

NOTE

Rigidity Theorems for Partial Linear Spaces

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Let $(\mathcal{X}, \mathcal{L})$ be a partial linear space such that each of its points is contained in at least $n+1$ lines. Fix a prime p . Let $\mathcal{C}_{\mathcal{L}} \subseteq \mathbb{F}_p^{\mathcal{X}}$ denote the p -ary code generated by lines in \mathcal{L} and $\mathcal{C}_{\mathcal{L}}^{\perp}$ denote its dual. In this article, we prove that the minimum weight of $\mathcal{C}_{\mathcal{L}}^{\perp}$ is at least $2n+2-\frac{2n}{p}$. (As a special case of our result we see that the minimum weight of the p -ary code dual to the code generated by lines in a projective plane of order n is at least $2n+2-\frac{2n}{p}$. This bound sharply improves upon the earlier known bounds as given in [1, Corollary 6.3.1 and Theorem 6.4.2]). We also prove that the induced structure on the support of a word of weight $2n+2-\frac{2n}{p}$ in $\mathcal{C}_{\mathcal{L}}^{\perp}$, if it exists, is isomorphic to the join of two Steiner 2-designs with $\binom{n}{p}$ lines through every point and p points on every line (Theorem 2).

For a projective plane Π of order n , let \mathcal{C} denote the p -ary code generated by lines in Π and \mathcal{C}^{\perp} its orthogonal. We denote by $Hull_p(\Pi)$ the intersection $\mathcal{C} \cap \mathcal{C}^{\perp}$. The plane Π is said to be *tame at p* if the minimum weight of $Hull_p(\Pi)$ is $2n$ and the words of minimum weight are scalar multiples of vectors of the form $\psi_L - \psi_M$, where ψ_L and ψ_M denote characteristic functions of two lines L and M in Π . It is called *tame* if it is tame at all the primes that divide n . Thus, our result also implies that a projective plane of order p is tame. As only desarguesian planes are believed to be tame (see [1, Theorem 6.9.1] and the discussion thereon), the results of this article provide strong evidence towards the validity of the prime order case of the Hamada-Sachar conjecture.

2. RIGIDITY THEOREMS

In this section, in Theorem 1, we prove a generalisation of Bagchi's conjecture for linear spaces of prime order. Let us fix a prime p .

Let \mathcal{X} be a non-empty set and $\mathcal{L} \subseteq P(\mathcal{X})$ be a collection of non-empty subsets of \mathcal{X} . The pair $(\mathcal{X}, \mathcal{L})$ is called a *partial linear space* if for any $A \neq B$ in \mathcal{X} , $\{A, B\}$ is a subset of at most one $L \in \mathcal{L}$. Elements of \mathcal{X} are called points and elements of \mathcal{L} are called lines. By a line through A , we mean a line containing the point A . The p -ary code generated by the lines in \mathcal{L} is denoted by $\mathcal{C}_{\mathcal{L}}$ and its dual by $\mathcal{C}_{\mathcal{L}}^{\perp}$. We call a partial linear space $(\mathcal{X}, \mathcal{L})$ *non-trivial at p* if the p -ary code $\mathcal{C}_{\mathcal{L}}$ is not equal to $\mathbb{F}_p^{\mathcal{X}}$.

Let $(\mathcal{X}, \mathcal{L})$ be a partial linear space such that at least $n+1$ lines pass through every point in \mathcal{X} , and let \mathcal{Y} be a subset of \mathcal{X} such that no line in \mathcal{L} intersects \mathcal{Y} in one point. Let \mathcal{L}' denote the set of lines in \mathcal{L} obtained by intersecting lines in \mathcal{L} with \mathcal{Y} . Then $(\mathcal{Y}, \mathcal{L}')$ is also a partial linear space with $\geq n+1$ lines through every point. This is called *the induced structure on \mathcal{Y}* . Let $(\mathcal{X}, \mathcal{L}_1)$ and $(\mathcal{Y}, \mathcal{L}_2)$ be two partial linear spaces with disjoint point sets. Their *join* is the partial linear space $(\mathcal{X} \cup \mathcal{Y}, \mathcal{L}_1 \cup \mathcal{L}_2 \cup \mathcal{L}_3)$ where \mathcal{L}_3 is the collection of sets $\{x, y\}$ for all $x \in \mathcal{X}$ and $y \in \mathcal{Y}$. If $(\mathcal{X}_1, \mathcal{L}_1)$ and $(\mathcal{X}_2, \mathcal{L}_2)$ are two partial linear spaces such that p divides the cardinality of every line in $\mathcal{L}_1 \cup \mathcal{L}_2$, then their join is non-trivial at p since the non-zero word w defined by $w(x) = (-1)^{|x|}$ for all $x \in \mathcal{X}_i$ is in the dual of the code generated by its lines.

