

# INDIAN STATISTICAL INSTITUTE

## End Semester Examination

M.Tech CS

*Computational Complexity*

Date: 24 November 2025

Maximum Marks: 50

Duration: 3 hours

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**General comment.** Answer as much as you can, but the maximum you can score is 50.

### **Part A ( $5 \times 5 = 25$ marks)**

1. (a) Define the classes **P**, **NP**, and **coNP**.  
(b) Give an example of a problem in **NP** that is *not known* to be in **P** or **NP**-complete.  
(c) State whether the following inclusions are known to be true or false (justify briefly):
    - i.  $P \subseteq NP$
    - ii.  $NP \subseteq PSPACE$
    - iii.  $P = NP$
  2. Explain what it means for a problem  $A$  to be *polynomial-time reducible* to a problem  $B$ . Give an example of one **NP**-complete problem and briefly describe how another problem can be reduced to it.
  3. Explain in your own words how the **Time Hierarchy Theorem** is proved using diagonalization. What is the main idea behind the construction?
  4. What does it mean for a proof technique to *relativize*? Give an example showing that relativization cannot resolve the  $P$  vs  $NP$  question.
  5. (a) Define the *deterministic decision tree complexity* of a Boolean function.  
(b) What is the deterministic decision tree complexity of the **OR** function on  $n$  variables? Justify briefly.
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## Part B (5 × 12 = 60 marks)

1. Show that **3SAT** is polynomial-time reducible to **CLIQUE**. Clearly describe the reduction and argue correctness. (You may assume 3SAT is NP-complete.)
2. State the **Time Hierarchy Theorem** formally. Explain why it implies that  $P \subsetneq EXP$ . What would happen if the theorem did not hold? Discuss briefly.
3. (a) State (without proof) the **PCP theorem**.  
(b) Explain in simple terms what it means for NP to have a PCP verifier.  
(c) Mention one important implication of the PCP theorem in the context of *hardness of approximation*.
4. Suppose there exists an oracle  $A$  such that  $P^A = NP^A$ , and another oracle  $B$  such that  $P^B \neq NP^B$ .
  - (a) What does this tell us about the limitations of relativizing proof techniques?
  - (b) Which of the following techniques are **non-relativizing**: diagonalization, arithmetization, or probabilistic checking?
  - (c) Why was the discovery of the PCP theorem considered a breakthrough in non-relativizing techniques?
5. (a) Define the **deterministic communication complexity** of a function  $f(x, y)$ .  
(b) For the **Equality** function

$$EQ(x, y) = \begin{cases} 1 & \text{if } x = y, \\ 0 & \text{otherwise,} \end{cases}$$

where  $x, y \in \{0, 1\}^n$ , prove that its deterministic communication complexity is at least  $n + 1$  bits. (Hint: Use the fooling set argument.)