \times \omega^\omega$ (corresponding to each recursive α for λ), as in Theorem 2.2.8, such that the set

$B = \bigcup \{H_n - H_m : |\varphi(\alpha, n)| \text{ is even} \ \& \ n <_\alpha m \ \& \ |\varphi(\alpha, n)| + 1 = |\varphi(\alpha, m)|\}$
separates E from E' .

Proof : By Lemma 2.2.1 there is a $\Delta_1^1 \cap \Sigma_{\xi+1}^0$ set B separating E from E' . The result now follows from Theorem 2.2.8.

sub-

2.3 The Kalmar hierarchy : In this section we shall study the Kalmar hierarchy for clopen subsets of ω^ω which are Δ_1^1 .

For any set $K \subseteq \omega^\omega$ and any $s \in \omega^{<\omega}$ let

$$K_{(s)} = \{\alpha \in \omega^\omega : s \hat{\ } \alpha \in K\}.$$

The Kalmar hierarchy $\{\mathcal{K}_\xi : \xi < \omega_1\}$ is defined by induction as follows :

$$\mathcal{K}_0 = \{\emptyset, \omega^\omega\}.$$

Having defined \mathcal{K}_η for every $\eta < \xi$, put

$$K \in \mathcal{K}_\xi \iff (\forall \eta) \left[K_{((\eta))} \in \bigcup_{\eta < \xi} \mathcal{K}_\eta \right].$$

It is not very difficult to see that

$$\bigcup_{\xi < \omega_1} \mathcal{K}_\xi = \Delta_1^1 \cap \omega^\omega.$$

Given a clopen set $\emptyset \neq K \subseteq \omega^\omega$ we shall associate a well-founded tree T_K as follows.

$$T_K = \{ \langle s \rangle \in \text{Seq} \mid K_{(s)} \neq \emptyset, \omega^\omega \vee K_{(s)} = \omega^\omega \ \& \ K_{(s')} \neq \omega^\omega \},$$

where $s' = (n_0, \dots, n_{\lambda-1})$ if $s = (n_0, \dots, n_{\lambda-1}, n_\lambda)$. If $K = \emptyset$, T_K is taken to be the empty tree. Conversely, given a well-founded tree T ,

there is a clopen set $K \subseteq \omega^\omega$ such that $T = T_K$. If T is a well-founded tree, then its length, denoted by $|T|$, is defined as in [38].

2.3.1 Proposition: Let $K \subseteq \omega^\omega$ be a clopen set. Then

$$|T_K| \leq \xi < \omega_1 \text{ iff } K \in \mathcal{K}_\xi.$$

Proof: This is proved by induction on ξ and is straightforward.

2.3.2 Proposition: If $K \subseteq \omega^\omega$ is Δ_1^1 and clopen, then there is a recursive $\xi < \omega_1^{ck}$ such that $K \in \mathcal{K}_\xi$.

Proof: Let K be Δ_1^1 and clopen. We shall first show that T_K is a Δ_1^1 well-founded tree. Consider the set

$$T'_K = \{ \langle s \rangle \in \text{Seq} : K_{(s)} = \omega^\omega \text{ \& } K_{(s')} \neq \omega^\omega \}.$$

$$\begin{aligned} \text{Then } s \in T'_K &\leftrightarrow N_s \subseteq K \text{ \& } N_{s'} \not\subseteq K \\ &\leftrightarrow N_s \subseteq K \text{ \& } (\exists t) (s' \subseteq t \text{ \& } N_t \subseteq K^c) \end{aligned}$$

This shows that T'_K is Π_1^1 . Similarly, T'_{K^c} is also Π_1^1 . Further

$$\begin{aligned} s \in T'_K &\leftrightarrow N_s \subseteq K \text{ \& } N_{s'} \not\subseteq K \\ &\leftrightarrow (\forall t) (s \subseteq t \rightarrow t \notin T'_{K^c}) \text{ \& } N_{s'} \not\subseteq K \end{aligned}$$

Thus, T'_K is also Σ_1^1 . Since

$$s \in T_K \leftrightarrow (\exists t) (s \subseteq t \text{ \& } t \in T'_K)$$

it follows that T_K is Δ_1^1 . Hence T_K is a Δ_1^1 well-founded tree and so $\xi = |T_K| < \omega_1^{ck}$. By Proposition 2.3.1, $K \in \mathcal{K}_\xi$. This completes the proof.

We shall now obtain a separation result analogous to Theorem 2.2.4. If α codes a well-founded tree T , we shall denote by $\|\alpha\|$ the length of the tree T . As in [44; 7.10] we can define a Π_1^1 relation $S_<$ such that if β codes a well-founded tree, then

$$S_<(\alpha, \beta) \leftrightarrow \alpha \text{ codes a well-founded tree } \& \|\alpha\| < \|\beta\|.$$

2.3.3 Theorem: Let P, Q be two disjoint Σ_1^1 sets of reals and $\lambda < \omega_1^{\text{ck}}$. Suppose there is a set $K \in \mathcal{K}_\lambda$ which separates P from Q . Then there is a Δ_1^1 set $K^* \in \mathcal{K}_\lambda$ which separates P from Q .