A partial linear space in which any pair of distinct points lie on a (unique) line is called a *linear space*. A *Steiner 2-design* is a linear space whose lines have same cardinality and the same number of lines pass through any point. The cardinality of a set S is denoted by $|S|$.

THEOREM 1. *If a partial linear space $(\mathcal{X}, \mathcal{L})$ is non-trivial at p and has at least $n+1$ lines through every point, then $|\mathcal{X}| \geq 2n+2 - \frac{2^n}{p}$. Moreover, the equality holds if and only if $(\mathcal{X}, \mathcal{L})$ is isomorphic to the join of two Steiner 2-designs with $\frac{n}{p}$ lines through every point and p points on every line.*

Proof. A Steiner 2-design with $\frac{n}{p}$ lines through every point and p points on each line has $n+1 - \frac{n}{p}$ points. Thus the join of two such designs is a linear space with $2n+2 - \frac{2^n}{p}$ points and exactly $n+1$ lines through every point. Moreover, by the above remark, it is non-trivial at p .

To prove the converse, let $(\mathcal{Y}, \mathcal{L}')$ be a partial linear space with $\geq n+1$ lines through every point, with $|\mathcal{Y}| \leq 2n+2 - \frac{2^n}{p}$ and which is non-trivial at p . Let \mathcal{X} denote the support of a word w of least weight in $\mathcal{C}_{\mathcal{L}'}^{\perp}$. The induced structure $(\mathcal{X}, \mathcal{L}_0)$ is a partial linear space with $\geq n+1$ lines through every point, with $|\mathcal{X}| \leq 2n+2 - \frac{2^n}{p}$ and such that $\mathcal{C}_{\mathcal{L}_0}^{\perp}$ is generated by the restriction of w to \mathcal{X} . To simplify the notation, we also denote the

restriction of w to its support by w . Thus $(\mathcal{X}, \mathcal{L}_0)$ is non-trivial at p and the characteristic function of a subset L' of \mathcal{X} is in $\mathcal{C}_{\mathcal{L}_0}$ if it is in the dual of $\langle w \rangle$. We now replace a line $L \in \mathcal{L}_0$ by two non-empty subsets L' and L'' whenever $L = L' \cup L''$ is a partition of L with the property that the characteristic function of L' (and hence of L'' also) is in $\mathcal{C}_{\mathcal{L}_0}$. When we reach a stage where no such replacements can be made, we get a set of lines on \mathcal{X} such that the characteristic function of no proper subset of a line is in $\mathcal{C}_{\mathcal{L}_0}$ (we thank the referee for pointing out this construction). We denote the new set of lines by \mathcal{L} . Note that the construction of \mathcal{L} from \mathcal{L}_0 is not canonical. However, for any such \mathcal{L} , the partial linear space $(\mathcal{X}, \mathcal{L})$ will have at least $n+1$ lines passing through every point and will have the following two properties:

1. $\mathcal{C}_{\mathcal{L}} = \mathcal{C}_{\mathcal{L}_0}$ and $\mathcal{C}_{\mathcal{L}}$ does not contain the characteristic function of a proper non-empty subset of a line in \mathcal{L} .
2. $\mathcal{C}_{\mathcal{L}}^\perp$ is one dimensional and \mathcal{X} is the support of its generator w .

We deal with such a partial linear space in the remainder of the proof.