Proof: We shall prove this by induction on λ . Now for each n , $P_{((n))}$ and $Q_{((n))}$ are Σ_1^1 sets and $K_{((n))} \in \bigcup_{\xi < \lambda} \mathcal{K}_\xi$ separates $P_{((n))}$ from $Q_{((n))}$. Hence by induction hypothesis, there is a Δ_1^1 real α , $\|\alpha\| < \lambda$, which codes a wellfounded tree with the associated clopen set separating $P_{((n))}$ and $Q_{((n))}$. Fix a Δ_1^1 γ such that γ codes a Δ_1^1 wellfounded tree and $\|\gamma\| = \lambda$ and define

$$\begin{aligned} R(n, \alpha) \leftrightarrow S_<(\alpha, \gamma) \& (\forall \beta) [\beta \in P_{((n))} \rightarrow (\exists s) (\forall k) (\alpha(s) = 0 \\ & \& \alpha(s^{\langle k \rangle}) \neq 0 \& \beta \in N_s)] \\ & \& (\forall \beta) [(\forall s) (\exists k) (\alpha(s) = 0 \& \alpha(s^{\langle k \rangle}) \neq 0 \& \beta \in N_s) \\ & \rightarrow \beta \notin Q_{((n))}]. \end{aligned}$$

R is Π_1^1 and $(\forall n) (\exists \alpha \in \Delta_1^1) R(n, \alpha)$. Hence by the Uniformization Lemma there is a Δ_1^1 -function f such that $\forall n R(n, f(n))$. Now define a tree T as follows:

$$T = \{ \langle n \rangle s \mid f(n)(s) = 0 \}.$$

T is wellfounded and $|T| \leq \lambda$. If K^* be the clopen set such that $T = T_{K^*}$, then it is not hard to see that K^* is Δ_1^1 and separates P from Q .

An immediate consequence of Lemma 2.2.1 and Proposition 2.3.2 is the following

2.3.4 Theorem: Let P and Q be two Σ_1^1 sets of reals and suppose there is a Δ_1^0 set which separates P and Q . Then there is a $\lambda < \omega_1^{ck}$ and a Δ_1^1 set $K \in \mathcal{H}_\lambda$ such that K separates P from Q . \square

2.4 Applications to Borel sets in product spaces: In this subsection we shall obtain consequences of the results obtained in the previous subsections to the boldface theory. As usual for convenience we shall work in ω^ω . Our first result is the generalization of Debs' result mentioned in the introduction.

2.4.1 Theorem: Let B be a Borel subset of $\omega^\omega \times \omega^\omega$ and λ, ξ two countable ordinals. Suppose for each $x \in \omega^\omega$ there is a decreasing transfinite sequence of π_ξ^0 sets $\{H_\mu : \mu < \lambda\}$, with $H_\mu = \bigcap_{\mu' < \mu} H_{\mu'}$ at limits, such that the vertical section $B_x = \bigcup_{\mu < \lambda, \mu \text{ even}} (H_\mu - H_{\mu+1})$.

Then there is a decreasing transfinite sequence of Borel sets $\{B^\mu : \mu < \lambda\}$ with sections in π_ξ^0 such that

$$B^\mu = \bigcap_{\mu' < \mu} B^{\mu'} \quad \text{if } \mu \text{ is limit}$$

$$\text{and } B = \bigcup_{\mu < \lambda, \mu \text{ even}} (B^\mu - B^{\mu+1}).$$

Proof: Fix z such that B is $\Delta_1^1(z)$ and $\xi, \lambda < \omega_1^z$. Then each vertical section B_x is $\Delta_1^1(\langle x, z \rangle)$ and is also the union of alternating differences of a decreasing λ -sequence of π_ξ^0 sets. Hence by Theorem 2.2.5 (relativized), there is a π_ξ^* ($\langle x, z \rangle$) set $H^* \subseteq \omega \times \omega^\omega$ such that

$$(a) H_n^* \subseteq H_m^* \text{ if } m \leq_\alpha n$$

$$(b) H_n^* = \bigcap_{n' <_\alpha n} H_{n'}^* \text{ if } |\varphi(\alpha, n)| \text{ is limit}$$

$$\text{and } (c) B_x = \bigcup \{ H_n^* - H_m^* : |\varphi(\alpha, n)| \text{ is even \& } n <_\alpha m \\ \& |\varphi(\alpha, n)| + 1 = |\varphi(\alpha, m)| \},$$

where α is a recursive in z code for λ . Just as in the proof of Theorem 2.2.2, get a $\Delta_1^1(z)$ -recursive function f such that for each x , $f(x) \in \omega_\xi^{\langle x, z \rangle}$ and $C_{\langle x, z \rangle}, f(x) \stackrel{\text{def}}{=} H^*$ satisfies (a) - (c) above.

(Here for each y , $\langle \omega_y, C_y \rangle$ is a coding pair for $\Delta_1^1(y) / \omega \times \omega^\omega$).

For each $\mu < \lambda$ fix $n^* \in \text{Field}(\leq_\alpha)$ such that $|\varphi(\alpha, n^*)| = \mu$ and define

$$B^\mu(x, y) \leftrightarrow C(\langle x, z \rangle, f(x), n^*, y).$$

It can be checked that $\{ B^\mu : \mu < \lambda \}$ satisfies the conclusion of the theorem.

2.4.2 Theorem: Let $P, Q \subseteq \omega^\omega \times \omega^\omega$ be two Σ_1^1 sets such that for each x , P_x is separated from Q_x by a $\Delta_{\xi+1}^0$ set, $1 \leq \xi < \omega_1$. Then there is a countable ordinal λ and a decreasing transfinite sequence of Borel sets $\{ B^\mu : \mu < \lambda \}$ as above such that the set

$$B = \bigcup_{\mu < \lambda, \mu \text{ even}} (B_\mu - B_{\mu+1})$$

separates P from Q .