For any $A \in \mathcal{X}$, let x_A, y_A, z_A denote the number of lines through A of cardinality 2, 3, 4 respectively. Fix a point Q of \mathcal{X} such that $x_Q \leq x_A$ for all $A \in \mathcal{X}$. We normalise w by assuming $w(Q) = -1$. We now colour \mathcal{X} by elements of \mathbb{F}_p using this w , wherein a point P gets the colour $w(P)$. As the characteristic function of a line is in the dual of $\langle w \rangle$, the sum of the colours occurring on any line is $0 \pmod p$. Also, by the construction of $(\mathcal{X}, \mathcal{L})$, the colours occurring on any proper subset of a line do not add up to $0 \pmod p$. Hence the lines containing a point each of colour α and $-\alpha$ have length 2 and the lines whose every point has same colour have length p . Let $\mathcal{S} = \{\alpha \in \mathbb{F}_p \mid w(P) = \alpha \text{ for some point } P\}$ denote the set of colours and let $X_\alpha \subseteq \mathcal{X}$ denote the set of points with colour α . Note that $0 \notin \mathcal{S}$ since \mathcal{X} is the support of w .

Let S_A denote the union of all the lines through a point A . Then $1 + x_A + 2y_A + 3z_A + 4(n + 1 - x_A - y_A - z_A) \leq |S_A| \leq 2n + 2 - \frac{2n}{p}$ so that

$$3x_A + 2y_A + z_A \geq 2n + 3 + \frac{2n}{p}. \tag{1}$$

Similarly, $1 + x_A + 2y_A + 3(n + 1 - x_A - y_A) \leq 2n + 2 - \frac{2n}{p}$. Hence

$$2x_A + y_A \geq n + 2 + \frac{2n}{p}. \tag{2}$$

A line of cardinality s is called an s -line. If L_1, \dots, L_m denote all the s -lines through A with $s > 3$, then $|S_A| \geq 1 + x_A + 2(n + 1 - x_A - m) + \sum_{i=1}^m (|L_i| - 1)$. The bound on $|\mathcal{X}|$ now implies that

$$x_A \geq 1 + \frac{2n}{p} + \sum_{i=1}^m (|L_i| - 3). \tag{3}$$

The theorem is trivial for $p = 2$. Thus we assume that $p \geq 3$. The number of 2-lines through any point is at least x_Q . Hence for every $\alpha \in \mathcal{S}$, at least x_Q points of \mathcal{X} are coloured $-\alpha$. As $x_Q > 0$, we have $\alpha \in \mathcal{S} \Rightarrow -\alpha \in \mathcal{S}$ and $|X_\alpha| \geq x_Q$ for all $\alpha \in \mathcal{S}$. Also, $|\mathcal{S}|$ is even as $0 \notin \mathcal{S}$. Let $|\mathcal{S}| = 2r$ for some r with $1 \leq r \leq \frac{p-1}{2}$. We then have

$$rx_Q \leq n + 1 - \frac{n}{p}. \tag{4}$$

First consider the case when $r = 1$. In this case $\mathcal{S} = \{1, -1\}$ and \mathcal{L} contains only 2-lines and p -lines. Also, since $x_Q \leq n + 1 - \frac{n}{p}$, the number of p -lines through Q is at least $\frac{n}{p}$. Therefore, we must have

$$2n + 2 - \frac{2n}{p} = 1 + \frac{n}{p}(p - 1) + \left(n + 1 - \frac{n}{p} \right) \leq |S_Q| \leq 2n + 2 - \frac{2n}{p}.$$

Thus $x_Q = n + 1 - \frac{n}{p}$ and the number of p -lines through Q must be $\frac{n}{p}$. Therefore, exactly $n + 1$ lines pass through Q and $|\mathcal{X}| = |S_Q| = 2n + 2 - \frac{2n}{p}$. Moreover, $|X_1| = |X_{-1}| = n + 1 - \frac{n}{p}$. As the number of 2-lines through a point $A \in X_i$ is at most $|X_{-i}|$, we must have $x_A = x_Q$ for all $A \in \mathcal{X}$. Thus the above argument for Q also holds for any $A \in \mathcal{X}$ so that the number of p -lines through any point of \mathcal{X} is $\frac{n}{p}$. Thus X_i with all the p -lines contained in it is isomorphic to a Steiner 2-design whose lines have length p and having $\frac{n}{p}$ lines through every point. As $\{x, y\} \in \mathcal{L}$ for all $x \in X_1$ and $y \in X_{-1}$, the partial linear space $(\mathcal{X}, \mathcal{L})$ is their join.

Thus $(\mathcal{X}, \mathcal{L})$ is a linear space. Since any substitution of a line in \mathcal{L}_0 by its partition would prevent $(\mathcal{X}, \mathcal{L})$ from being a linear space, we get $\mathcal{L} = \mathcal{L}_0$. Also, $2n + 2 - \frac{2n}{p} = |\mathcal{X}| \leq |\mathcal{Y}| \leq 2n + 2 - \frac{2n}{p}$. Hence $\mathcal{X} = \mathcal{Y}$ so that $\mathcal{L}_0 = \mathcal{L}'$. Thus $(\mathcal{X}, \mathcal{L})$ is the partial linear space $(\mathcal{Y}, \mathcal{L}')$ that we started off with. This proves the theorem for $r = 1$.

Thus we may assume that $r > 1$. To complete the proof of the theorem, we wish to prove that no such $(\mathcal{X}, \mathcal{L})$ exists.

As $1 < r \leq \frac{p-1}{2}$, we have $p > 3$ so that $(p - 1)(1 + \frac{2n}{p}) > 2n + 2 - \frac{2n}{p}$. Hence, by (3) and (4), $r \neq \frac{p-1}{2}$. We thus have $1 < r < \frac{p-1}{2}$.

Let G_Q denote the graph whose vertex set is \mathcal{S} and whose edges are given by the rule

$$\alpha \text{ and } \beta \text{ are adjacent if and only if } \alpha + \beta = 0 \text{ or } 1 \text{ in } \mathbb{F}_p.$$

As the sum of the colours occurring on any line is zero, a 3-line having non-empty intersections with X_{-1} and X_α also has non-empty intersection with $X_{1-\alpha}$. If α has degree one in G_Q , $X_{1-\alpha}$ is empty. Thus if α has degree one in G_Q , then $L \cap X_\alpha$ is either empty or is $\{Q\}$ for a 3-line L through Q . Since $0 \notin \mathcal{S}$, the degree of 1 is one in G_Q . The only possible loop of G_Q is at the vertex $\frac{p+1}{2}$ and this loop occurs if and only if $\frac{p+1}{2}$ belongs to \mathcal{S} .

If $\alpha_1 \cdots \alpha_m \alpha_1$ is a non-trivial cycle in G_Q , m must be even, as the edges of types $\{\alpha, -\alpha\}$ and $\{\alpha, 1-\alpha\}$ must alternate in the cycle. Also,

$$(\alpha_1 + \alpha_2) + \cdots + (\alpha_{m-1} + \alpha_m) = (\alpha_2 + \alpha_3) + \cdots + (\alpha_m + \alpha_1).$$

Hence m must be a multiple of $2p$, as one of these sums is zero and the other is $\frac{m}{2}$. However, $m \leq |\mathcal{S}| < p-1$, and hence G_Q does not contain any cycles. Therefore, each connected component of G_Q is a path. In case $\frac{p+1}{2} \in \mathcal{S}$, one of these paths has a loop at one end.

Case 1. The graph G_Q is connected.

In this case, G_Q is a path. If 1 is the only vertex of degree one, then the other end of this path must be a loop at $\frac{p+1}{2}$. Therefore, in this case G_Q is the path $1(-1)2(-2)\cdots(\frac{p-1}{2})(\frac{p+1}{2})(\frac{p+1}{2})$. However, this forces $|\mathcal{S}| = p-1$ which cannot happen. Therefore, G_Q must have two vertices of degree one and it is the path $1(-1)2(-2)\cdots r(-r)$ with $1 < r < \frac{p-1}{2}$.