Proof : By Theorem 2.2.2 there is a Borel set $B \subseteq \omega^\omega \times \omega^\omega$ with sections in $\Delta_{\xi+1}^0$ which separates P from Q . Fix z such that B is $\Delta_1^1(z)$ and $\xi < \omega_1^Z$. Thus for each x , B_x is $\Delta_1^1(\langle x, z \rangle)$ and $\Delta_{\xi+1}^0$. Hence by Theorem 2.2.8, relativized, there is an $\alpha \in \omega^\omega$ recursive in $\langle x, z \rangle$ and a $\pi_\xi^*(\langle x, z \rangle)$ set $H^* \subseteq \omega \times \omega^\omega$ as in the proof of Theorem 2.4.1 above. By the Uniformization Lemma and the boundedness theorem there is a $\lambda < \omega_1$ such that each $\alpha \in \omega^\omega$ above has length less than λ . An application of Theorem 2.4.1 gives a sequence $\{B^\mu : \mu < \lambda\}$ of Borel sets with required properties such that $B = \bigcup_{\mu < \lambda, \mu \text{ even}} (B^\mu - B^{\mu+1})$.

Remark : The above result for $\xi = 1$ is due to Burgess [9].

The following is now immediate.

2.4.3 Corollary : Let B be a Borel subset of $\omega^\omega \times \omega^\omega$ with each vertical section in $\Delta_{\xi+1}^0$, $1 \leq \xi < \omega_1$. Then there is a countable ordinal λ and a decreasing sequence of Borel sets $\{B^\mu : \mu < \lambda\}$ as above such that $B = \bigcup_{\mu < \lambda, \mu \text{ even}} (B^\mu - B^{\mu+1})$.

Finally we establish the following unpublished result of Addison and Harrington.

2.4.4 Corollary : There is no Borel set $B \subseteq \omega^\omega \times \omega^\omega$ which is universal for $\Delta_{\xi+1}^0 \upharpoonright \omega^\omega$, $\xi < \omega_1$.

Proof : Suppose there is such a Borel set B . Fix z such that B is $\Delta_1^1(z)$ and assume $\xi \geq 1$. Then for each x , B_x is $\Delta_1^1(\langle x, z \rangle)$ and $\Delta_{\xi+1}^0$. Hence by Theorem 2.2.8, relativized, there is an $\alpha \in \omega^\omega$ recursive

in $\langle x, z \rangle$ and a $\pi_{\xi}^*(\langle x, z \rangle)$ set $H \subseteq \omega \times \omega^{\omega}$ such that

$$(i) \quad H_n \subseteq H_m \quad \text{if } m \leq_{\alpha} n$$

$$(ii) \quad H_n = \bigcap_{n' <_{\alpha} n} H_{n'} \quad \text{if } |\varphi(\alpha, n)| \text{ is limit}$$

$$\text{and (iii) } B_x = \bigcup \left\{ H_n - H_m : |\varphi(\alpha, n)| \text{ is even \& } n <_{\alpha} m \right. \\ \left. \& |\varphi(\alpha, n)| + 1 = |\varphi(\alpha, m)| \right\}.$$

By the usual selection and boundedness argument there is $\lambda < \omega_1$ such that each $\alpha \in \omega_0$ as above has length less than λ . This would imply that each set in $\Delta_{\xi+1}^0$ is available at the λ th level of the difference hierarchy of $\Delta_{\xi+1}^0$. This is clearly a contradiction (of [28, §37 IV]). The proof for $\xi = 0$ is similar. \square

APPENDIX

Proof of Lemma 7.5 of Part I:

With each operation $R_\rho, \rho < \omega_1$, every family $\{E_n\}$ of subsets of X and every $x \in X$ we will associate a (closed) game on the integers, $G_\rho(x)$, played between \forall and \exists such that

$$x \in R_\rho(\{E_n\}) \leftrightarrow \exists \text{ wins } G_\rho(x).$$

(Observe that $G_\rho(x)$ depends on $\{E_n\}$ though our notation does not indicate this).

The game $G_\rho(x)$ will be defined by induction on ρ ; it will be of length $\lambda + 1$ for some limit ordinal $\lambda < \omega_1$ with \forall making a move at position λ . (This move will be called the stopping move of the game $G_\rho(x)$).

Corresponding to each limit ordinal $\lambda < \omega_1$ fix, for the rest of this section, a sequence $\varphi_\lambda(n) \uparrow \lambda$.

For $\rho = 0$, the game $G_\rho(x)$ is played with \exists playing successively a_0, a_1, \dots while at position ω \forall plays a k (a stopping move) thereby producing, what we shall call a relevant sequence, $s = \langle a_0, a_1, \dots, a_{k-1} \rangle$. We say that \exists wins iff $x \in E_s$.

Now suppose $\rho = \mu + 1$. Then $R_\rho = \text{IR } R_\mu R_\mu^0$. Set $R_\mu = \emptyset$ and define

$$E = \Phi_n^x(\{E_n\})$$

$$E^u = \Phi_n^* (\{E_{u*n}\})$$

$$E^{<u_0, \dots, u_{k-1}, s>} = \Phi_m^o (\{ \Phi_n^* (\{ E_{<u_0, \dots, u_{k-1}>* \langle\langle s, m \rangle\rangle * n \}) \} \}). \quad (i)$$

The game $G_\rho(x)$ is played as follows. First \forall and \exists play $G_\mu(x)$ corresponding to the family $\{E^{<k>}\}$ with \forall making a (stopping) move at the last position thereby producing a relevant sequence s_0 . Then \forall and \exists play the dual of the game $G_\mu(x)$ (denoted by $G_\mu^o(x)$ which is played as in $G_\mu(x)$ by interchanging the moves of \forall and \exists) corresponding to the family $\{ \sim E_{\langle\langle s_0, m \rangle\rangle} \}_m$ with \exists making a (stopping) move at the last position to produce a relevant sequence t_0 . Let $u_0 = \langle s_0, t_0 \rangle$. Next \forall and \exists play $G_\mu(x)$ corresponding to the family $\{ E_{\langle u_0, n \rangle} \}_n$ with \forall making a (stopping) move at the last position to produce a relevant sequence s_1 . Then \forall and \exists play the dual game $G_\mu^o(x)$ associated with the family $\{ \sim E_{\langle u_0, \langle s_1, m \rangle \rangle} \}_m$ where \exists makes the last (stopping) move to produce a relevant sequence t_1 . Set $u_1 = \langle s_1, t_1 \rangle$; \exists and \forall continue to play alternately $G_\mu(x)$ and $G_\mu^o(x)$ as described above. At the end \forall plays some k so that $\langle u_0, \dots, u_{k-1} \rangle$ is the relevant sequence produced. Then \exists wins iff $x \in E_{\langle u_0, \dots, u_{k-1} \rangle}$.

Now suppose ρ is limit. Then $R_\rho = \mathbb{R}\Phi\Phi^0$, where Φ is defined by

$$\Phi(\{F_n\}) = \bigcap_i R_{\Phi_\rho(i)}(\{F_{\langle i, n \rangle}\}), \quad (ii)$$

for any family $\{F_n\}$. The game $G_\rho(x)$ is defined as follows. First \forall makes a move i_0 (called a choice move) which specifies the game to be played. Then \forall and \exists play the game $G_{\Phi_\rho(i_0)}(x)$ associated with the family $\{E_{\langle i_0, n \rangle}\}_n$ (see (i)) at the end of which \forall makes a (stopping) move producing a relevant sequence s_0 . Next \exists plays j_0 (a choice move) and then \forall, \exists start playing the dual game $G_{\Phi_\rho(j_0)}(x)$ associated with $\{\sim E_{\langle\langle i_0, s_0 \rangle, \langle j_0, m \rangle\rangle}\}_m$ with \exists moving at the last position producing a relevant sequence t_0 . Let $u_0 = \langle\langle i_0, s_0 \rangle, \langle j_0, t_0 \rangle\rangle$. Then \forall makes a choice move i_1 signaling the start of the game $G_{\Phi_\rho(i_1)}(x)$ associated with $\{E_{\langle u_0, \langle i_1, n \rangle \rangle}\}_n$ and the above procedure is repeated. At the end a stopping move k is made by \forall so that $\langle u_0, \dots, u_{k-1} \rangle$ is the relevant sequence produced. Then \exists wins iff $x \in E_{\langle u_0, \dots, u_{k-1} \rangle}$.

It is quite straightforward, though tedious, to show that for every $\rho < \omega_1$, if $x \in R_\rho(\{E_n\})$, then \exists wins $G_\rho(x)$.

For the converse, recall the definition of E_s^u from I:2.6. Define for every $s \in \text{Seq}$, the game $G_{\rho, s}(x)$ exactly as above with the following winning condition for \exists : if u_0, u_1, \dots

are the relevant sequences produced and if \forall plays k at the last position (stopping move), then \exists wins iff $x \in E_{s^* \langle u_0, u_1, \dots, u_{k-1} \rangle}$. Then one can show by induction on μ that if $x \notin E_s^\mu$ then \exists does not win $G_{\rho, s}(x)$. It then follows from Theorem 2.7 that if $x \notin E$ then \exists does not win $G_{\rho, e}(x) = G_\rho(x)$. Thus we have

$$x \in R_\rho(\{E_n\}) \leftrightarrow \exists \text{ wins } G_\rho(x).$$

Note that if $\rho < \omega$, then $G_\rho(x)$ is of length $\omega^{\rho+1} + 1$.

We now define by induction on ρ a game $G_\rho^1(x)$ of length $\omega + 1$ (and can thus be regarded as an ordinary Gale-Stewart game with the last move contributing to the pay-off set) associated with $\{E_n\}$, $x \in X$ such that

$$\exists \text{ wins } G_\rho^1(x) \leftrightarrow \exists \text{ wins } G_\rho(x).$$

In this game also the \forall player makes a choice move at the last position ω to produce a relevant sequence.

At the base step, $G_\rho^1(x)$ is the same as $G_\rho(x)$. We shall now describe $G_\rho^1(x)$ for successor ordinals ρ only, the limit case being much similar. So suppose $\rho = \mu + 1$. The game $G_\rho^1(x)$ will be such that if both players play according to some specified rules, then at the end \forall plays a k producing a relevant sequence $\langle u_0, \dots, u_{k-1} \rangle$ and \exists wins iff $x \in E_{\langle u_0, \dots, u_{k-1} \rangle}$. To see how $G_\rho^1(x)$ is played recall how

the game $G'_\rho(x)$ is described. We shall first define an intermediate game $G''_\rho(x)$ of length ω^2+1 and then obtain $G'_\rho(x)$. In the game $G''_\rho(x)$, first \forall and \exists play the game $G'_\mu(x)$ (associated with $\{E^{<n>}\}$) with \forall making a move at position ω (thereby producing some relevant sequence s_0). Next \forall and \exists play $(G'_\mu(x))^0$, the dual of the game $G'_\mu(x)$ (associated with $\{\sim E^{<<s_{0,m}>>}\}_m$), with \exists making a stopping move at the last position so as to produce some relevant sequence t_0 . Let $u_0 = \langle s_0, t_0 \rangle$. The game then proceeds with the two players alternately playing $G'_\mu(x)$ and $(G'_\mu(x))^0$. At position ω^2+1 , \forall plays k and produces the relevant sequence $\langle u_0, \dots, u_{k-1} \rangle$. If none of the players violate the prescribed rules in each of the above subgames, then \exists wins iff $x \in E_{\langle u_0, \dots, u_{k-1} \rangle}$.