Since $r > 1$, $Q \notin X_{-r}$, so that $|\{Q\} \cup X_{-r}| \geq 1 + x_Q$. Let $T = \mathcal{X} \setminus (\{Q\} \cup X_{-r})$ and let l denote the number of s -lines through Q with $s > 2$ and which contain at most one point from T . Note that every 2-line through Q contains exactly one point from T . Counting points of T which lie on lines through Q , we get

$$2(n+1-x_Q-l)+x_Q \leq |T| \leq 2n+2-\frac{2n}{p}-(1+x_Q).$$

This means $l \geq \frac{n}{p} + \frac{1}{2} > \frac{n}{p}$.

Let L be one of these l lines. If $L \subset X_{-r} \cup \{Q\}$, then L must contain at least $\frac{p-1}{r}$ points from X_{-r} , as the colours on L add up to $0 \pmod{p}$. As $Q \notin X_{-r}$, in this case $|L| \geq \frac{p+r-1}{r}$. If L contains one point from T , the colour of that point in L must be $(1+(|L|-2)r) \pmod{p}$. As an integer, $1+(|L|-2)r > r$ as $|L| > 2$. If this number is greater than p , then we must have $|L|-2 > \frac{p-1}{r}$ so that $|L| > \frac{p+2r-1}{r}$. If it is less than p , then this number itself represents the colour of the remaining point (which is not $p-r$ by our assumption). Thus, as $\mathcal{S} = \{1, \dots, r, p-r, \dots, p-1\}$, we must have $p-r+1 \leq 1+(|L|-2)r$. This forces $|L|-2 \geq \frac{p-r}{r}$, so that $|L| \geq \frac{p+r}{r}$. Considering all the cases, we see that if a line L through Q contains at most one

point from T , then $|L| \geq \frac{p+r-1}{r}$. As $r < \frac{p-1}{2}$, we have $|L| > 3$. We may then apply (3) and the fact that $l > \frac{n}{p}$ to get

$$x_Q > 1 + \frac{2n}{p} + \frac{n}{p} \left(\frac{p+r-1}{r} - 3 \right) = 1 + \frac{n}{r} - \frac{n}{pr}.$$

However, this bound contradicts (4).

Case 2. G_Q is disconnected.

In this case, let $\mathcal{S}' \subseteq \mathcal{S}$ denote the set of colours which have degree one in G_Q . As G_Q is disconnected, we must have $|\mathcal{S}'| \geq 3$. Let $T = (\bigcup_{\alpha \in \mathcal{S}'} X_\alpha) \cup 3L_Q$ where $3L_Q$ is the set consisting of points $P \in \mathcal{X} \setminus \{Q\}$ which lie on 3-lines through Q . Since $3L_Q \cap X_\alpha$ is empty for all $\alpha \in \mathcal{S}'$, this union is disjoint. Hence $|T| \geq 2y_Q + |\mathcal{S}'|x_Q$. The bound on $|\mathcal{X}|$ together with (2) now implies that $|\mathcal{S}'| < 4$.

Thus $|\mathcal{S}'| = 3$ and we must have $2y_Q + 3x_Q \leq |\mathcal{X}| \leq 2n + 2 - \frac{2n}{p}$. By (2), we also have $2y_Q + 4x_Q \geq 2n + 4 + \frac{4n}{p}$. Hence $x_Q \geq 2 + \frac{6n}{p}$ so that (4) implies that $|\mathcal{S}| = 2r$ with $1 < r < \frac{p}{6}$. Also, G_Q must contain a loop as the number of vertices in G_Q of degree 1 is odd. Thus the graph G_Q consists of two components. One is

$$1(-1)2(-2)\cdots t(-t) \quad \text{for some } t \text{ such that } 1 \leq t < r,$$

and the other is

$$\left(\frac{p+1}{2}\right)\left(\frac{p+1}{2}\right)\left(\frac{p-1}{2}\right)\left(\frac{3-p}{2}\right)\left(\frac{p-3}{2}\right)\cdots\left(\frac{p+1-(2r-2t)}{2}\right).$$

Thus 1, $-t$ and $\alpha = \frac{p+1-(2r-2t)}{2} = \frac{p+1}{2} - (r-t)$ are the three vertices of degree one in G_Q .