The game $G'_\rho(x)$ is now easy to describe. Though its total length is $\omega+1$ we think of it as consisting of potentially infinite sequence of subgames each consisting of potentially infinite sequence of rounds. Each of these subgames "corresponds" to the corresponding subgame in $G''_\rho(x)$. If in any play of $G'_\rho(x)$ the j -th subgame actually goes through infinitely many rounds, then the $(j+1)$ th subgame never gets started. This keeps the total length within bounds.

In the j -th subgame of $G'_\rho(x)$, \forall and \exists make moves maintaining the rules of the j -th subgame of $G''_\rho(x)$. A move

of V followed by a move of \exists (or vice versa, depending on the order of the moves in the j -th subgame of $G_\rho''(x)$) will be called a round.

The i -th round of the $2j$ -th subgame opens with V signalling (by a choice of 0 or 1) either a challenge or a pass. If he challenges, the whole $2j$ -th subgame ends at once and the players proceed to the $(2j+1)$ th subgame; in this case we record $u_j \in \text{Seq}$ which codes the moves of the players in this subgame. If V passes, then the two players play the i -th round (according to the rules of $2j$ -th subgame in $G_\rho''(x)$); then the players proceed to the $(i+1)$ th round.

The $(2j+1)$ th subgame is played exactly like the $2j$ -th subgame with \exists (instead of V) signalling a challenge or a pass at the start of each round.

If some subgame goes on forever because \exists or V (as the case may be) fails to challenge on any round, then he forfeits the game. If this provision does not apply, a sequence $u_0, v_0; u_1, v_1; \dots$ (of relevant sequences) is generated. Then at position ω , V plays a j and \exists wins iff $x \in E_{\langle \langle u_0, v_0 \rangle, \dots, \langle u_{j-1}, v_{j-1} \rangle \rangle}$.

One can check that \exists wins $G_\rho'(x)$ iff \exists wins $G_\rho(x)$. In fact any winning strategy σ for \exists in $G_\rho'(x)$ gives rise in a canonical way to a winning strategy σ^*

for \exists in the game $G_\rho(x)$; and vice versa. Moreover, $G'_\rho(x)$ is a Σ_3^0 game. (For details at the finite levels, see [13, § 11]).

The game $J_\rho(x)$: We shall now define a game $J_\rho(x)$ associated with normal families $\{E_n\}$, analogous to $G_\rho(x)$, such that

\exists wins $J_\rho(x)$ iff E^x is comeager.

Now observe that in each ω -block in the game $G_\rho(x)$ (apart from the choice moves and the stopping moves) either \exists plays successively or \forall plays successively. Each move by \exists , which is not a choice move or a stopping move, in $G_\rho(x)$ will "correspond" (in the game $J_\rho(x)$) to a sequence move (i.e., a choice from $\omega^{<\omega}$) by \forall followed by a sequence move and an integer move by \exists . Similarly, each move by \forall in $G_\rho(x)$ which is not a choice move or a stopping move, will correspond to a sequence move and an integer move by \forall followed by a sequence move by \exists in $J_\rho(x)$. A choice move by \forall (\exists) in $G_\rho(x)$ corresponds to a choice move by \forall (\exists) in $J_\rho(x)$ and a stopping move by \forall (\exists) in $G_\rho(x)$ corresponds to a stopping move by \forall (\exists) in $J_\rho(x)$. Thus each λ -block, λ limit, in $J_\rho(x)$ will give rise to two relevant sequences, one from the integer moves and the other from the sequence moves. At the end of the game $J_\rho(x)$ \forall plays some k thus producing two relevant sequences of the form $u = \langle u_0, \dots, u_{k-1} \rangle$ and

$v = v_0 * v_1 * \dots * v_{k-1} \dots$. We say that \exists wins iff E_u^x is comeager in $\Sigma(v)$.

More formally we define $J_\rho(x)$ as follows. First recall the definition of the Vaught operation \mathcal{V} and the operations S_ρ from §6.5. For convenience we replace Seq by $\omega^{<\omega}$. Let $\{E_n\}$ be a normal family of subsets of $\omega^\omega \times \omega^\omega$, s a finite sequence and $x \in \omega^\omega$. When $\rho = 0$, $J_{\rho,s}(x)$ is the Vaught game and is played as follows: First \forall chooses a finite sequence s_0 , then \exists replies with a finite sequence t_0 and a natural number k_0 ; then \forall chooses a finite sequence s_1 while \exists replies with t_1, k_1 and so on. In the end \forall plays an integer i (a stopping move) to produce two relevant sequences $u = \langle k_0, \dots, k_{i-1} \rangle$ and $v = s_0 * t_0 * \dots * s_{i-1} * t_{i-1}$. \exists wins iff E_u^x is comeager in $\Sigma(s * v)$.

Now suppose $\rho = \mu + 1$. Then $S_\rho = \text{IRS}_{\mu} S_\mu^0$. Now \forall and \exists first play the game $J_{\mu,s}(x)$ associated with the normal family $\{E^{<n>}\}$ (defined as in (i) above) with \forall making a (stopping) move at the end so as to produce two relevant sequences $s_0 \in \text{Seq}$ and $t_0 \in \omega^{<\omega}$. Then \forall plays a finite sequence t and the two players play the dual game $J_{\mu, s * t_0 * t}(x)$ (associated with $\{\sim E^{<<s_0, m>>}\}_m$) with \exists playing at the end to produce two relevant sequences $s'_0 \in \text{Seq}$, $t'_0 \in \omega^{<\omega}$ with $t'_0 = t * t''_0$ for some t''_0 . Let

$u_0 = \langle s_0, s'_0 \rangle$, $v_0 = t_0 * t'_0$. Next \exists plays a $t' \in \omega^{<\omega}$ and two players play $J_{\mu, s * v_0 * t'}(x)$ (associated with $\{E_{\langle u_0, n \rangle}^x\}_n$) with \forall making the last move to produce two relevant sequences s_1, t_1 and then \forall plays $t'' \in \omega^{<\omega}$ and they play the dual game $J_{\mu, s * v_0 * t_1 * t''}^0(x)$. Thus the two players alternately play $J_{\mu}(x)$ and $J_{\mu}^0(x)$ as described above producing two sequences $\{u_i\}$ and $\{v_i\}$ of relevant sequences. At the end of the play \forall plays a k to produce the relevant sequences $u = \langle u_0, \dots, u_{k-1} \rangle \in \text{Seq}$ and $v = v_0 * v_1 * \dots * v_{k-1}$. Then \exists wins iff E_u^x is comeager in $\Sigma(s * v * v')$, where v' is the next sequence move (played by \forall).

As in the case of $G_{\rho}(x)$ we can now define exactly as above, a game $J'_{\rho, s}(x)$ of length $\omega+1$ with \forall playing in the end to produce two relevant sequences such that

$$\exists \text{ wins } J'_{\rho, s}(x) \leftrightarrow \exists \text{ wins } J_{\rho, s}(x).$$

We now show that E^x is comeager in $\Sigma(s)$ iff \exists wins $J_{\rho, s}(x)$. This is really Theorem I:6.6 in disguise.

Let $\tilde{J}_{\rho, s}(x)$ be a variant of the game $J_{\rho, s}(x)$ which is played exactly as in $J_{\rho, s}(x)$ such that if $u_0, u_1, \dots, v_0, v_1, \dots$ are the two sequences of relevant sequences produced before \forall plays an integer k in the last position, then \exists wins $\tilde{J}_{\rho, s}(x)$ iff $s * v_0 * v_1 * \dots \in E_{\langle u_0, \dots, u_{k-1} \rangle}^x$.

Lemma 1: With notation as above, we have E^X is comeager in $\Sigma(s)$ iff \exists wins $J_{\rho,s}(x)$ iff \exists wins $\tilde{J}_{\rho,s}(x)$.

Proof. We shall prove this by induction on ρ . We prove it only for the successor case, the limit case being similar.

Without loss of generality assume that $\{E_n\}$ is regular.

Suppose $\rho = \mu + 1$. Then $R_\rho = \mathbb{R}R_\mu R_\mu^0$. Put $\phi = R_\mu$.

1°. Since $E = \phi(\{E_n^{<n>}\})$, E^X is comeager in $\Sigma(s)$ implies

\exists wins $J_{\mu,s}(x)$ associated with $\{E_n^{<n>}\}$, by induction

hypothesis. So to win the game $J_{\rho,s}(x)$ \exists first plays with

a winning strategy σ_0 , say, in $J_{\mu,s}(x)$. If in a complete

run of this subgame s_0, t_0 are the relevant sequences produced,

then $(E^{<s_0>})^X$ is comeager in $\Sigma(s*t_0)$. Now

$E^{<s_0>} = \phi_m^0(\{E_m^{<<s_0,m>>}\})$. Thus $(E^{<s_0>})^X$ is comeager in $\Sigma(s*t_0)$

+ $(\phi_m(\{\sim E_m^{<<s_0,m>>}\}))^X$ is meager in $\Sigma(s*t_0)$

+ $(\forall u) \sim [(\phi_m(\{\sim E_m^{<<s_0,m>>}\}))^X \text{ is comeager in } \Sigma(s*t_0*u)]$

+ $(\forall u) [\exists \text{ does not win the game } J_{\mu,s*t_0*u}(x)$

associated with $\{\sim E_m^{<<s_0,m>>}\}]$

+ $(\forall u) [\exists \text{ wins the dual game } J_{\mu,s*t_0*u}^0(x)$

associated with $\{\sim E_m^{<<s_0,m>>}\}]$.

Thus \exists wins the second subgame of $J_{\rho, s}(x)$ with strategy τ_0 , say, when the first subgame has been played with s_0, t_0 as the generated relevant sequences. Thus \exists plays with τ_0 and if s'_0, t'_0 are the relevant sequences produced at the end of a run, then we have (since \exists wins the dual game)

$$(E_{\langle\langle s_0, s'_0 \rangle\rangle}^x)^x \text{ is nonmeager in } \Sigma(s*t_0*t'_0)$$

and so for some $v \in \omega^{<\omega}$,

$$(E_{\langle\langle s_0, s'_0 \rangle\rangle}^x)^x \text{ is comeager in } \Sigma(s*t_0*t'_0*v).$$

By regularity $E_{\langle\langle s_0, s'_0 \rangle\rangle}^x$ is comeager in $\Sigma(s*t_0*t'_0*v)$. In the next subgame \exists plays this v and then continues as described above. This clearly defines a winning strategy for \exists in the game $J_{\rho, s}(x)$.

2°. It is clear that if \exists wins $J_{\rho, s}(x)$ then \exists wins $J_{\rho, s}(x)$. For, without loss of generality, $\{E_n^x\}$ may be taken to be a family of clopen sets and so E_u^x is comeager in $\Sigma(v)$ iff $\Sigma(u) \subseteq E_u^x$.