Let $T = \{Q\} \cup 2L_Q \cup 3L_Q \cup X_{-t} \cup X_\alpha$, where iL_Q denotes the set of points in $\mathcal{X} \setminus \{Q\}$ which lie on i -lines through Q . Clearly, $|T| \geq 2y_Q + 3x_Q$. If every 4-line through Q contains a point of $\mathcal{X} \setminus T$, then we get

$$|\mathcal{X}| \geq z_Q + |T| \geq z_Q + 2y_Q + 3x_Q.$$

Since this, by (1), contradicts the bound on $|\mathcal{X}|$, we must have a 4-line L through Q contained in T . As $2L_Q$ and $3L_Q$ consist of the points on 2 and 3-lines through Q , respectively, $L \subset \{Q\} \cup X_{-t} \cup X_\alpha$.

Let L contain i points from X_α , where $0 \leq i \leq 3$. Then the sum of the colours occurring on L is $-1 + i\alpha - (3-i)t$ which must be $0 \pmod{p}$. Substituting the value of α , for i varying from 0 to 3, we infer that one of the following four integers is a multiple of p :

$$3t + 1, \quad 2(r+t) + 1, \quad 2r - t, \quad 6(r-t) - 1.$$

However, as $1 \leq t < r < \frac{p}{6}$, none of the above integers can be a multiple of p . This completes the proof of the theorem.

THEOREM 2. *Let $(\mathcal{X}, \mathcal{L})$ be a partial linear space whose every point is contained in at least $n+1$ lines. If it is non-trivial at p , then the minimum weight of $\mathcal{C}_{\mathcal{L}}^{\perp}$ is at least $2n+2-\frac{2^n}{p}$. If this weight is attained, then the induced structure on the support of a word of weight $2n+2-\frac{2^n}{p}$ in $\mathcal{C}_{\mathcal{L}}^{\perp}$ is isomorphic to the join of two Steiner 2-designs with $\frac{n}{p}$ lines through every point and p points on every line.*

Proof. The induced structure on the support of a non-zero word w of $\mathcal{C}_{\mathcal{L}}^{\perp}$ is a partial linear space with $\geq n+1$ lines through every point that is non-trivial at p . Theorem 1 now implies the theorem.

THEOREM 3. *A projective plane of order p is tame.*

Proof. This is the special case of the above theorem when $n = p$ and $(\mathcal{X}, \mathcal{L})$ is a projective plane.

A partial linear space of order n is a partial linear space with exactly $n+1$ lines through every point.

Remark 1. If a Steiner 2-design with order $\frac{n}{p}-1$ and line length p is to exist, then the standard divisibility properties of its parameters imply that $n \equiv 0$ or $p \pmod{p^2}$. Therefore, in case n does not satisfy this requirement, the minimum weight of $\mathcal{C}_{\mathcal{L}}^{\perp}$ is more than $2n+2-\frac{2^n}{p}$. The join of a projective plane of order 2 with a set S of cardinality 4 (with no lines!) is a partial linear space of order $n = 6$. With $p = 3$, its cardinality is $2n+3-\frac{2^n}{p} = 11$ and it is non-trivial at p .

Remark 2. By arguments similar to those in the proof of Theorem 1, we can prove that the minimum weight of $\mathcal{C}_{\mathcal{L}}^{\perp}$ corresponding to a linear space of order $2p$ is at least $4p$ and can obtain an explicit description of words of weight $4p$ in $\mathcal{C}_{\mathcal{L}}^{\perp}$ if they exist. One consequence of our result is that a projective plane of order $2p$, if it exists, must be tame at p . (The reader interested in knowing the detailed proof may contact the author.)

Remark 3. The join of an affine plane of order p and a set S of cardinality $p(p-1)$ with one line $L = S$ is an example of a linear space of order p^2 which is non-trivial at p and whose cardinality is strictly less than $2p^2$.

The results that we have obtained so far prompt us to believe in:

Conjecture 1. The minimum weight of the p -ary code $\mathcal{C}_{\mathcal{L}}^{\perp}$ corresponding to a linear space of order $n < p^2$ is at least $2n$.

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