3°. Now suppose \exists wins $J_{\rho, s}(x)$ with strategy σ . Now suppose the two players play the first two subgames with \exists playing according to σ . Suppose u_0, v_0 are the relevant sequences produced. In the next subgame \exists plays some v as dictated by σ . If for some $t \in \omega^{<\omega}$,

$$E_{u_0}^x \cap \Sigma(s*v_0*v*t) = \emptyset, \text{ then } \forall \text{ can beat } \sigma \text{ by playing}$$

this t as his next move. (Apart from the choice move, if there is one, it is \forall 's turn to play a finite sequence). Thus for all $t \in \omega^{<\omega}$, $E_{u_0}^x \cap \Sigma(s * v_0 * v * t) \neq \emptyset$; and this is true at every stage. This shows that if $\{u_i\}$ and $\{v_i\}$ are the two sequences produced before \forall plays at the last position, then $s * v_0 * v_1 * \dots \in E_{\langle u_0, \dots, u_{k-1} \rangle}^x$ for every k and so \exists wins $\overline{J}_{\rho, s}(x)$ with the same strategy σ .

4^o. It now remains to show that if \exists wins $J_{\rho, s}(x)$ then E^x is comeager in $\Sigma(s)$. To this end we show that if E^x is not comeager in $\Sigma(s)$, then \exists wins the dual game $J_{\rho, s}^o(x)$. This is shown as in 1^o using the decomposition in Lemma I:6.2. We omit the proof. ||

With notation as above we now have the following

Lemma 2: E^x is comeager iff \exists wins $J_{\rho, e}(x)$. Moreover, with any winning strategy σ for \exists in $J_{\rho, e}(x)$ one can associate a winning strategy σ^* such that if $\{u_i\}$ and $\{v_i\}$ are two relevant sequences produced (before \forall plays a k in the last position) in a complete run of $J_{\rho, e}(x)$ when \exists plays with σ^* , the sequence $\delta = v_0 * v_1 * v_2 * \dots$ is in E^x .

Proof. The first part of the lemma is simply Lemma 1. Now E^x is comeager iff player II wins the Banach-Mazur game $\sim E^x$ i.e., iff

$$\begin{aligned}
& (\forall s_0) (\exists t_0) (\forall s_1) (\exists t_1) \dots \{ \delta \in E^X \} \\
\leftrightarrow & (\forall s_0) (\exists t_0) (\forall s_1) (\exists t_1) \dots \{ \exists \text{ wins } G_\rho(x, \delta) \} \\
\leftrightarrow & (\forall s_0) (\exists t_0) (\forall s_1) (\exists t_1) \dots \{ \exists \text{ wins } G'_\rho(x, \delta) \}
\end{aligned}$$

where $\delta = s_0 * t_0 * s_1 * t_1 * \dots$

Now, \exists wins $J_{\rho, e}(x) = J_\rho(x)$ iff \exists wins $\tilde{J}_\rho(x)$ (with same strategy). Further, the game $\tilde{J}_\rho(x)$ is obtained from $G_\rho(x, \delta)$ by interlacing the integer moves of $G_\rho(x, \delta)$ with sequence moves such that the sequence moves of a complete run in $\tilde{J}_\rho(x)$ generate δ . Thus if $\tilde{J}'_\rho(x)$ is obtained from $\tilde{J}_\rho(x)$ in the same way as $J'_\rho(x)$ is obtained from $J_\rho(x)$ or $G'_\rho(x)$ is obtained from $G_\rho(x)$; then observe that $\tilde{J}'_\rho(x)$ is simply $G'_\rho(x, \delta)$ interlaced with a Banach-Mazur game. Thus we have

$$(\forall s_0) (\exists t_0) (\forall s_1) (\exists t_1) \dots \{ \exists \text{ wins } G'_\rho(x, \delta) \} \leftrightarrow \exists \text{ wins } \tilde{J}'_\rho(x),$$

which is of the form of the equivalence in Lemma I:7.2. Now suppose \exists wins $J_\rho(x)$. Then he wins $\tilde{J}'_\rho(x)$ with some strategy σ . Now modify σ to σ^* so that any complete play in $\tilde{J}'_\rho(x)$ corresponds to a complete play in the (left-side) Banach-Mazur game as in Lemma I:7.2. The strategy σ^* gives rise to a strategy σ^{**} for \exists in $J_\rho(x)$ such that if $\{u_i\}$ and $\{v_i\}$ are two sequences of relevant sequences formed when \exists plays σ^{**} , then the same two relevant sequences can be obtained, by

simulation, in a play of $\mathcal{J}'_{\rho}(x)$ when \exists plays σ^* . This can be easily seen from the way $\mathcal{J}'_{\rho}(x)$ is obtained from $\mathcal{J}_{\rho}(x)$. Consequently, $\delta = v_0 * v_1 * v_2 * v_3 * \dots$ is generated by a complete play of the (left-side) Banach-Mazur game when player II plays with a winning strategy. Thus $\delta \in E^X$. ||

As we have mentioned earlier the above lemma is really Theorem I:6.6, stated in a different form. Thus it is not surprising that for every ρ and every normal family $\{E_n\}$, one can show that

$$\exists \text{ wins } J_{\rho, S}(x) \leftrightarrow (\exists \eta \in M_{\rho}) (\forall n \in \eta) [E_{f_{\rho}}^X(x) \text{ is comeager in } \Sigma(s * g_{\rho}(n))],$$

where $\Phi_{M_{\rho}} = S_{\rho}$ and f_{ρ} and g_{ρ} are suitable functions independent of $\{E_n\}$. We shall now show that if \exists wins $J_{\rho, S}(x)$ ($J_{\rho, S}^{\circ}(x)$) with some strategy σ_x , then there is some $\eta \in M_{\rho}$ (M_{ρ}°) such that for each $n \in \eta$, there is a complete play in $J_{\rho, S}(x)$ ($J_{\rho, S}^{\circ}(x)$) with \exists playing with σ_x such that $f_{\rho}(n)$ and $g_{\rho}(n)$ are relevant sequences produced. We will show this by induction on ρ . Suppose $\rho = 0$ and \exists wins $J_{\rho, S}(x)$ with strategy σ_x . Then \exists wins the following game with strategy σ_x :

$$(\forall s_0) (\exists t_0) (\exists k_0) (\forall s_1) (\exists t_1) (\exists k_1) \dots$$

$$\dots (\forall i) [E_{k_0, \dots, k_{i-1}}^X \text{ is comeager in } \Sigma(s * s_0 * t_0 * \dots * s_{i-1} * t_{i-1})]$$

Let $\eta = \{\langle\langle s_0, \langle k_0, t_0 \rangle \rangle, \dots, \langle s_{i-1}, \langle k_{i-1}, t_{i-1} \rangle \rangle \rangle\}$:

$s_0, t_0, k_0; s_1, t_1, k_1; \dots; s_{i-1}, t_{i-1}, k_{i-1}$ is an initial segment of a play in the above game when \exists follows σ_x .

It is easy to see that η is in the canonical base for S_0 .

Define

$$f_0(\langle\langle s_0, \langle k_0, t_0 \rangle \rangle, \dots, \langle s_{i-1}, \langle k_{i-1}, t_{i-1} \rangle \rangle \rangle) = \langle k_0, \dots, k_{i-1} \rangle$$

$$f_0(n) = 1 \text{ otherwise;}$$

$$g_0(\langle\langle s_0, \langle k_0, t_0 \rangle \rangle, \dots, \langle s_{i-1}, \langle k_{i-1}, t_{i-1} \rangle \rangle \rangle) = s_0 * t_0 * \dots * s_{i-1} * t_{i-1}.$$

$$g_0(n) = 1 \text{ otherwise.}$$

The dual case is shown similarly.

We now prove our assertion for successor ordinals ρ , the limit case being similar. So let $\rho = \mu + 1$ and suppose \exists wins $J_{\rho, s}(x)$ with strategy σ_x . We show that \exists wins the game (ii) below with a strategy σ'_x with the following property: If $\eta_0, n_0, u_0, \xi_0, m_0, v_0, \dots$ is a complete play consistent with σ'_x then for each k there is a complete play in $J_{\rho, s}(x)$ consistent with σ_x such that

$$\langle\langle f_\mu(n_0), f_\mu(m_0) \rangle, \dots, \langle f_\mu(n_{k-1}), f_\mu(m_{k-1}) \rangle \rangle \text{ and}$$

$$g_\mu(n_0) * u_0 * g_\mu(m_0) * v_0 * \dots * g_\mu(m_{k-1}) * v_{k-1}$$

are the corresponding relevant sequences produced.

$(\exists \eta_0 \in M_\rho) (\forall n_0 \in \eta_0) (\forall u_0) (\forall \xi_0 \in M_\mu) (\exists m_0 \in \xi_0) (\exists v_0) \dots$

$\dots (\forall k) [E^x \langle \langle f_\mu(n_0), f_\mu(m_0) \rangle, \dots, \langle f_\mu(n_{k-1}), f_\mu(m_{k-1}) \rangle \rangle$

is comeager in $\Sigma(s * g_\mu(n_0) * u_0 * g_\mu(m_0) * v_0 * \dots * g_\mu(m_{k-1}) * v_{k-1})]$

... (ii)

Since \exists wins $J_{\rho, s}(x)$, he wins $J_{\mu, s}(x)$ (associated with the normal family $\{E^{<n>}\}$). So for some $\eta_0 \in M_\mu$ and for every $n \in \eta_0$ there is a complete play in the subgame $J_{\mu, s}(x)$ (consistent with σ_x) such that $f_\mu(n)$ and $g_\mu(n)$ are the relevant sequences produced. So in the game (ii) \exists plays such an η_0 . Suppose \forall plays $n_0 \in \eta_0$, a finite sequence u_0 and $\xi_0 \in M_\mu$. Now \exists wins the dual game $J_{\mu, s * g_\mu(n_0) * u_0}^0(x)$ associated with $\{\sim E^{\langle \langle f_\mu(n_0), m \rangle \rangle}_m\}$. Hence for every $\xi_0 \in M_\mu$ there is a $m_0 \in \xi_0$ such that $\langle f_\mu(n_0), f_\mu(m_0) \rangle$ and $g_\mu(n_0) * u_0 * g_\mu(m_0)$ are the relevant sequences (consistent with σ_x) produced in the dual game and consequently

$E^{\langle \langle f_\mu(n_0), f_\mu(m_0) \rangle \rangle}$ is comeager in $\Sigma(s * g_\mu(n_0) * u_0 * g_\mu(m_0) * v_0)$

for some $v_0 \in \omega^{<\omega}$. Thus \exists plays such m_0, v_0 . By regularity,

$E^{\langle \langle f_\mu(n_0), f_\mu(m_0) \rangle \rangle}$ is comeager in $\Sigma(s * g_\mu(n_0) * u_0 * g_\mu(m_0) * v_0)$.

Proceeding thus it is clear that \exists can win the game (i). Finally observe that the function $x \mapsto \sigma_x$ can be chosen to be \mathcal{BR}_0^p -measurable (see the proof of Theorem II:1.1 and Theorem I:3.8). Consequently, $x \mapsto \sigma'_x$ can also be chosen to be \mathcal{BR}_0^p -measurable, since to describe σ'_x one needs to describe relevant sequences with respect to σ_x . Finally, for each $x \in E^*$ we choose σ_x (in a \mathcal{BR}_0^p -measurable way) to satisfy the conclusion of Lemma 2 above. Thus for each $x \in E^*$ we can choose (in a \mathcal{BR}_0^p -measurable way) σ'_x such that σ'_x is a winning strategy in the game (ii) above and satisfies the conclusion of Lemma I:7.5. \parallel

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